

Symmetry-based Task Reduction for Relaxed Reachability Analysis

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Reachability

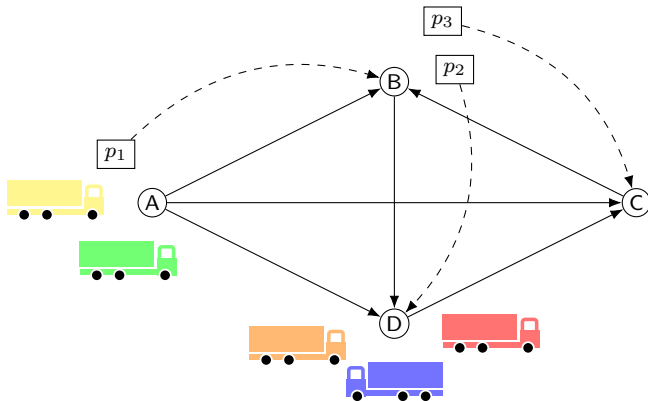
Question: Which atoms **can become true**
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Reachability

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in the **reachable** part of the state space?

- Relevant for grounding, mutexes (pairs of atoms), ...
- As **hard** as the planning problem
- Usually: relaxation-based over-approximation

Example Task



Idea

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Perform analysis for fewer trucks and packages.

Illustration on Example Task

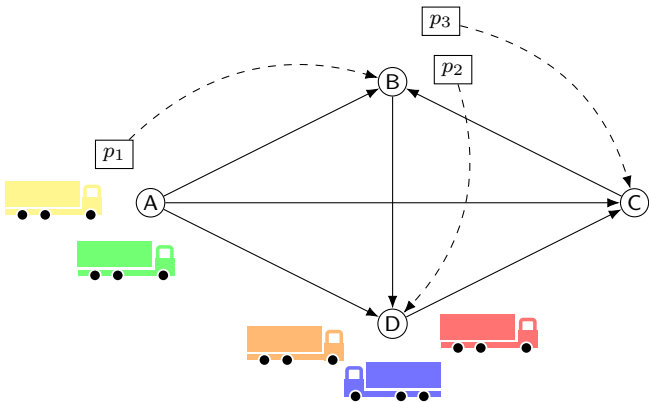


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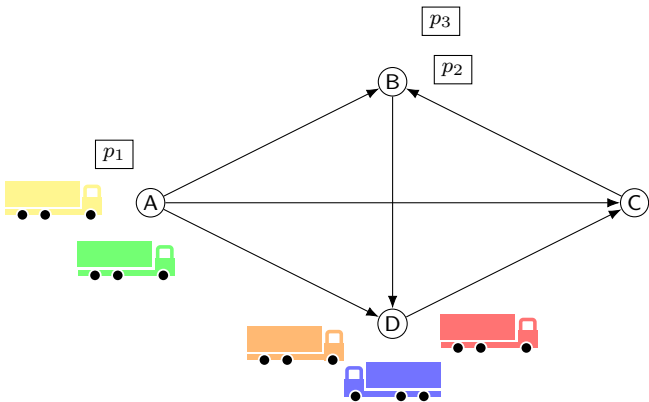


Illustration on Example Task

reduction:

