

It ain't necessarily so: Basic sequent systems for negative modalities

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Abstract

We look at non-classical negations and their corresponding adjustment connectives from a modal viewpoint, over complete distributive lattices, and apply a very general mechanism in order to offer adequate analytic proof systems to logics that are based on them. Defining non-classical negations within usual modal semantics automatically allows one to treat equivalent formulas as synonymous, and to have a natural justification for a global version of the contraposition rule. From that perspective, our study offers a particularly useful environment in which negative modalities and their companions may be used for dealing with inconsistency and indeterminacy. After investigating modal logics based on arbitrary frames, we extend the results to serial frames, reflexive frames, functional frames, and symmetric frames. In each case we also investigate when and how classical negation may thereby be defined.

Keywords: negative modalities, sequent systems, cut-admissibility, analyticity.

1 Capturing the impossible, and its dual

Many well-known subclassical logics—including intuitionistic logic and several many-valued logics—share the conjunction-disjunction fragment of classical logic, but disagree about the exact notion of opposition and the specific logical features to be embodied in *negation*. In contrast, modal logics are often thought of as superclassical, and are obtained by the addition of *identity*-like ‘positive modalities’ \Box and \Diamond . For various well-known cases, such modalities fail to have a finite-valued characterization. Notwithstanding, each m -ary connective \Diamond of a modal logic is typically *congruential* (with respect to the underlying consequence relation \vdash), in treating equivalent formulas as synonymous: if $\alpha_i \vdash \beta_i$ and $\beta_i \vdash \alpha_i$, for every $1 \leq i \leq m$, then $\Diamond(\alpha_1, \dots, \alpha_m) \vdash \Diamond(\beta_1, \dots, \beta_m)$. To

logical systems containing only such sort of connectives one might associate semantics in terms of neighborhood frames (see ch.5 of [23]), and the same applies if one uses 1-ary ‘negative modalities’ instead, as in [19]. But normal modal logics make their 1-ary positive modalities respect indeed a stronger property: if $\alpha \vdash \beta$ then $\diamond(\alpha) \vdash \diamond(\beta)$. Such monotone behavior may be captured by semantics based on Kripke frames, and the same applies to the antitone behavior that characterize negative modalities, namely: if $\alpha \vdash \beta$ then $\diamond(\beta) \vdash \diamond(\alpha)$.

In [7] an investigation of negative modalities is accomplished on top of the $\wedge\vee\top\perp$ -fragment of classical logic, and the same base language had already been considered in [18] for the combination of positive and negative modalities. Typically, in studies of positive and negative modalities the so-called compatibility (bi-relational) frames are used, and certain appropriate conditions upon the commutativity of diagrams involving their two relations are imposed, having as effect the heredity of truth (i.e., its persistence towards the future) with respect to one of the mentioned relations (assumed to be a partial order).

There are a number of studies (e.g. [21,6]) in which the above mentioned languages for dealing with negative modalities are upgraded in order to count on an (intuitionistic or classical) implication, and sometimes also its dual, co-implication (cf. [17]). If one may count on classical implication, however, it suffices to add to it the modal paraconsistent negation given by ‘unnecessity’ (cf. [13]), and all other connectives of normal modal logics turn out to be definable from such impoverished basis (indeed, where \smile is a primitive symbol for unnecessity and \rightarrow represents classical implication, we have that $\sim\alpha := \alpha \rightarrow \smile(\alpha \rightarrow \alpha)$ behaves as the classical negation of α , and $\Box\alpha := \smile\smile\alpha$ behaves as the usual positive modality box).

Our intuition about the relation between a paracomplete (a.k.a. ‘intuitionistic-like’) negation and a paraconsistent negation is that the former would be expected to be more demanding than the latter, while classical negation should sit between the two (whenever it also turns out to be expressible). It takes indeed more effort to assert a negated statement constructively, while such statements are more readily asserted should some contradictions be allowed to subsist; in other words, negations in a paracomplete logic come at a greater cost than classical negations, while paraconsistent logics indulge on negations in which classical logic would show greater restraint. The presence of a classical negation, however, often makes it too easy to forget that there are two distinct kinds of deviations equally worth studying, concerning non-classical negation, as one of these deviations may then be recovered in the standard way as the dual of the other. In order to get a better grasp of the duality between paraconsistent and paracomplete modal negations (namely, unnecessity vs. impossibility), we purposefully make an effort to prevent the underlying language from being sufficiently expressive so as to allow for the definition of a classical negation (or a classical implication) — whenever that goal lies within reach. Here we do however in all cases enrich our object language with certain ‘adjustment connectives’ expressing negation-consistency and negation-determinacy, allowing for the simulation of usual features of classical negation and for the

(partial) recovery of classical reasoning. It should be noted, however, that as a byproduct of the presence of such adjustment connectives truth will no longer be hereditary in our Kripke models, that is, it will not in general be preserved for all compound formulas towards the future, in contrast with what happens with models of compatibility frames.

In what follows, first and foremost we will concentrate on the logic PK , determined by the class of all Kripke frames, which has been introduced and received a presentation as a sequent system in [5]. We show here that it can be reintroduced in terms of a so-called ‘basic sequent system’, which allows one to take advantage of general techniques developed in [10], including a method for obtaining sound and complete Kripke semantics and a uniform recipe for semantic proofs of cut-admissibility or analyticity. The next section adopts a semantical perspective to explain why and how our study is done.

2 On negative modalities

We briefly recall the now familiar elements of a Kripke semantics. A *frame* is a structure consisting of a nonempty set W (of ‘worlds’) and a binary (‘accessibility’) relation R on W . A *model* $\mathcal{M} = \langle \mathcal{F}, V \rangle$ is based on a frame $\mathcal{F} = \langle W, R \rangle$ and on a *valuation* $V : W \times \mathcal{L} \rightarrow \{f, t\}$ that assigns truth-values to worlds $w \in W$ and sentences φ of a propositional language \mathcal{L} generated over a denumerable set of propositional variables \mathcal{P} . The valuations must satisfy certain conditions that are induced by the fixed interpretation of the connectives of the language. When $V(w, \varphi) = t$ we say that V *satisfies* φ at w , and denote this by $\mathcal{M}, w \Vdash \varphi$; otherwise we write $\mathcal{M}, w \not\Vdash \varphi$ and say that V *leaves* φ *unsatisfied* at w . The connectives from the positive fragment of classical logic receive their standard boolean interpretations locally, world-wise, by recursively setting:

$$\begin{aligned} [\text{ST}] \quad & \mathcal{M}, w \Vdash \top \\ [\text{S}\wedge] \quad & \mathcal{M}, w \Vdash \varphi \wedge \psi \quad \text{iff} \quad \mathcal{M}, w \Vdash \varphi \text{ and } \mathcal{M}, w \Vdash \psi \\ [\text{S}\vee] \quad & \mathcal{M}, w \not\Vdash \varphi \vee \psi \quad \text{iff} \quad \mathcal{M}, w \not\Vdash \varphi \text{ and } \mathcal{M}, w \not\Vdash \psi \end{aligned}$$

Given formulas $\Gamma \cup \Delta$ of \mathcal{L} , and given a class of frames \mathcal{E} , we say that Γ *entails* Δ *in* \mathcal{E} , and denote this by $\Gamma \models_{\mathcal{E}} \Delta$, if for each model \mathcal{M} based on a frame $\mathcal{F} \in \mathcal{E}$ and each world w of \mathcal{M} we have either $\mathcal{M}, w \not\Vdash \gamma$ for some $\gamma \in \Gamma$ or $\mathcal{M}, w \Vdash \delta$ for some $\delta \in \Delta$. The assertion $\Gamma \models_{\mathcal{E}} \Delta$ will be called a *consecution*. As usual, in what follows we will focus most of the time on consecutions $\Gamma \models_{\mathcal{E}} \Delta$ involving a singleton Δ , and in the next section we will extend the notion of entailment so as to cover sequents instead of formulas. The subscript \mathcal{E} shall be omitted in what follows whenever there is no risk of ambiguity.

In the following subsections we extend the above language with connectives whose modal interpretations will be useful for the investigation of non-classical negations.

2.1 Adding negations

Our first extension of the above language proceeds by the addition of a 1-ary connective \neg , to be interpreted non-locally as follows:

$[S\sim]$ $\mathcal{M}, w \Vdash \sim\varphi$ iff $\mathcal{M}, v \nVdash \varphi$ for some $v \in W$ such that wRv

Accordingly, a formula $\sim\varphi$ is said to be satisfied at a given world of a model precisely when the formula φ fails to be satisfied at some world accessible from this given world. In the following paragraph we will show that \sim respects some minimal conditions to deserve being called a ‘negation’, namely, we will demonstrate its ability to invert truth-values assigned to certain formulas (at certain worlds).

Let $\#$ represent an arbitrary 1-ary connective, and let $\#^j$ abbreviate a j -long sequence of $\#$'s. The least we will demand from $\#$ to call it a *negation* is that, for every $p \in \mathcal{P}$ and every $k \in \mathbb{N}$:

$[[\textit{falsificatio}]]$ $\#^k p \nVdash \#^{k+1} p$ $[[\textit{verificatio}]]$ $\#^{k+1} p \nVdash \#^k p$

To witness $[[\textit{falsificatio}]]$, some sentence φ is to be satisfied while the sentence $\#\varphi$ is not simultaneously satisfied; for $[[\textit{verificatio}]]$ some sentence φ is left unsatisfied while at the same time $\#\varphi$ is satisfied. To check that the connective \sim fulfills such requisites, it suffices for instance to build a frame in which $W = \{w_n : n \in \mathbb{N}\}$ and wRv iff $v = w^{++}$ (namely, v is the successor of w), and consider a valuation V such that $V(w_n, p) = t$ iff n is odd.

