

Subset-Saturated Cost Partitioning for Optimal Classical Planning

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- optimal classical planning
- A* search + admissible heuristic
- multiple heuristics
- cost partitioning

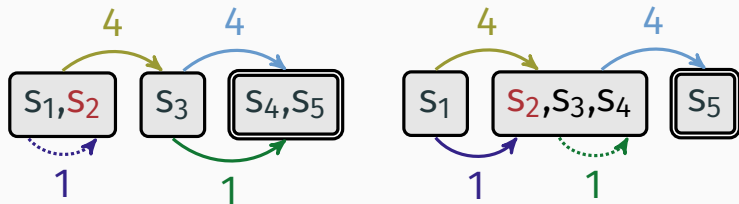
- optimal classical planning
- A* search + admissible heuristic
- multiple heuristics
- **saturated** cost partitioning

- optimal classical planning
- A* search + admissible heuristic
- multiple heuristics
- **subset-saturated** cost partitioning

Saturated cost partitioning

Saturated cost partitioning algorithm

- order heuristics, then for each heuristic h :
 - use minimum costs preserving all estimates of h
 - use remaining costs for subsequent heuristics



$$\max(h_1(s_2), h_2(s_2)) = \max(5, 4) = 5$$

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Saturated cost partitioning

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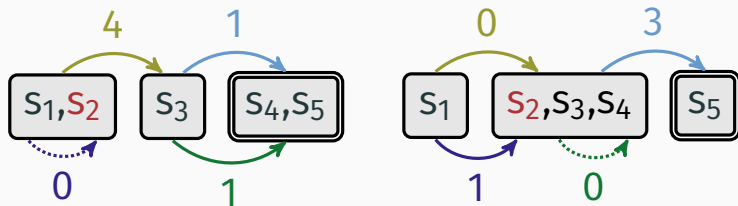
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Saturated cost partitioning

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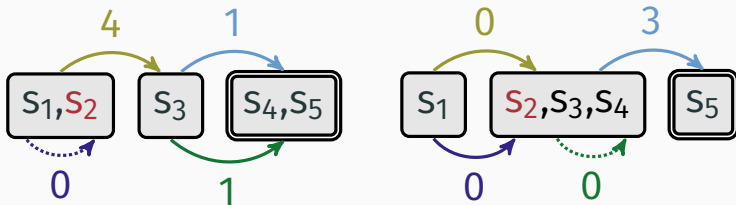
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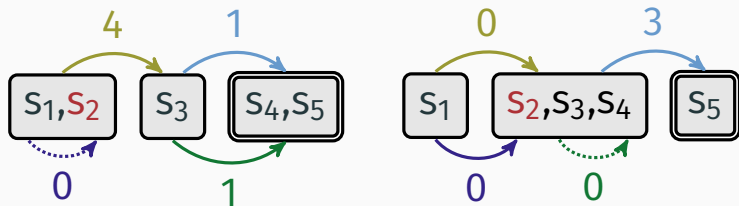
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$$h_{\langle h_1, h_2 \rangle}^{\text{SCP}}(s_2) = 5 + 3 = 8$$

Saturated cost function

1. $\text{scf}(o) \leq \text{cost}(o)$ for all operators $o \in \mathcal{O}$
2. $h(\text{scf}, s) = h(\text{cost}, s)$ for all states $s \in S$

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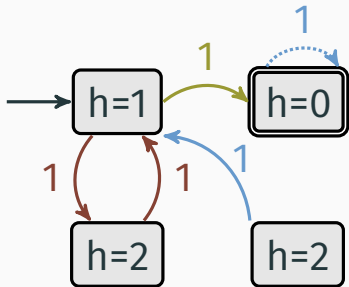
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→ 4 saturators