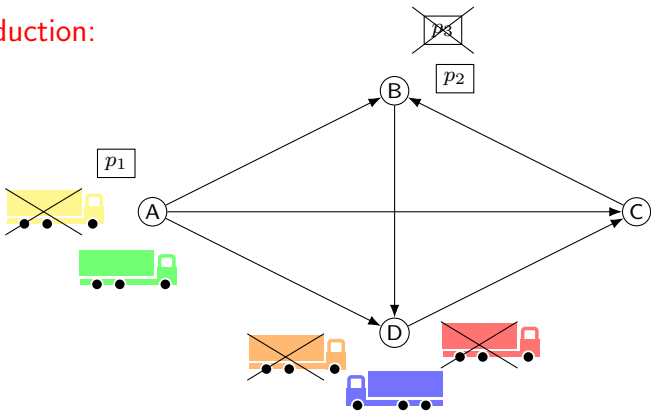
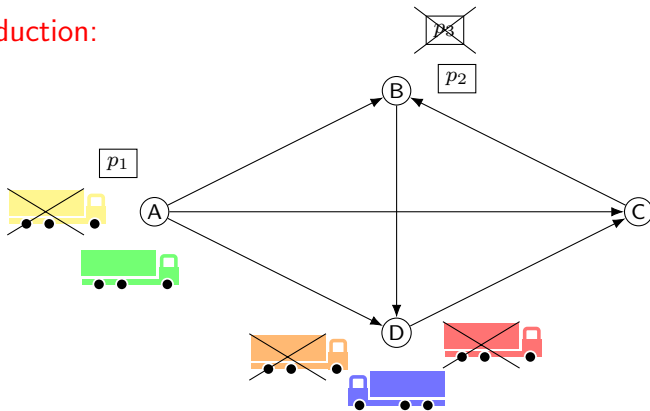


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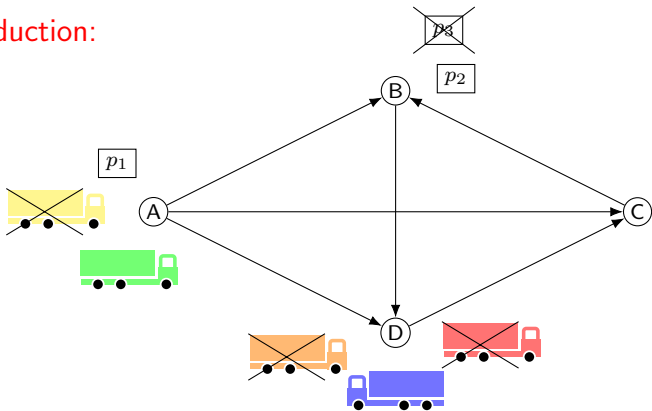
reduction:



analysis: Blue truck can reach B, C, and D

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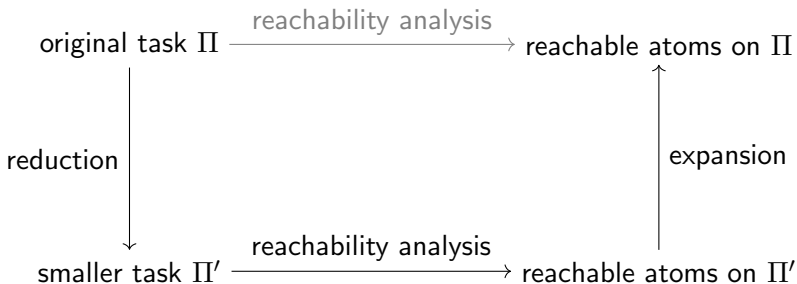
analysis: Blue truck can reach B, C, and D

expansion: Also orange and red truck can reach B, C, and D

More General Idea

original task Π $\xrightarrow{\text{reachability analysis}}$ reachable atoms on Π

More General Idea



Symmetries

- We consider a **lifted** task representation.
- As we only consider **reachability** we can **ignore the goal**.

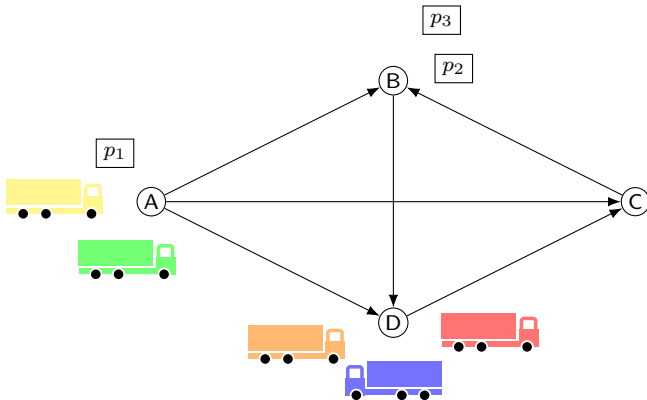
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- Two objects are **symmetric** if swapping them in the task description does not change it (up to ordering of elements).
- **Symmetric constant set**: set of **pairwise symmetric objects**