It is very easy to see that our connective \sim satisfies *global contraposition* in the sense that $\alpha \models \beta$ implies $\sim\beta \models \sim\alpha$. Indeed, assume $\alpha \models \beta$ and suppose that $\mathcal{M}, w \Vdash \sim\beta$ for some world w of an arbitrary model \mathcal{M} . Then, $[S\sim]$ informs us that there must be some world v in \mathcal{M} such that wRv and $\mathcal{M}, v \nVdash \beta$. By the definition of entailment, the initial assumption gives us $\mathcal{M}, v \nVdash \alpha$. Using again $[S\sim]$ we conclude that $\mathcal{M}, w \Vdash \sim\alpha$. As a byproduct of this, if one defines an equivalence relation \equiv on \mathcal{L} by setting $\alpha \equiv \beta$ whenever both $\alpha \models \beta$ and $\beta \models \alpha$, then an easy structural induction on \mathcal{L} establishes that \equiv is not only compatible with \sim but also with the other connectives that are used in constructing the algebra of formulas; in other words, \equiv constitutes a congruence relation on \mathcal{L} .

It is straightforward to see that any 1-ary connective $\#$ satisfying global contraposition is such that, given $p, q \in \mathcal{P}$:

(DM1.1 $\#$) $\#(p \vee q) \models \#p \wedge \#q$ (DM2.1 $\#$) $\#p \vee \#q \models \#(p \wedge q)$

If $\#$ also respects the following consecutions, then it is said to be a *full type diamond-minus connective*:

(DM2.2 $\#$) $\#(p \wedge q) \models \#p \vee \#q$ (DT $\#$) $\#\top \models p$

Note that \sim is a full type diamond-minus connective. To check that \sim satisfies (DM2.2 $\#$), indeed, suppose that $\mathcal{M}, w \Vdash \sim(p \wedge q)$ for some arbitrary world w of an arbitrary model \mathcal{M} . By $[S\sim]$ we know that there is some world v such that wRv and $\mathcal{M}, v \nVdash p \wedge q$. It follows by $[S\wedge]$ that $\mathcal{M}, v \nVdash p$ or $\mathcal{M}, v \nVdash q$. Using $[S\sim]$ again we conclude that $\mathcal{M}, w \Vdash \sim p$ or $\mathcal{M}, w \Vdash \sim q$ and $[S\vee]$ gives us $\mathcal{M}, w \Vdash \sim p \vee \sim q$. In addition, to check that \sim satisfies (DT $\#$) one may invoke $[S\sim]$ and $[S\top]$. Note that satisfying (DT $\#$) means that the nullary connective \perp taken as an abbreviation of $\#\top$ is interpretable by setting, for every world w of every model \mathcal{M} :

$[S\perp]$ $\mathcal{M}, w \nVdash \perp$

Given a negation $\#$, we call the logic containing it $\#$ -*paraconsistent* if the following consecution fails, for $p, q \in \mathcal{P}$:

$$\llbracket \# \text{-explosion} \rrbracket \quad p, \#p \vdash q$$

This means that there must be valuations that satisfy both some sentence φ and the sentence $\#\varphi$ while not satisfying every other sentence. It is worth noticing that $\llbracket \sim \text{-explosion} \rrbracket$ holds good in frames containing exclusively worlds that are accessible to themselves, and themselves only (call such worlds ‘narcissistic’) and worlds that do not access any other world (call them ‘dead ends’): in the former case, it is impossible to simultaneously satisfy both φ and $\sim\varphi$; in the latter case, the sentence $\sim\varphi$ is never satisfied. Note moreover that in the class of all narcissistic frames the connective \sim happens to behave like classical negation, i.e., it behaves like the symbol \sim in the following semantic clause:

$$[S\sim] \quad \mathcal{M}, w \Vdash \sim\varphi \quad \text{iff} \quad \mathcal{M}, w \nVdash \varphi$$

In contrast, in the class of all frames whose worlds are all dead ends the connective \sim does not respect [verification], and cannot be said thus to be a negation.

We now make a further extension of the above language by adding a 1-ary connective \frown , non-locally interpreted as follows:

$$[S\frown] \quad \mathcal{M}, w \Vdash \frown\varphi \quad \text{iff} \quad \mathcal{M}, v \nVdash \varphi \text{ for every } v \in W \text{ such that } wRv$$

It is not difficult to check that again we have a connective that qualifies as a negation, and satisfies global contraposition. To reinforce the meta-theoretical duality between the latter negation and the negation introduced above through $[S\sim]$, we will henceforth refer to the previous interpretation clause in the following equivalent form:

$$[S\sim] \quad \mathcal{M}, w \nVdash \frown\varphi \quad \text{iff} \quad \mathcal{M}, v \Vdash \varphi \text{ for some } v \in W \text{ such that } wRv$$

A *full type box-minus connective* is a 1-ary connective $\#$ that respects:

$$(DM1.2\#) \quad \#p \wedge \#q \vdash \#(p \vee q) \quad (DF\#) \quad p \vdash \#\perp$$

One may easily check that \frown is indeed a full type box-minus connective.

Given a negation $\#$, we call the logic containing it $\#$ -*paracomplete* if it fails the following consecution, for $p, q \in \mathcal{P}$:

$$\llbracket \# \text{-implosion} \rrbracket \quad q \vdash \#p, p$$

Such failure will clearly be the case for $\# = \frown$ as soon as we entertain frames that contain worlds that are neither dead ends nor narcissistic. Otherwise, we see that \frown will behave either like classical negation (if all worlds are narcissistic) or like \top (if all worlds are dead ends).

In the following sections, unless noted otherwise, we will no longer consider classes of frames containing only frames with worlds that are either dead ends or narcissistic — so we will only consider entailment relations that are \sim -paraconsistent and \frown -paracomplete, for the negative modalities \sim (assumed to be full-type diamond-minus) and \frown (assumed to be full-type box-minus).

2.2 Recovering negation-consistency and negation-determinacy

In what follows we will call a model *dadaistic* when it contains some world in which all formulas are satisfied, and call it *nihilistic* if it leaves all formulas unsatisfied at some world. It is straightforward to see that the language based on $\wedge\vee\top\sim$, with the above interpretations, admits dadaistic models, while the language based on $\wedge\vee\perp\sim$ admits nihilistic models.

Recall that a $\#$ -paraconsistent logic allows for valuations that satisfy certain formulas φ and $\#\varphi$ while leaving some other formula ψ unsatisfied (at some fixed world). There might be reasons for disallowing this phenomenon to occur with an arbitrary φ , or for restricting to certain formulas ψ but not others. A particularly useful way of keeping a finer control over which ‘inconsistencies’ of the form φ and $\#\varphi$ are to be acceptable within non-dadaistic models is to mark down the formula thereby involved so as to recover a ‘gentle’ version of $\llbracket\#\text{-explosion}\rrbracket$. Concretely, for us here, a 1-ary connective \oplus that strongly internalizes the meta-theoretic ‘consistency assumption’ at the object language level will be such that:

$$[\text{SC}\#] \quad \mathcal{M}, w \Vdash \oplus\varphi \quad \text{iff} \quad \mathcal{M}, w \nVdash \varphi \text{ or } \mathcal{M}, w \nVdash \#\varphi$$

It is easy to check that any connective \oplus respecting $[\text{SC}\#]$ is such that:

$$(\text{C1}\#) \quad \oplus p, p, \#p \models \quad (\text{C2}\#) \quad \models p, \oplus p \quad (\text{C3}\#) \quad \models \#p, \oplus p$$

Note in particular that (C1#) guarantees that there are no valuations that satisfy (at a fixed world) both p and $\#p$ if these are put in the presence of $\oplus p$. Thus, in case $\#$ fails $\llbracket\#\text{-explosion}\rrbracket$ we may look at the latter formula involving \oplus as guaranteeing that a weaker form of explosion is available. On these grounds we shall call the connective \oplus an *adjustment* companion to $\#$: it allows one to recover explosion from within a non- $\#\text{-explosive}$ (i.e., paraconsistent) logical context, and adjust the consecutions of the underlying logic so as to allow for the simulation of the consecutions that would otherwise be justified by reference to $\llbracket\#\text{-explosion}\rrbracket$. Semantically, the presence of such connective also guarantees that dadaistic models are not admissible over the language based on $\wedge\vee\top\sim\ominus$, with the above interpretations. This is because a formula of the form $\ominus\varphi \wedge (\varphi \wedge \sim\varphi)$ is equivalent to a formula \perp respecting $[\text{S}\perp]$.