Saturate for all states

Saturator **all**

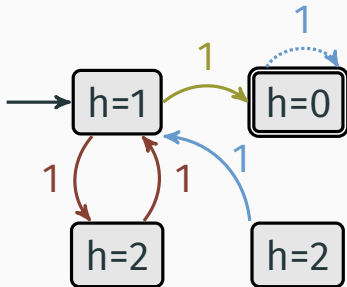
$$\text{scf}(o) = \max_{a \xrightarrow{o} b \in T} (h(a) - h(b))$$



Saturate for reachable states

Saturator **reach**

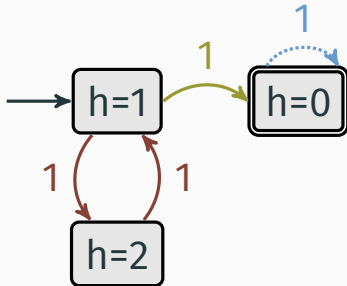
- remove unreachable states and transitions
- use **all** saturator for remaining states



Saturate for reachable states

Saturator **reach**

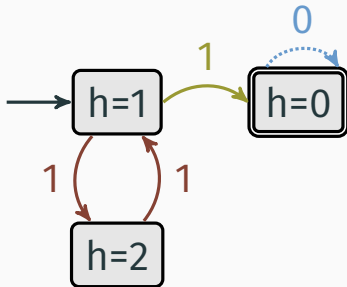
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Saturate for reachable states

Saturator **reach**

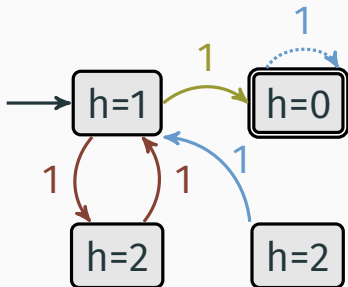
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Saturate for a perimeter

Saturator **perim**

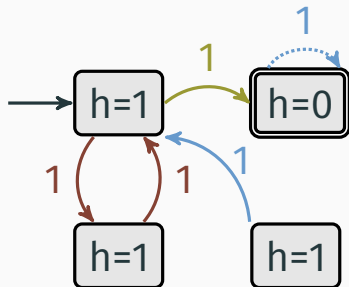
- use perimeter k around goal states ($k = h(s)$)
- cap all distances at k
- use **all** saturator



Saturate for a perimeter

Saturator **perim**

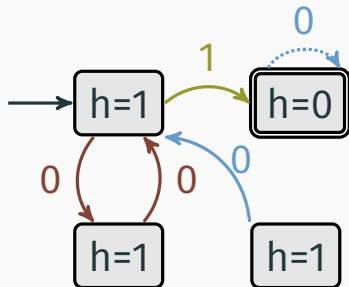
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Saturate for a perimeter