Truck Example



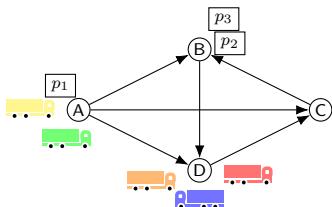
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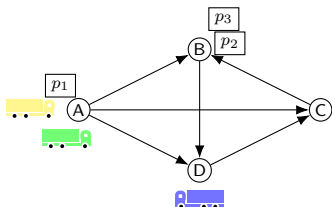


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R  (II):

$\{\dots, \dots, \dots, \dots\} \downarrow \{\dots, \dots\}$



Expansion

Expansion $E_C(L)$ extends a set of atoms L with all atoms that can be generated by permuting elements of C in a literal from L .

Example (Expansion)

$$E_{\{o_1, o_2, o_3\}}(\{P(o_1, o_2, o_2), Q(o_1, o_4)\}) = \\ \{P(o_1, o_2, o_2), P(o_1, o_3, o_3), \\ P(o_2, o_1, o_1), P(o_2, o_3, o_3), \\ P(o_3, o_1, o_1), P(o_3, o_2, o_2), \\ Q(o_1, o_4), Q(o_2, o_4), Q(o_3, o_4)\}$$

Reduction and Expansion

For symmetric constant set C and C -symmetric set of atoms

- $E_C(R_{C \downarrow C'}(L)) \subseteq L$

Reduction and Expansion

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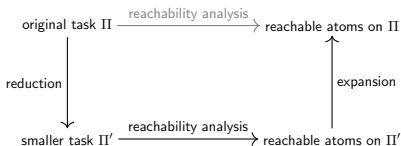
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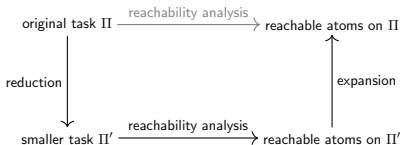
- $E_C(R_{C \downarrow C'}(L)) \subseteq L$
- $L = E_C(R_{C \downarrow C'}(L))$ for **sufficiently large C'**
 - maximal number of different constants from C in one literal

Symmetry-based Task Reduction



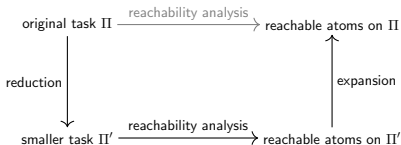
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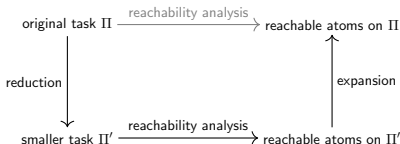
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Symmetry-based Task Reduction



- Bounds on number of elements that must be preserved from a symmetric constant set
- Overall bounds depend on reachability system.
- b_C^{lit} : **upper** bound on the number of objects from C that **can occur together** in a reachable ground **literal**
- $b_C^{\text{op}}, b_C^{\text{ax}}$: analogously for ground operators and axioms

Example: Relaxed Reachability of Literals

Definition (Relaxed Reachability of Literals)

The set of **k -reachable** ground literals ℓ ($k \in \mathbb{N}_0$) is the smallest set that contains literal ℓ if

- ℓ is true in the initial state, or
- ℓ is the default value of an axiom, or
- $k > 0$ and there is a ground operator o such that
 - o has an effect $\varphi \triangleright \ell$, and
 - each literal in φ and in $pre(o)$ is $k - 1$ -reachable, or
- there is a ground axiom $\ell \leftarrow \psi$ such that each literal in ψ is k -reachable.