Dually, a $\#$ -paracomplete logic allows for valuations that leave the formulas φ and $\#\varphi$ both unsatisfied (at some fixed world), while satisfying some other formula ψ . A particular way of keeping a finer control over which ‘indeterminacies’ of the form φ and $\#\varphi$ are to be acceptable within non-nihilistic models is to allow for a ‘gentle’ version of $\llbracket\#\text{-implosion}\rrbracket$, where a 1-ary connective \oplus internalizes the meta-theoretic ‘determinacy assumption’ at the object language level, in such a way that:

$$[\text{SD}\#] \quad \mathcal{M}, w \nVdash \oplus\varphi \quad \text{iff} \quad \mathcal{M}, w \Vdash \varphi \text{ or } \mathcal{M}, w \Vdash \#\varphi$$

Clearly, any connective \oplus respecting $[\text{SD}\#]$ is such that:

$$(\text{D1}\#) \quad \models \#p, p, \oplus p \quad (\text{D2}\#) \quad \oplus p, p \models \quad (\text{D3}\#) \quad \oplus p, \#p \models$$

Note that a formula of the form $(\varphi \vee \sim\varphi) \vee \ominus\varphi$ is equivalent to a formula \top respecting $[\text{S}\top]$. Note, moreover, that whenever it turns out that a connective $\#$

respects $\llbracket \# \text{-explosion} \rrbracket$ and at the same time its adjustment companion \oplus respects $\llbracket \text{SC}\# \rrbracket$, then the formula $\oplus\varphi$ is equivalent to \top . In an analogous way, whenever a connective $\#$ respects $\llbracket \# \text{-implosion} \rrbracket$ and at the same time its adjustment companion \oplus respects $\llbracket \text{SD}\# \rrbracket$, the formula $\oplus\varphi$ is equivalent to \perp . This stresses the fact that the adjustment connectives with which we deal in this subsection are more interesting when they accompany the respective non-classical negations to whose meaning they contribute.

At this point we have finally finished constructing the richest language that will be used throughout the rest of the paper: It will contain the connectives $\wedge\vee\top\perp\sim\ominus\wedge\odot$, disciplined by the $\llbracket \text{S}\# \rrbracket$ conditions above. In the following subsection we will explain precisely when a classical negation, that is a 1-ary connective \sim subject to condition $\llbracket \text{S}\sim \rrbracket$, is definable with the use of our language. Fixed such language, the logic characterized over it by the class \mathcal{E} of all frames will be called *PK*; the logic characterized by the class $\mathcal{E}_{\mathbf{D}}$ of all frames with serial accessibility relations will be called *PKD*; the logic characterized by the class $\mathcal{E}_{\mathbf{T}}$ of all frames with reflexive accessibility relations will be called *PKT*; the logic characterized by the class $\mathcal{E}_{\mathbf{Fun}}$ of all frames whose accessibility relations are total functions will be called *PKF*; the logic characterized by the class $\mathcal{E}_{\mathbf{B}}$ of all symmetric frames (those with symmetric accessibility relations) will be called *PKB*.

2.3 Around classical negation

According to the intuitions laid down at Section 1, one could expect that in general (a) $\neg\alpha \vdash \sim\alpha$ and (b) $\sim\alpha \vdash \neg\alpha$. It is easy to see that these consecutions are sanctioned by *PKT*, for the classical negation \sim that may be defined by setting $\sim\varphi := \neg\varphi \wedge \ominus\varphi$ (alternatively, one may set $\sim\varphi := \neg\varphi \vee \odot\varphi$).

Meanwhile, in the deductively weaker logic *PKD* one cannot in general prove (a) nor (b), even though a classical negation may be defined in this logic by setting $\sim\varphi := (\neg\varphi \wedge \ominus\varphi) \vee \odot\varphi$. However, one can still easily prove in *PKD* that (c) $\neg\alpha \vdash \sim\alpha$. In the logic *PKF*, deductively stronger than *PKD* (but neither stronger nor weaker than *PKT*) one may also prove the converse consecution, (d) $\sim\alpha \vdash \neg\alpha$. Indeed, suppose $\mathcal{M}, w \Vdash \sim\alpha$. There is, by the fact that the accessibility relation is a total function, a single world v such that wRv . Then $\mathcal{M}, v \not\Vdash \alpha$, by $\llbracket \text{S}\sim \rrbracket$. For a similar reason, invoking now $\llbracket \text{S}\neg \rrbracket$ we conclude that $\mathcal{M}, w \Vdash \neg\alpha$. Note that (c) and (d) together make our two modal non-classical negations indistinguishable from the viewpoint of *PKF*, yet there would still be no reason for them to collapse into classical negation.

The situation concerning classical negation and its relation to its non-classical neighbours gets even more interesting if one acknowledges that *no* classical negation is definable in *PK*, the weakest of our logics, but also no classical negation is definable in the fragment of *PKT* without neither of the adjustment connectives, or in the fragment of *PKF* (or *PKD*) without either one of the adjustment connectives, or in *PKB*. Detailed proofs concerning the mentioned results about (non)definability of classical negation in the weak modal logics that constitute our present object of study may be found in Section 6.

Notice that in PKD and its extensions there are no negated formulas that happen to be true or false at a given world just because there are no worlds accessible from it. Note also that the logic PKT is: paraconsistent but not paracomplete with respect to the connective \sim ; paracomplete but not paraconsistent with respect to \wedge (even though we will not prove it here, this logic is indeed the least extension of the positive implicationless fragment of classical logic with the latter mentioned properties). The logic PKT will have its word in the following sections, for it also allows for the straightforward application of the techniques that will be hereby illustrated. In the other four mentioned logics, in contrast, both non-classical negations behave at once as paracomplete and paraconsistent negations (recall, though, that each is associated to a different adjustment connective). We take the cases among these in which no classical negation is available to be particularly attractive for the task of revealing the ‘uncontaminated’ nature of non-classical negation. Establishing well-behaved proof theoretical counterparts for such logics, as we shall do in what follows, is meant to allow for them to be better understood and dealt with.

3 A proof system for PK

A sequent calculus for PK , that we denote by PK , was introduced in [5], and consists of the following rules:

$$\begin{array}{ll}
[id] \frac{}{\Gamma, \varphi \Rightarrow \varphi, \Delta} & [cut] \frac{\Gamma, \varphi \Rightarrow \Delta \quad \Gamma \Rightarrow \varphi, \Delta}{\Gamma \Rightarrow \Delta} \\
[W \Rightarrow] \frac{\Gamma \Rightarrow \Delta}{\Gamma, \varphi \Rightarrow \Delta} & [\Rightarrow W] \frac{\Gamma \Rightarrow \Delta}{\Gamma \Rightarrow \varphi, \Delta} \\
[\perp \Rightarrow] \frac{}{\Gamma, \perp \Rightarrow \Delta} & [\Rightarrow \top] \frac{}{\Gamma \Rightarrow \top, \Delta} \\
[\wedge \Rightarrow] \frac{\Gamma, \varphi, \psi \Rightarrow \Delta}{\Gamma, \varphi \wedge \psi \Rightarrow \Delta} & [\Rightarrow \wedge] \frac{\Gamma \Rightarrow \varphi, \Delta \quad \Gamma \Rightarrow \psi, \Delta}{\Gamma \Rightarrow \varphi \wedge \psi, \Delta} \\
[\vee \Rightarrow] \frac{\Gamma, \varphi \Rightarrow \Delta \quad \Gamma, \psi \Rightarrow \Delta}{\Gamma, \varphi \vee \psi \Rightarrow \Delta} & [\Rightarrow \vee] \frac{\Gamma \Rightarrow \varphi, \psi, \Delta}{\Gamma \Rightarrow \varphi \vee \psi, \Delta} \\
[\sim \Rightarrow] \frac{\Gamma \Rightarrow \varphi, \Delta}{\wedge \Delta, \sim \varphi \Rightarrow \sim \Gamma} & [\Rightarrow \sim] \frac{\Gamma, \varphi \Rightarrow \Delta}{\wedge \Delta \Rightarrow \wedge \varphi, \sim \Gamma} \\
[\odot \Rightarrow] \frac{\Gamma \Rightarrow \varphi, \Delta \quad \Gamma \Rightarrow \sim \varphi, \Delta}{\Gamma, \odot \varphi \Rightarrow \Delta} & [\Rightarrow \odot] \frac{\Gamma, \varphi, \sim \varphi \Rightarrow \Delta}{\Gamma \Rightarrow \odot \varphi, \Delta} \\
[\oslash \Rightarrow] \frac{\Gamma \Rightarrow \varphi, \wedge \varphi, \Delta}{\Gamma, \oslash \varphi \Rightarrow \Delta} & [\Rightarrow \oslash] \frac{\Gamma, \varphi \Rightarrow \Delta \quad \Gamma, \wedge \varphi \Rightarrow \Delta}{\Gamma \Rightarrow \oslash \varphi, \Delta}
\end{array}$$

Above, sequents are taken to have the form $\Sigma \Rightarrow \Pi$ where Σ and Π are finite sets of formulas, and given a unary connective $\#$ and $\Psi \subseteq \mathcal{L}$, by $\#\Psi$ we denote the set $\{\#\psi \mid \psi \in \Psi\}$. We write $S \vdash_{PK} s$ to say that there is a derivation in PK of a sequent s from a set S of sequents. That establishes a consequence relation between sequents. A consequence relation between formulas is defined by setting $\Gamma \vdash_{PK} \varphi$ if $\vdash_{PK} \Gamma' \Rightarrow \varphi$ for some finite subset Γ' of Γ . The overloaded

notation \vdash_{PK} will always be resolved by the pertinent context.