Saturator **perim**

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Saturate for a single state

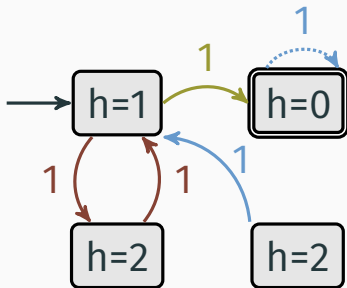
Saturator **lp**

$$H_a \leq 0 \quad \text{for all } a \in S_*$$

$$H_a \leq C_o + H_b \quad \text{for all } a \xrightarrow{o} b \in T$$

$$C_o \leq \text{cost}(o) \quad \text{for all } o \in \mathcal{O}$$

$$H_{\alpha(s)} = h(s)$$



Saturate for a single state

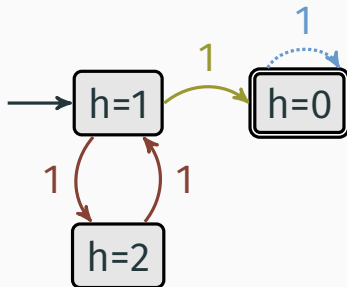
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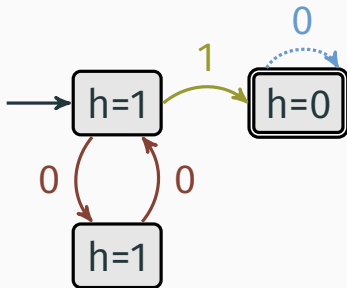
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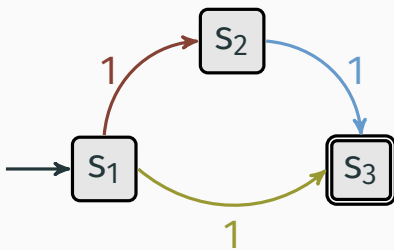
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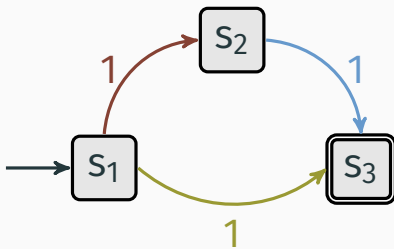
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No unique minimum saturated cost function



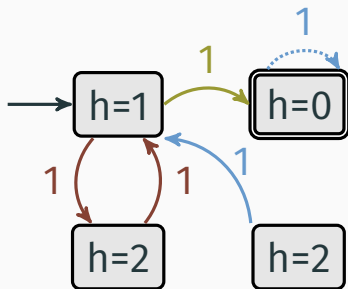
No unique minimum saturated cost function



- large Pareto frontier
- \rightarrow chain saturators

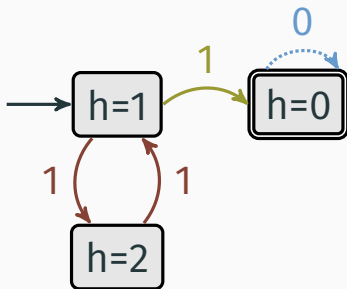
Chaining saturators

reach, perim:



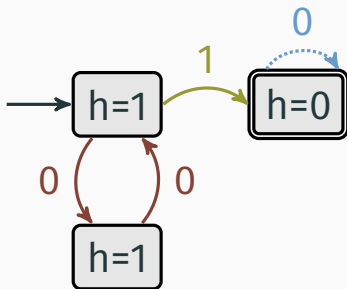
Chaining saturators

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Chaining saturators

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Experiments

- 30 minutes, 3.5 GiB
- hill climbing and systematic PDBs, Cartesian abstractions
- online and offline
- non-negative and general costs

Heuristic value for initial state – non-negative costs (1506 tasks)

	all	reach	perim	lp	all, lp	reach, lp	perim, lp
all	-	2	2	88	2	2	2
reach	15	-	5	88	2	2	2
perim	433	430	-	142	15	15	6
lp	493	490	229	-	79	79	84
all, lp	507	505	219	179	-	0	12
reach, lp	507	505	219	179	0	-	12
perim, lp	512	510	217	185	16	16	-

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Online – non-negative (nn) vs. general costs (gen)

	Coverage		Evals/sec		h(s ₀) higher	
	nn	gen	nn	gen	nn	gen
all	680	703	940.9	946.2	7	195
reach	657	679	501.4	514.1	5	222
perim	722	726	897.7	906.6	4	99
lp	360	320	8.7	6.9	493	144
all, lp	385	354	10.0	8.1	138	180
reach, lp	384	355	9.9	8.1	120	181
perim, lp	392	367	11.0	8.8	31	134

Offline subset-saturated cost partitioning

	all	reach	perim	lp	Coverage
all	-	1	9	36	1136
reach	0	-	8	36	1134
perim	1	2	-	36	1117
lp	0	0	0	-	694

Offline subset-saturated cost partitioning

	all	reach	perim	lp	perim [*]	Coverage
all	-	1	9	36	1	1136
reach	0	-	8	36	1	1134
perim	1	2	-	36	0	1117
lp	0	0	0	-	0	694
perim [*]	4	4	11	36	-	1144

Summary

- generalization of saturated cost partitioning
- three new saturators
- stronger heuristics