Preserve $\max\{b_C^{lit}, b_C^{op}, b_C^{ax}\}$ objects from C

Example: h^2 Mutexes

Definition (Relaxed Reachability of Pairs of Literals)

For $k \in \mathbb{N}_0$, the set M_k of **k -reachable** pairs of ground literals is the smallest set that contains pair $\{\ell, \ell'\}$ if one of the following holds:

- $\ell \wedge \ell'$ is true in the initial state.
- $k > 0$ and there is a ground operator o such that
 - o has effects $\varphi \triangleright \ell$ and $\varphi' \triangleright \ell'$, and
 - ...
- ...

Preserve $\max\{b_C^{\text{lit}}, b_C^{\text{op}}, b_C^{\text{ax}}\} + b_C^{\text{lit}}$ objects from C

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- tightening the bounds
 - use logic program to compute over-approximation of relaxation
- combination of several symmetric constant sets
 - unproblematic if they are disjoint

Implementation

- translator component of Fast Downward
- grounding: use existing implementation
- h^2 mutexes: add logic program

Results

77 domains with 2518 tasks from IPC benchmarks
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Grounding:

- symmetry reduction applicable to 1004 tasks from 49 domains
- however: regular grounding is so fast that reduction and expansion is not faster

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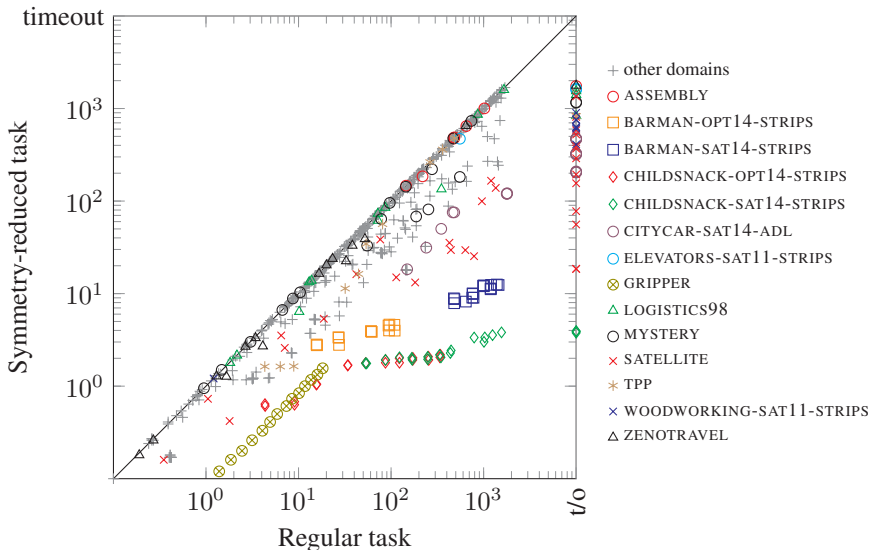
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h^2 mutexes: reduction applicable to 610 tasks from 38 domains

Results – h^2 mutexes

Summary

- With symmetric constant sets. . .
- . . . we can reduce the size of a task . . .
- . . . perform a reachability analysis on the smaller task. . .
- . . . and reconstruct the original result with an expansion.

Future Work

- Formulation for general rule-based systems

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- Rintanen (AAAI 2017):
Schematic Invariants by Reduction to Ground Invariants

Definition 14 (Limited Instantiation) For a given action set A , predicate set P , domain function D , type t , and integer $N \geq 1$, define

$$L_t^N(A, P) = \max(\max_{a \in A} prms_t(a), \max_{p \in P} prms_t(p)) \\ + (N - 1) \cdot (\max_{p \in P} prms_t(p))$$

→ Clarify relationship and applicability to
a wider range of invariant synthesis algorithms