Next, we utilize in what follows the general mechanisms and methods applicable to the so-called ‘basic systems’ of [10] in order to prove soundness, completeness and cut-admissibility. From the viewpoint of basic systems, each sequent is seen as a union of a ‘main sequent’ and a ‘context sequent’. For example, in $[\Rightarrow\vee]$, the main sequent of the premise is $\Rightarrow \varphi, \psi$; the main sequent of the conclusion is $\Rightarrow \varphi \vee \psi$; and the context sequent of both is $\Gamma \Rightarrow \Delta$. Note that in the rules for \smile and \frown , the context sequent of the premise is different from the one of the conclusion. Accordingly, [10] introduces the notion of a *basic rule*, whose premises take the form $\langle s, \pi \rangle$, where s is a sequent that corresponds to the main sequent of the premise, and π is a relation between singleton-sequents (that is, sequents of the form $\varphi \Rightarrow$ or $\Rightarrow \varphi$) called a *context relation* that determines the behavior of the context sequents. The sequent calculus PK may be naturally regarded as a basic system that employs two context relations, namely: $\pi_0 = \{\langle p_1 \Rightarrow; p_1 \Rightarrow \rangle, \langle \Rightarrow p_1; \Rightarrow p_1 \rangle\}$, and $\pi_1 = \{\langle p_1 \Rightarrow; \Rightarrow \smile p_1 \rangle, \langle \Rightarrow p_1; \frown p_1 \Rightarrow \rangle\}$. The rules of PK may then be presented as particular instances of basic rules. For example, the following are the basic rules for \wedge, \smile, \frown and \ominus :

$$\begin{array}{ll} [\Rightarrow\wedge] \langle \Rightarrow p_1; \pi_0 \rangle, \langle \Rightarrow p_2; \pi_0 \rangle / \Rightarrow p_1 \wedge p_2 & [\wedge\Rightarrow] \langle p_1, p_2 \Rightarrow; \pi_0 \rangle / p_1 \wedge p_2 \Rightarrow \\ [\smile\Rightarrow] \langle \Rightarrow p_1; \pi_1 \rangle / \smile p_1 \Rightarrow & [\Rightarrow\smile] \langle p_1 \Rightarrow; \pi_1 \rangle / \Rightarrow \smile p_1 \\ [\ominus\Rightarrow] \langle \Rightarrow p_1; \pi_0 \rangle, \langle \Rightarrow \smile p_1; \pi_0 \rangle / \ominus p_1 \Rightarrow & [\Rightarrow\ominus] \langle p_1, \smile p_1 \Rightarrow; \pi_0 \rangle / \Rightarrow \ominus p_1 \end{array}$$

In applications of $[\Rightarrow\ominus]$, the context sequent is left unchanged, as two singleton-sequents relate to each other (with respect to π_0) iff they are the same. In contrast, applications of $[\smile\Rightarrow]$ are based on π_1 . A sequent $\Gamma_1 \Rightarrow \Delta_1$ relates (with respect to π_1) to a sequent $\Gamma_2 \Rightarrow \Delta_2$ iff $\Gamma_2 = \frown \Delta_1$ and $\Delta_2 = \smile \Gamma_1$.

We extend the notion of satisfaction from Section 2 to sequents by setting $\mathcal{M}, w \Vdash \Gamma \Rightarrow \Delta$ if $\mathcal{M}, w \not\models \gamma$ for some $\gamma \in \Gamma$ or $\mathcal{M}, w \Vdash \delta$ for some $\delta \in \Delta$. Semantics for PK may then be obtained using the general method introduced in [10], by having each derivation rule and each context relation match a semantic condition, and the semantics of the system is obtained by conjoining all these semantic conditions. For example, the basic rule $[\smile\Rightarrow]$ induces the condition: “if $\mathcal{M}, v \Vdash \Rightarrow \varphi$ for every world v such that wRv , then $\mathcal{M}, w \Vdash \smile \varphi \Rightarrow$ ”, which is equivalent to: “If $\mathcal{M}, w \Vdash \smile \varphi$ then $\mathcal{M}, v \not\models \varphi$ for some $v \in W$ such that wRv ”. This is half of clause $[S\smile]$, from Section 2. Furthermore, the context relation π_1 induces an additional semantic condition: “if wRv then $\mathcal{M}, w \Vdash \Rightarrow \smile \varphi$ whenever $\mathcal{M}, v \Vdash \varphi \Rightarrow$ ”. This amounts to the other half of clause $[S\smile]$, namely: “ $\mathcal{M}, w \Vdash \smile \varphi$ whenever $\mathcal{M}, v \not\models \varphi$ for some $v \in W$ such that wRv ”. Systematically applying this semantic reading to all rules and all context relations of PK (according to Definitions 4.5 and 4.12 of [10]), one obtains soundness and completeness with respect to the class of all Kripke models $\langle \mathcal{F}, V \rangle$, where \mathcal{F} is an arbitrary frame and each valuation $V : W \times \mathcal{L} \rightarrow \{f, t\}$ respects the following conditions, for every $w \in W$ and $\varphi, \psi \in \mathcal{L}$:

$$\begin{array}{ll} [\mathbf{T}\top] & \mathbf{T}_w(\top) \\ [\mathbf{F}\perp] & \mathbf{F}_w(\perp) \end{array}$$

[T ∧]	if $\mathbf{T}_w(\varphi)$ and $\mathbf{T}_w(\psi)$, then $\mathbf{T}_w(\varphi \wedge \psi)$
[F ∧]	if $\mathbf{F}_w(\varphi)$ or $\mathbf{F}_w(\psi)$, then $\mathbf{F}_w(\varphi \wedge \psi)$
[T ∨]	if $\mathbf{T}_w(\varphi)$ or $\mathbf{T}_w(\psi)$, then $\mathbf{T}_w(\varphi \vee \psi)$
[F ∨]	if $\mathbf{F}_w(\varphi)$ and $\mathbf{F}_w(\psi)$, then $\mathbf{F}_w(\varphi \vee \psi)$
[T ¬]	if $\mathbf{F}_v(\varphi)$ for some $v \in W$ such that wRv , then $\mathbf{T}_w(\neg\varphi)$
[F ¬]	if $\mathbf{T}_v(\varphi)$ for every $v \in W$ such that wRv , then $\mathbf{F}_w(\neg\varphi)$
[T ¬]	if $\mathbf{F}_v(\varphi)$ for every $v \in W$ such that wRv , then $\mathbf{T}_w(\neg\varphi)$
[F ¬]	if $\mathbf{T}_v(\varphi)$ for some $v \in W$ such that wRv , then $\mathbf{F}_w(\neg\varphi)$
[T ⊙]	if $\mathbf{F}_w(\varphi)$ or $\mathbf{F}_w(\neg\varphi)$, then $\mathbf{T}_w(\odot\varphi)$
[F ⊙]	if $\mathbf{T}_w(\varphi)$ and $\mathbf{T}_w(\neg\varphi)$, then $\mathbf{F}_w(\odot\varphi)$
[T ⊙]	if $\mathbf{F}_w(\varphi)$ and $\mathbf{F}_w(\neg\varphi)$, then $\mathbf{T}_w(\odot\varphi)$
[F ⊙]	if $\mathbf{T}_w(\varphi)$ or $\mathbf{T}_w(\neg\varphi)$, then $\mathbf{F}_w(\odot\varphi)$

where we take ‘ $\mathbf{T}_u(\alpha)$ ’ as abbreviating ‘ $V(u, \alpha) = t$ ’, and ‘ $\mathbf{F}_u(\alpha)$ ’ as abbreviating ‘ $V(u, \alpha) = f$ ’. If alternatively one just *rewrites* $V(v, \alpha) = t$ as $\mathcal{M}, v \Vdash \alpha$ and rewrites $V(v, \alpha) = f$ as $\mathcal{M}, v \nVdash \alpha$, where $\mathcal{M} = \langle \langle W, R \rangle, V \rangle$, what results thereby is a collection of conditions that are essentially identical to the [S#] clauses introduced in our Section 2.

Two brief comments are in order here. First, our valuation functions assign truth-values to *every* formula in every world. However, as the values of compound formulas are uniquely determined by the values of their subformulas, we could have rested content above with assigning truth-values to propositional variables. Second, given that for the above valuations $\mathbf{T}_u(\alpha)$ is the case iff $\mathbf{F}_u(\alpha)$ fails to be the case, the semantic conditions [**T**#] and [**F**#], for each connective #, are clearly the converse of each other. In setting the two conditions apart, we have just given them directionality, pointing from less complex to more complex formulas, and have separated between conditions induced by rules from those induced by context relations. While neither of these manoeuvres are very useful here, they will allow us to more easily relate, in Section 4, valuations to ‘quasi valuations’ that have non truth-functional semantics.

Fix in what follows a Kripke model $\mathcal{M} = \langle \langle W, R \rangle, V \rangle$. We say that $w, v \in W$ *agree with respect to the formula α , according to V* , if either ($\mathbf{T}_w(\alpha)$ and $\mathbf{T}_v(\alpha)$) or ($\mathbf{F}_w(\alpha)$ and $\mathbf{F}_v(\alpha)$). We say that \mathcal{M} is *differentiated* if we have $w = v$ whenever w and v agree with respect to every $\alpha \in \mathcal{L}$, according to V . We call \mathcal{M} a *strengthened model* if wRv iff ($\mathbf{T}_v(\alpha)$ implies $\mathbf{F}_w(\neg\alpha)$) and ($\mathbf{F}_v(\alpha)$ implies $\mathbf{T}_w(\neg\alpha)$), for every $\alpha \in \mathcal{L}$. It is worth stressing that the accessibility relation of a strengthened model is uniquely determined by the underlying collection of worlds and valuation. The following result follows directly from Corollary 4.26 in [10], thus there is no need to prove it again here:

Theorem 3.1 *PK is sound and complete with respect to any class of Kripke models that: (i) contains only models that satisfy all the above [**T**#] and [**F**#] conditions; and (ii) contains all strengthened differentiated models that satisfy all the above [**T**#] and [**F**#] conditions.*

This theorem provides a mechanism that will be recycled in the subsequent

sections, when we consider extensions of PK. The following result from [5] comes as a byproduct of it:

Corollary 3.2 $\Gamma \models_{\mathcal{E}} \varphi$ iff $\Gamma \vdash_{\text{PK}} \varphi$ for every $\Gamma \cup \{\varphi\} \subseteq \mathcal{L}$, where \mathcal{E} denotes the class of all frames.

4 (Almost) Free Lunch: cut-elimination and analyticity

In this section we make further use of the powerful machinery introduced in [10] to prove that PK enjoys strong cut-admissibility, in other words, we show that $S \vdash_{\text{PK}} s$ implies that there is a derivation in PK of the sequent s from the set of sequents S such that in every application of the cut rule the cut formula φ appears in S . In particular, $\vdash_{\text{PK}} s$ implies that s is derivable in PK without any use of the cut rule. The proof is done in two steps. *First*, we present an adequate semantics for the cut-free fragment of PK. *Second*, we show that a countermodel in this new semantics entails the existence of a countermodel in the form of a Kripke model as defined in the previous section. This, together with Corollary 3.2, entails that PK is equivalent to its cut-free fragment.

Step 1. Semantics for cut-free PK

Semantics for cut-free basic systems may be obtained through the use of ‘quasi valuations’. Models based on quasi valuations differ from usual Kripke models in two main aspects: (a) the underlying interpretation is three-valued; (b) the underlying interpretation is non-deterministic — the truth-value of a compound formula in a given world is *not* always uniquely determined by the truth values of its subformulas in the collection of worlds of the underlying frame.

To obtain such semantics for PK, as before, one reads off a semantic condition on quasi valuations from each derivation rule and from each context relation. As per Theorems 5.24 and 5.31 of [10], we know that the class of models based on quasi valuations that respect all these conditions is sound and complete for the cut-free fragment of PK. Concretely, given a frame $\mathcal{F} = \langle W, R \rangle$, a *quasi valuation* over it is a function $QV : W \times \mathcal{L} \rightarrow \{\{f\}, \{t\}, \{f, t\}\}$ satisfying precisely the same semantic conditions laid down in Section 3, where we now take ‘ $\mathbf{T}_u(\alpha)$ ’ as abbreviating ‘ $t \in V(u, \alpha)$ ’, and ‘ $\mathbf{F}_u(\alpha)$ ’ as abbreviating ‘ $f \in V(u, \alpha)$ ’. Whenever we need to distinguish between a semantic condition on a tuple $\langle w, \varphi \rangle$ as constraining a valuation V or a quasi valuation QV , we will use $\mathbf{X}_w(\varphi)$ for the former and $\mathbf{X}_w^Q(\varphi)$ for the latter, where $\mathbf{X} \in \{\mathbf{T}, \mathbf{F}\}$. A *quasi model* is a structure $\mathcal{QM} = \langle \mathcal{F}, QV \rangle$, where QV is a quasi valuation over \mathcal{F} . The notions of a differentiated quasi model and of a strengthened quasi model are defined as before, assuming the same abbreviations.

Step 2. Semantic cut-admissibility

The next step is to show that the existence of a countermodel in the form of a strengthened differentiated quasi model implies the existence a countermodel in the form of an ordinary Kripke model (following Corollary 5.48 of [10]). For this purpose we define an *instance* of a quasi model $\mathcal{QM} = \langle \langle W, R \rangle, QV \rangle$ as any model of the form $\mathcal{M} = \langle \langle W, R' \rangle, V \rangle$ such that $\mathbf{X}_w^Q(\varphi)$ whenever $\mathbf{X}_w(\varphi)$,

for every $\mathbf{X} \in \{\mathbf{T}, \mathbf{F}\}$, every $w \in W$ and every $\varphi \in \mathcal{L}$. Note that a quasi model and its instances may have different accessibility relations.

In what follows, the construction of appropriate instances is done by a recursive definition over the following well-founded relation \prec on the set of formulas: $\alpha \prec \beta$ if either (i) α is a proper subformula of β ; (ii) $\alpha = \neg\gamma$ and $\beta = \odot\gamma$ for some $\gamma \in \mathcal{L}$; or (iii) $\alpha = \neg\gamma$ and $\beta = \ominus\gamma$ for some $\gamma \in \mathcal{L}$.

Lemma 4.1 *Every quasi model has an instance.*

Proof. Let $\mathcal{QM} = \langle \mathcal{F}, QV \rangle$ be a quasi model based on a frame $\mathcal{F} = \langle W, R \rangle$. We set us now an appropriate valuation $V : W \times \mathcal{L} \rightarrow \{f, t\}$. For every world w and formula φ , the valuation V is inductively defined (with respect to \prec) on φ as follows: (R1) if $\mathbf{T}_w^Q(\varphi)$ fails for QV , we postulate $\mathbf{F}_w(\varphi)$ to be the case for V ; (R2) if $\mathbf{F}_w^Q(\varphi)$ fails for QV , we postulate $\mathbf{T}_w(\varphi)$ to be the case for V ; (R3) otherwise both $\mathbf{T}_w^Q(\varphi)$ and $\mathbf{F}_w^Q(\varphi)$ hold good for QV , and in this case we postulate $\mathbf{T}_w(\varphi)$ to be the case for V if one of the following holds:

- (M1) φ is a propositional variable or φ is \top
- (M2) $\varphi = \varphi_1 \wedge \varphi_2$, and both $\mathbf{T}_w(\varphi_1)$ and $\mathbf{T}_w(\varphi_2)$
- (M3) $\varphi = \varphi_1 \vee \varphi_2$, and either $\mathbf{T}_w(\varphi_1)$ or $\mathbf{T}_w(\varphi_2)$
- (M4) $\varphi = \neg\psi$, and $\mathbf{F}_v(\psi)$ for some $v \in W$ such that wRv
- (M5) $\varphi = \neg\psi$, and $\mathbf{F}_v(\psi)$ for every $v \in W$ such that wRv
- (M6) $\varphi = \odot\psi$, and either $\mathbf{F}_w(\psi)$ or $\mathbf{F}_w(\neg\psi)$
- (M7) $\varphi = \odot\psi$, and both $\mathbf{F}_w(\psi)$ and $\mathbf{F}_w(\neg\psi)$

Otherwise, we postulate $\mathbf{F}_w(\varphi)$ to be the case for V . Obviously, $\mathbf{X}_w(\varphi)$ implies $\mathbf{X}_w^Q(\varphi)$ for every $w \in W$, every $\varphi \in \mathcal{L}$ and every $\mathbf{X} \in \{\mathbf{T}, \mathbf{F}\}$. It is routine to verify that $\langle \mathcal{F}, V \rangle$ is a model. We show here that the semantic conditions for \neg and \odot hold:

[**Case of \neg**] Let $\psi \in \mathcal{L}$. Suppose first that $\mathbf{F}_v(\psi)$ is the case for some $v \in W$ such that wRv . Then $\mathbf{F}_v^Q(\psi)$. Since \mathcal{QM} is a quasi model, then $\mathbf{T}_w^Q(\neg\psi)$ is the case. If, on the one hand, $\mathbf{F}_w^Q(\neg\psi)$ fails, then we must have $\mathbf{T}_w(\neg\psi)$, by (R2). If, on the other hand, neither $\mathbf{T}_w^Q(\neg\psi)$ nor $\mathbf{F}_w^Q(\neg\psi)$ fail, we are in case (R3). Since we have $\mathbf{F}_v(\psi)$ and wRv we conclude by (M4) that $\mathbf{T}_w(\neg\psi)$ must be the case. Suppose now that $\mathbf{T}_v(\psi)$ is the case for every world v such that wRv . Then we have $\mathbf{T}_v^Q(\psi)$ for every such world. Since \mathcal{QM} is a quasi model, it follows that $\mathbf{F}_w^Q(\neg\psi)$ is the case. If, on the one hand, $\mathbf{T}_w^Q(\neg\psi)$ fails, then we must have $\mathbf{F}_w(\neg\psi)$, by (R1). If, on the other hand, neither $\mathbf{T}_w^Q(\neg\psi)$ nor $\mathbf{F}_w^Q(\neg\psi)$ fail, we are in case (R3). Since we have $\mathbf{T}_v(\psi)$ for every world v such that wRv we conclude that none of (M1)–(M7) applies, thus $\mathbf{F}_w(\neg\psi)$ must be the case.

[**Case of \odot**] Let $\psi \in \mathcal{L}$. Suppose first that either $\mathbf{F}_w(\psi)$ or $\mathbf{F}_w(\neg\psi)$ are the case for some $w \in W$. Then either $\mathbf{F}_w^Q(\psi)$ or $\mathbf{F}_w^Q(\neg\psi)$. Since \mathcal{QM} is a quasi model, it follows that $\mathbf{T}_w^Q(\odot\psi)$. If, on the one hand, $\mathbf{F}_w^Q(\odot\psi)$ fails, then we must have $\mathbf{T}_w(\odot\psi)$, by (R2). If, on the other hand, neither $\mathbf{T}_w^Q(\odot\psi)$ nor $\mathbf{F}_w^Q(\odot\psi)$ fail, we are in case (R3) and we conclude by (M6) that $\mathbf{T}_w(\odot\psi)$ must be the case. Suppose now that both $\mathbf{T}_w(\psi)$ and $\mathbf{T}_w(\neg\psi)$ are the case for some $w \in W$. Then $\mathbf{T}_w^Q(\psi)$ and $\mathbf{T}_w^Q(\neg\psi)$. Since \mathcal{QM} is a quasi model, then $\mathbf{F}_w^Q(\odot\psi)$. If, on the

one hand, $\mathbf{T}_w^Q(\ominus\psi)$ fails, then we must have $\mathbf{F}_w(\ominus\psi)$, by (R1). If, on the other hand, neither $\mathbf{T}_w^Q(\ominus\psi)$ nor $\mathbf{F}_w^Q(\ominus\psi)$ fail, we are in case (R3) and $\mathbf{F}_w(\ominus\psi)$ must be the case because none of (M1)–(M7) applies. \square

Since the class of all quasi models contains the strengthened differentiated quasi models, it follows that:

Corollary 4.2 *PK enjoys strong cut-admissibility.*

Corollary 4.3 *PK is \prec -analytic: If a sequent s is derivable from a set S of sequents in PK, then there is a derivation of s from S such that every formula φ that occurs in the derivation satisfies $\varphi \prec \psi$ for some ψ in $S \cup s$.*

Proof. By induction on the length of the derivation of s from S in PK: In all rules except for (*cut*), the premises include only formulas φ that satisfy $\varphi \prec \psi$ for some formula ψ in the conclusion. \square

5 Some special classes of frames

In this section we present three very natural deductive extensions of PK. Given a property X of binary relations, we call a frame $\langle W, R \rangle$ an X frame if R enjoys X . A (quasi) model $\langle \mathcal{F}, V \rangle$ is called an X (quasi) model if \mathcal{F} is an X frame. In addition, and similarly to what we did in the case of PK, for every proof system Y we write $S \vdash_Y s$ if there is a derivation of s from S in Y .

5.1 Seriality

Let PKD be the system obtained by augmenting PK with the following rule:

$$[\mathbf{D}] \frac{\Gamma \Rightarrow \Delta}{\neg\Delta \Rightarrow \neg\Gamma}$$

This rule may be formulated as the basic rule: $\langle \Rightarrow ; \pi_1 \rangle / \Rightarrow$. Since its premise is the empty sequent, the semantic condition it imposes (following [10]) is seriality: indeed, respecting $[\mathbf{D}]$ in a world w of a model \mathcal{M} based on a frame $\langle W, R \rangle$ means that if $\mathcal{M}, v \Vdash \Rightarrow$ for every world v such that wRv , then also $\mathcal{M}, w \Vdash \Rightarrow$. Since the empty sequent is not satisfied at any world, this condition would hold iff for every world w there exists a world v such that wRv . A similar argument shows that every serial frame satisfies this semantic condition.

As in Corollary 3.2, we obtain a completeness theorem for PKD with respect to serial models:

Corollary 5.1 $\Gamma \models_{\mathcal{E}_D} \varphi$ iff $\Gamma \vdash_{\text{PKD}} \varphi$ for every $\Gamma \cup \{\varphi\} \subseteq \mathcal{L}$, where \mathcal{E}_D is the class of serial models.

Additionally, we may prove cut-admissibility also for PKD, going through serial quasi models.

Lemma 5.2 *Every serial quasi model has a serial instance.*

Proof. The proof is the same as the proof of Lemma 4.1. Note indeed that no property of the accessibility relation was assumed, and the constructed instance has the same accessibility relation as the original quasi model. \square

Corollary 5.3 *PKD enjoys cut-admissibility and is \prec -analytic.*

5.2 Reflexivity

Let PKT be the system obtained by augmenting PK with the following rules:

$$[\Rightarrow\smile] \frac{\Gamma, \varphi \Rightarrow \Delta}{\Gamma \Rightarrow \smile\varphi, \Delta} \qquad [\wedge\Rightarrow] \frac{\Gamma \Rightarrow \varphi, \Delta}{\Gamma, \wedge\varphi \Rightarrow \Delta}$$

These rules may be formulated as the basic rules: $\langle p_1 \Rightarrow ; \pi_0 \rangle / \Rightarrow \smile p_1$ and $\langle \Rightarrow p_1 ; \pi_0 \rangle / \wedge p_1 \Rightarrow$. It should be clear that PKT allows thus for the derivation of the consecutions representing $\llbracket \smile\text{-implosion} \rrbracket$ and $\llbracket \wedge\text{-explosion} \rrbracket$.

Semantically, they impose reflexivity not on all models, but only on *strengthened* models. Indeed, since the underlying context relation is π_0 , for every model $\mathcal{M} = \langle \mathcal{F}, V \rangle$ based on a frame $\mathcal{F} = \langle W, R \rangle$ that respects $[\Rightarrow\smile]$ and $[\wedge\Rightarrow]$, and every world w , if $\mathcal{M}, w \vDash \varphi \Rightarrow$ then $\mathcal{M}, w \vDash \smile\varphi$ and if $\mathcal{M}, w \vDash \wedge\varphi$ then $\mathcal{M}, w \vDash \varphi \Rightarrow$. To put it otherwise, if $\mathbf{F}_w(\varphi)$ then $\mathbf{T}_w(\smile\varphi)$, and if $\mathbf{T}_w(\varphi)$ then $\mathbf{F}_w(\wedge\varphi)$. Clearly, every reflexive model satisfies these conditions. To show that every strengthened model that satisfies them is reflexive, consider an arbitrary strengthened model $\mathcal{M} = \langle \langle W, R \rangle, V \rangle$. Then for every world $w \in W$ we have that for every formula φ , ($\mathbf{T}_w(\varphi)$ implies $\mathbf{F}_w(\smile\varphi)$) and ($\mathbf{F}_w(\varphi)$ implies $\mathbf{T}_w(\wedge\varphi)$), which in strengthened models means precisely that wRw . We obtain thus a completeness theorem for PKT with respect to reflexive models:

Corollary 5.4 $\Gamma \vDash_{\mathcal{E}_T} \varphi$ iff $\Gamma \vdash_{\text{PKT}} \varphi$ for every $\Gamma \cup \{\varphi\} \subseteq \mathcal{L}$, where \mathcal{E}_T is the class of reflexive models.

Such semantics for PKT allows one to easily confirm that the full type diamond-minus connective \smile fails (DM1.2#), and that the full type box-minus connective \wedge fails (DM2.2#). These properties transfer to the weaker logics PKD and PK, of course.

Cut-admissibility for PKT may be obtained using arguments similar to those used in proving Lemma 4.1. It follows thus that:

Lemma 5.5 Every reflexive strengthened quasi model has a reflexive instance.

Corollary 5.6 PKT enjoys cut-admissibility and is \prec -analytic.

5.3 Functionality

In this section we address functional frames, that is, frames whose accessibility relations are *total functions*. In every model $\langle \langle W, R \rangle, V \rangle$ of a functional frame and world $w \in W$, we have $\mathbf{T}_w(\smile\varphi)$ iff $\mathbf{T}_w(\wedge\varphi)$. Hence \smile and \wedge are indistinguishable. Accordingly, here we consider a restricted language, without \wedge .

Let PKF be the system obtained from PK by substituting \smile for \wedge in rules $[\Rightarrow\smile]$ and $[\smile\Rightarrow]$, and replacing both rules $[\Rightarrow\wedge]$ and $[\wedge\Rightarrow]$ with the single rule:

$$[\mathbf{Fun}] \frac{\Gamma \Rightarrow \Delta}{\smile\Delta \Rightarrow \smile\Gamma}$$

It is straightforward to see that rule $[\mathbf{Fun}]$ may be formulated as the following basic rule: $\langle \Rightarrow ; \pi_2 \rangle / \Rightarrow$, for $\pi_2 = \{ \langle \varphi \Rightarrow ; \Rightarrow \smile\varphi \rangle, \langle \Rightarrow \varphi ; \smile\varphi \Rightarrow \rangle \}$.

The latter rule and context relation impose functionality on *differentiated* models. Indeed, respecting the basic rule $[\mathbf{Fun}]$ corresponds to seriality, similarly to the case of the rule $[\mathbf{D}]$. Additionally, the context relation π_2 forces

the accessibility relation to be a partial function: respecting π_2 in a world w of a model $\mathcal{M} = \langle \langle W, R \rangle, V \rangle$ means that for every $v_1, v_2 \in W$ such that wRv_1 and wRv_2 and for every formula φ we have that $\mathbf{T}_{v_1}(\varphi)$ iff $\mathbf{F}_w(\sim\varphi)$ iff $\mathbf{T}_{v_2}(\varphi)$. When \mathcal{M} is differentiated, this implies that $v_1 = v_2$. Now, every functional model satisfies these semantic conditions and every differentiated model that satisfies them is functional. We thus obtain a completeness result for PKF with respect to functional models:

Corollary 5.7 $\Gamma \models_{\mathcal{E}_{\mathbf{Fun}}} \varphi$ iff $\Gamma \vdash_{\text{PKF}} \varphi$ for every $\Gamma \cup \{\varphi\} \subseteq \mathcal{L}$, where $\mathcal{E}_{\mathbf{Fun}}$ is the class of functional models.

In contrast with what was the case for PKT, within such semantics for PKF there are no longer countermodels for (DM1.2 \sim) or for (DM2.2 \wedge). At any rate, it should be clear that PKF extends PKD, but does not extend PKT.

Going through quasi models we may prove cut-admissibility also for PKF. However, unlike in previous cases, considering functional quasi models will not suffice. Indeed, there exist differentiated strengthened quasi models that respect **[Fun]** whose accessibility relation is not a total function. Let a **Fun** quasi model $\mathcal{QM} = \langle \mathcal{F}, QV \rangle$ based on a frame $\mathcal{F} = \langle W, R \rangle$ be a serial quasi model in which for every $w, v \in W$ such that wRv we have, for every $\varphi \in \mathcal{L}$, both $(\mathbf{F}_v^Q(\varphi)$ implies $\mathbf{T}_w^Q(\sim\varphi)$) and $(\mathbf{T}_v^Q(\varphi)$ implies $\mathbf{F}_w^Q(\sim\varphi)$). We note that, although the accessibility relation in **Fun** quasi models may not be a total function, we are still able to extract a functional model from it:

Lemma 5.8 *Every Fun quasi model has a functional instance.*

Proof. Let $\mathcal{QM} = \langle \mathcal{F}, QV \rangle$ be an **Fun** quasi model based on a frame $\langle W, R \rangle$. Since \mathcal{QM} is an **Fun** quasi model, we have in particular that R is serial. Therefore, there exists some total function $R' : W \rightarrow W$ such that $R' \subseteq R$. Let $\mathcal{F}' = \langle W, R' \rangle$. We define an appropriate valuation $V : W \times \mathcal{L} \rightarrow \{f, t\}$ as in Lemma 4.1, while disregarding (M5), and using the following instead of (M4) and (M7):

- (M4') $\varphi = \sim\psi$, and $\mathbf{F}_{R'(w)}(\psi)$
(M7') $\varphi = \ominus\psi$, and $\mathbf{F}_w(\varphi)$ and $\mathbf{F}_w(\sim\varphi)$

The proof then carries on in a similar fashion to the proof of Lemma 4.1. \square

Corollary 5.9 *PKF enjoys cut-admissibility and is \prec' -analytic, where \prec' is the restriction of \prec to the \wedge -free fragment of \mathcal{L} , with an additional clause according to which $\sim\varphi \prec \ominus\varphi$.*

We include a word about further developments which could not be included here for reasons of space. It is easy to see that \sim and \wedge may be defined using the customary presentation of the modal logic **K** by $\sim\varphi := \sim\Box\varphi$ and $\wedge\varphi := \Box\sim\varphi$. When considering only functional frames (like in PKF), we get a translation to KF — the ordinary modal logic of functional Kripke models. For the $\ominus\ominus$ -free fragment of this logic, we may apply the general reduction to SAT proposed in [11], which in particular means that the derivability problem for it is in co-NP. We further note that if one dismisses $[\vee\Rightarrow]$ from the proof system, derivability

can be decided in linear time, by producing SAT-instances that consist solely of Horn clauses. Such ‘half-disjunction’ was also suggested in the context of primal infon logic [2], to obtain a linear time decision procedure.

5.4 Symmetry

Let PKB be the system obtained from PK by replacing $[\sim\Rightarrow]$ and $[\Rightarrow\sim]$ with the following rules:

$$[\mathbf{B}_1] \frac{\Gamma, \sim\Gamma', \varphi \Rightarrow \Delta, \wedge\Delta'}{\wedge\Delta, \Delta' \Rightarrow \wedge\varphi, \sim\Gamma, \Gamma'} \quad [\mathbf{B}_2] \frac{\Gamma, \sim\Gamma' \Rightarrow \varphi, \Delta, \wedge\Delta'}{\wedge\Delta, \Delta', \sim\varphi \Rightarrow \sim\Gamma, \Gamma'}$$

These correspond to the following basic rules: $\langle p_1 \Rightarrow ; \pi_3 \rangle / \Rightarrow \wedge p_1$ and $\langle \Rightarrow p_1 ; \pi_3 \rangle / \wedge p_1 \Rightarrow$, for the context relation $\pi_3 = \{ \langle p_1 \Rightarrow ; \Rightarrow \wedge p_1 \rangle, \langle \wedge p_1 \Rightarrow ; \Rightarrow p_1 \rangle, \langle \Rightarrow p_1 ; \wedge p_1 \Rightarrow \rangle, \langle \Rightarrow \wedge p_1 ; p_1 \Rightarrow \rangle \}$. This relation satisfies the following property: $s \pi_3 q$ iff $\bar{q} \pi_3 \bar{s}$, where $(\Rightarrow \bar{\varphi})$ denotes $(\varphi \Rightarrow)$ and $(\bar{\varphi} \Rightarrow)$ denotes $(\Rightarrow \varphi)$. By Proposition 4.28 of [10], the semantic condition these rules impose on strengthened models is symmetry of the accessibility relation. In addition, every symmetric model respects these rules, as well as the context relation π_3 . It follows that:

Corollary 5.10 $\Gamma \models_{\mathcal{E}_B} \varphi$ iff $\Gamma \vdash_{\text{PKB}} \varphi$ for every $\Gamma \cup \{\varphi\} \subseteq \mathcal{L}$, where \mathcal{E}_B is the class of symmetric models.

Symmetric frames are relevant from the viewpoint of sub-classical properties of negation. They validate, for instance, the consecutions $\sim\sim p \models p$ and $p \models \wedge\sim p$. Paraconsistent logics based on symmetric (and reflexive) frames are also studied in [1], a paper that investigates in detail a conservative extension of the corresponding logic, obtained by the addition of a classical implication (but without primitive \wedge and \odot), and offers for this logic a sequent system for which cut is not eliminable.

Quasi models for PKB are not necessarily symmetric, making it harder to convert them into instances in the form of symmetric models. This is why cut-admissibility for our system PKB is here left open as a matter for further research. However, using a similar technique of basic systems, it can be straightforwardly shown that PKB is \prec -analytic. This does not require quasi models at all: one only has to show that every *partial* model, whose valuation’s domain is closed under \prec -subformulas, may be extended to a full model (see Corollary 5.44 in [10]).

6 Definability of classical negation

In this section we investigate definability of classical negation in the modal logics studied in this paper. Given a set C of connectives and a logic \mathbf{L} , we denote by $\mathbf{L}\text{-}C$ the C -free fragment of \mathbf{L} , that is, the restriction of \mathbf{L} to the language without the connectives in C .

Theorem 6.1

- (i) *Classical negation is definable in the logics: $\text{PKT}\text{-}\{\sim, \odot\}$, $\text{PKT}\text{-}\{\wedge, \odot\}$, PKD , and PKF .*

(ii) *Classical negation is not definable in the logics: PK, PKB, PKT- $\{\ominus, \odot\}$, PKD- $\{\ominus\}$, PKD- $\{\odot\}$, PKF- $\{\ominus\}$, and PKF- $\{\odot\}$.*

Proof.

(i) For $PKT\text{-}\{\smile, \ominus\}$ we set $\sim\varphi := \smile\varphi \vee \odot\varphi$, for $PKT\text{-}\{\wedge, \odot\}$ we set $\sim\varphi := \smile\varphi \wedge \odot\varphi$, and for PKD and PKF we set $\sim\varphi := (\smile\varphi \wedge \odot\varphi) \vee \odot\varphi$. It is easy to see that $\Rightarrow\varphi, \sim\varphi$ and $\varphi, \sim\varphi \Rightarrow$ are derivable in each system for the defined connective \sim . Using cut, one obtains the usual sequent rules for classical negation. Table 1 provides the derivations for PKD . (Given that PKF is a deductive extension of PKD , the derivation in Table 1 is also good for PKF .)

$\frac{\frac{\frac{\varphi \Rightarrow \varphi}{\varphi \Rightarrow \varphi} [\mathbf{D}] \quad \frac{\varphi \Rightarrow \varphi}{\sim\varphi \Rightarrow \smile\varphi} [\mathbf{D}]}{\varphi, \sim\varphi, \odot\varphi \Rightarrow} [\odot\Rightarrow] \quad \frac{\varphi \Rightarrow \varphi, \smile\varphi}{\varphi, \odot\varphi \Rightarrow} [\odot\Rightarrow]}{\frac{\varphi, \sim\varphi, \odot\varphi \Rightarrow}{\varphi, \smile\varphi \wedge \odot\varphi \Rightarrow} [\wedge\Rightarrow] \quad \frac{\varphi \Rightarrow \varphi, \smile\varphi}{\varphi, \odot\varphi \Rightarrow} [\odot\Rightarrow]}{\varphi, (\smile\varphi \wedge \odot\varphi) \vee \odot\varphi \Rightarrow} [\vee\Rightarrow]$
$\frac{\frac{\varphi \Rightarrow \varphi}{\Rightarrow \varphi, \smile\varphi, \odot\varphi} [\Rightarrow\odot] \quad \frac{\varphi, \smile\varphi \Rightarrow \varphi, \odot\varphi}{\Rightarrow \varphi, \odot\varphi, \odot\varphi} [\Rightarrow\odot]}{\Rightarrow \varphi, (\smile\varphi \wedge \odot\varphi), \odot\varphi} [\Rightarrow\wedge]}{\Rightarrow \varphi, (\smile\varphi \wedge \odot\varphi) \vee \odot\varphi} [\Rightarrow\vee]$

Table 1

(ii) Let $X \in \{PK, PKB, PKT\text{-}\{\ominus, \odot\}, PKD\text{-}\{\ominus\}, PKD\text{-}\{\odot\}, PKF\text{-}\{\ominus\}, PKF\text{-}\{\odot\}\}$. Suppose for the sake of contradiction that classical negation \sim is definable in X . Let $p \in \mathcal{P}$ and let φ be $\sim(p)$. Then both $\Rightarrow\varphi, p$ and $p, \varphi \Rightarrow$ are valid in X . Consider a set W that consists of two worlds, w and v , and a valuation V such that $V(w, q) = 1$ and $V(v, q) = 0$ for every atomic formula q (including p). Now, for each relation R_X on W , consider the model $\mathcal{M}_X = \langle\langle W, R_X \rangle, V\rangle$. If \mathcal{M}_X belongs to the class of models that semantically characterize X , then we must have that $\mathcal{M}_X, w \Vdash \varphi, p \Rightarrow$ and $\mathcal{M}_X, v \Vdash \Rightarrow p, \varphi$. Since in \mathcal{M}_X we have $\mathbf{T}_w(p)$ and $\mathbf{F}_v(p)$, we must then have $\mathbf{F}_w(\varphi)$ and $\mathbf{T}_v(\varphi)$. We show that this is impossible, by structural induction on φ . More precisely, we claim that if $\mathbf{F}_w(\varphi)$ then $\mathbf{F}_v(\varphi)$. To show this, we consider the possible values for X , and define the accessibility relation R_X in each case. For $X \in \{PK, PKB\}$ define $R_X = \emptyset$, for $X = PKT\text{-}\{\ominus, \odot\}$ define $R_X = W \times W$, for $X \in \{PKD\text{-}\{\ominus\}, PKF\text{-}\{\ominus\}\}$ define $R_X = \{\langle w, v \rangle, \langle v, v \rangle\}$, and for $X \in \{PKD\text{-}\{\odot\}, PKF\text{-}\{\odot\}\}$ define $R_X = \{\langle w, w \rangle, \langle v, w \rangle\}$. We describe in detail only the third case. For this case, note that since R_X is a total function, \mathcal{M}_X belongs to the appropriate class of models, and \smile and \wedge are indistinguishable, hence we may choose to consider \wedge instead of \smile . The cases where φ is atomic, a conjunction, or a disjunction are trivial. If $\varphi = \wedge\psi$ for some ψ and $\mathbf{F}_w(\varphi)$, then we must have $\mathbf{T}_v(\psi)$ by $[\mathbf{T}\wedge]$, which implies by $[\mathbf{F}\wedge]$ that $\mathbf{F}_v(\varphi)$. If $\varphi = \odot\psi$ for some ψ , then $\mathbf{F}_v(\varphi)$ must hold good: indeed, if on

the one hand $\mathbf{T}_w(\neg\psi)$ then $\mathbf{F}_v(\psi)$ by $[\mathbf{F}\neg]$, and hence $\mathbf{T}_v(\neg\psi)$ by $[\mathbf{T}\neg]$, which implies by $[\mathbf{F}\odot]$ that $\mathbf{F}_v(\varphi)$; if on the other hand $\mathbf{F}_w(\neg\psi)$ then $\mathbf{T}_v(\psi)$ by $[\mathbf{T}\neg]$, and hence again $\mathbf{F}_v(\varphi)$ follows by $[\mathbf{F}\odot]$. \square

7 This is possibly not the end

In contrast to the usual ‘positive modalities’ of normal modal logics, which are monotone with respect to the underlying notion of consequence, we have devoted this paper to antitone connectives known as ‘negative modalities’ — specifically, to full type box-minus and full-type diamond-minus connectives.

Be they monotone or antitone on each of their arguments, the connectives of normal modal logics are always congruential: they treat equivalent formulas as synonymous. The phenomenon seems to be an exception rather than the rule if many-valued logics with non-classical negations are involved. For instance, Kleene’s 3-valued logic fails to be congruential, as $p \wedge \neg p$ is equivalent to $q \wedge \neg q$, but their respective negations, $\neg(p \wedge \neg p)$ and $\neg(q \wedge \neg q)$, are not equivalent. Also, the earliest paraconsistent logic in the literature (cf. [8]) fails to be congruential, in spite of having been defined in terms of a translation into a fragment of the modal logic $S5$, and this failure remained unknown for decades (cf. [12]). The same holds for the other early paraconsistent logics developed later on, containing extra ‘strong negations’ that live in the vicinity of classical negation (cf. [15,4]). Of course, there are important ‘non-exceptions’: intuitionistic logic and other intermediate logics constitute congruential paracomplete logics. For another example perhaps more to the point, consider the four-valued logic of FDE, whose semantics may be formulated having as truth-values $\{\mathbf{t}, \mathbf{b}, \mathbf{n}, \mathbf{f}\}$, where $\{\mathbf{t}, \mathbf{b}\}$ are designated, the transitive reflexive closure of the order \leq such that $\mathbf{f} \leq \mathbf{b}, \mathbf{n} \leq \mathbf{t}$ may be used to define \wedge and \vee , respectively, as its meet and its join, while $\neg\langle \mathbf{t}, \mathbf{b}, \mathbf{n}, \mathbf{f} \rangle := \langle \mathbf{f}, \mathbf{b}, \mathbf{n}, \mathbf{t} \rangle$. It is not hard to see that this logic is congruential and by defining the operators $\ominus\langle \mathbf{t}, \mathbf{b}, \mathbf{n}, \mathbf{f} \rangle := \langle \mathbf{t}, \mathbf{n}, \mathbf{b}, \mathbf{t} \rangle$ and $\odot\langle \mathbf{t}, \mathbf{b}, \mathbf{n}, \mathbf{f} \rangle := \langle \mathbf{f}, \mathbf{n}, \mathbf{b}, \mathbf{f} \rangle$ it gets conservatively extended into another congruential logic that deductively extends our logic PKF (but does not deductively extend PKT), if we interpret \sim as \neg . It is worth noting that the latter logic is equivalent to the expansion of FDE by the addition of a classical negation.

Some terminological conventions and some concepts used in the present paper were borrowed or adapted from other fonts, sometimes without explicit reference. For instance, in Section 2, dadaistic and nihilistic models come from [13], and that paper also introduces the connectives \ominus and \odot of the so-called Logics of Formal Inconsistency (cf. [3]) and the dual Logics of Formal Undeterminedness (cf. [13], where the adjustment connectives are called connectives ‘of perfection’). The minimal conditions on negation, called $[\textit{falsificatio}]$ and $[\textit{verificatio}]$, come from [14]. What we in the present paper call ‘determinacy’ has in [5] been called ‘determinedness’. The ‘strengthened models’ from Section 3 correspond to models with strongly-legal valuations in the terminology of [10]. In Section 5, Rule $[\mathbf{D}]$ may be thought of as a variation on the following well-known sequent rule for the modal logic KD : $\Gamma \Rightarrow / \Box\Gamma \Rightarrow$, and rules for PKT are variations on the usual sequent rule for the modal logic

KT: $\Gamma, \varphi \Rightarrow \Delta / \Gamma, \Box \varphi \Rightarrow \Delta$ (cf. [22]). Also, the rule for PKF is a variation on the sequent rule from [9] for the ‘Next’ operator in the temporal logic *LTL*, namely: $\Gamma \Rightarrow \Delta / \Box \Gamma \Rightarrow \Box \Delta$. We have not been able to find in the literature the obvious rules $\Gamma, \varphi \Rightarrow \Delta / \Box \Gamma, \Diamond \varphi \Rightarrow \Diamond \Delta$ and $\Gamma \Rightarrow \varphi, \Delta / \Box \Gamma \Rightarrow \Box \varphi, \Diamond \Delta$ for the modal logic *K* of which our rules $[\smile \Rightarrow]$ and $[\Rightarrow \smile]$ from Section 3 would be variations on. In Section 4, the trick behind using three-valued models for addressing the admissibility of the cut rule goes at least as far back as [20].

The main feature of our approach here has been to rely on theoretical technology built elsewhere and show how it may be adapted to the present case. Our hope is that this should prove a beneficial methodology, and that the idea of obtaining completeness and cut admissibility as particular applications of more general results will become more common, rather than proceeding always through *ad hoc* completeness and cut elimination theorems.

While we have directed our attention, in this paper, to classes of frames that turned out to be particularly significant from the viewpoint of the relation between negative modalities of different types, we see two very natural ways of extending such study. The *first* natural extension would be to look at other classes of frames that prove to be relevant from the viewpoint of sub-classical properties of negation. For instance, it is easy to see that the class of frames with the Church-Rosser property validates $\smile \smile p \models \smile \smile p$, pinpointing an interesting consecution involving the interaction between negations of different types. Some other classes of frames deserving study do not seem to show the same amount of promise, from the viewpoint of paraconsistency or para-completeness. For instance, euclidean frames validate $\llbracket \smile\text{-explosion} \rrbracket$ if in the set of formulas $\{\smile p, p\}$ one replaces p by $\smile r$, and validate $\llbracket \smile\text{-implosion} \rrbracket$ if in $\{\smile p, p\}$ one replaces p by $\smile r$; also, transitive frames cause a similar behavior, but now swapping the roles of $\smile r$ and $\smile r$ in replacing p . Alternatively, a *second* avenue worth exploring would lead us into logics containing more than one negative modality of the same type (as it has been done for logics with multiple para-complete negations in [18]). One could for instance consider not only the ‘forward-looking’ negative modalities defined by the semantic clauses $[S\smile]$ and $[S\smile]$, but also ‘backward-looking’ negative modalities \smile^{-1} and \smile^{-1} defined by the clauses obtained from the latter ones by replacing wRv by vRw (such ‘converse modalities’ have been studied in the context of temporal logic [16], as well as in the context of the so-called Heyting-Brouwer logic [17]). The interaction between the various negations would then be witnessed, in such extended language, by the validity over arbitrary frames of ‘pure’ consecutions such as $\smile^{-1}\smile p \models p$ and $\smile\smile^{-1}p \models p$ (as well as $p \models \smile^{-1}\smile p$ and $p \models \smile\smile^{-1}p$), and the validity over symmetric frames of ‘mixed’ consecutions such as $\smile^{-1}\smile p \models p$ and $\smile\smile^{-1}p \models p$ (as well as $p \models \smile^{-1}\smile p$ and $p \models \smile\smile^{-1}p$). In our view, it seems worth the effort applying the machinery employed in the present paper to the above mentioned systems, and still others, in order to investigate results analogous to the ones we have here looked at.¹

¹ The authors acknowledge partial support by the Marie Curie project GeTFun (PIRSES-

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