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**HSDIP 2015** 

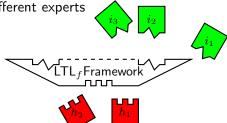
## Motivation

Motivation

### Goal

### framework for describing information about the search space

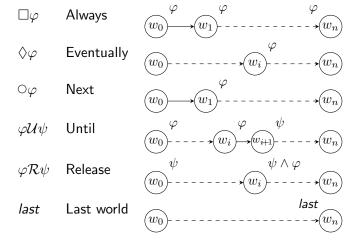
- combining information from different sources
- decoupling the derivation and exploitation of information
  - → split work between different experts



# Linear Temporal Logic on Finite Traces (LTL<sub>f</sub>)

- evaluated over a linear sequence of worlds (= variable assignments)
- extends propositional logic with:

•000000



## What if we only know the beginning of the sequence?

## Definition (Progression)

For an LTL $_f$  formula  $\varphi$  and a world sequence  $\langle w_0,\ldots,w_n\rangle$  with n>0 it holds that  $\langle w_1,\ldots,w_n\rangle\models \operatorname{progress}(\varphi,w_0)$  iff  $\langle w_0,\ldots,w_n\rangle\models \varphi.$ 

## Example

 $\mathsf{progress}\big(a \land \bigcirc e \land \Box(c \lor d) \land (b \ \mathcal{U}d), \{a, d\}\big) = e \land \Box(c \lor d)$ 

# LTL<sub>f</sub> Formulas in the Search Space

variable	$\leftrightarrow$	STRIPS variable or action
world	$\leftrightarrow$	node in search space (with incoming action)
world sequence	$\leftrightarrow$	path to a goal node

### $\mathsf{LTL}_f$ formulas associated to nodes

ightarrow express conditions all optimal paths to a goal need to fulfill



## Definition (Feasibility for nodes)

An LTL<sub>f</sub> formula  $\varphi$  is feasible for n if for all paths  $\rho$  such that

- ullet  $\rho$  is applicable in n,
- ullet the application of ho leads to a goal state  $(G\subseteq s[
  ho])$ , and
- $g(n) + c(\rho) = h^*$

it holds that  $w_{\rho}^{s} \models \varphi$ .

(where  $\boldsymbol{w_{a}^s} = \langle \{a_1\} \cup s[a_1], \{a_2\} \cup s[\langle a_1, a_2 \rangle], \dots, \{a_n\} \cup s[\rho], s[\rho] \rangle$ )

# Adding and Propagating Information during the Search

How can we add/propagate information while preserving feasibility?

- new information during the search directly added to the corresponding node with conjunction
- of formulas can be propagated with progression to successor nodes

### $\mathsf{Theorem}$

Let  $\varphi$  be a feasible formula for a node n, and let n' be the successor node reached from n with action a. Then  $\operatorname{progress}(\varphi, \{a\} \cup s(n'))$  is feasible for n'.

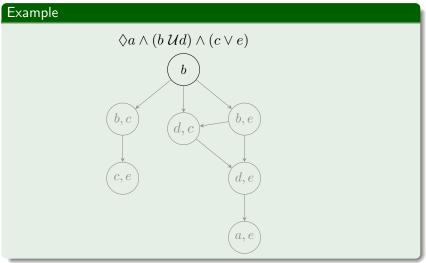
# Adding and Propagating Information during the Search

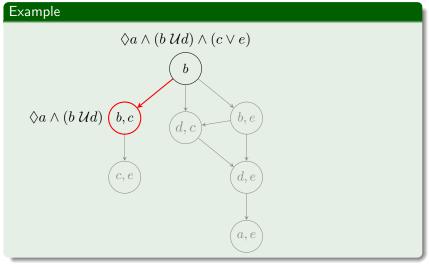
How can we add/propagate information while preserving feasibility?

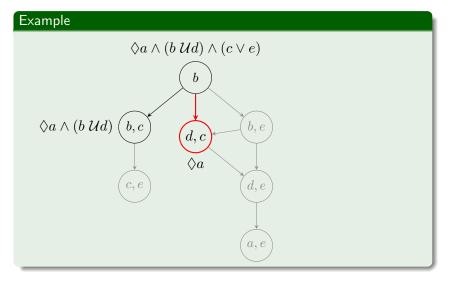
4 duplicate elimination: conjunction of formulas of "cheapest" nodes is feasible

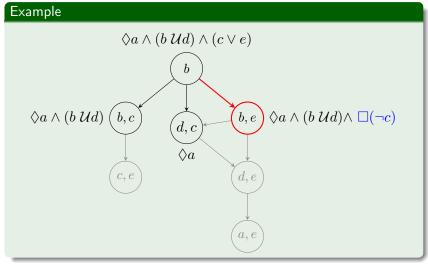
### $\mathsf{Theorem}$

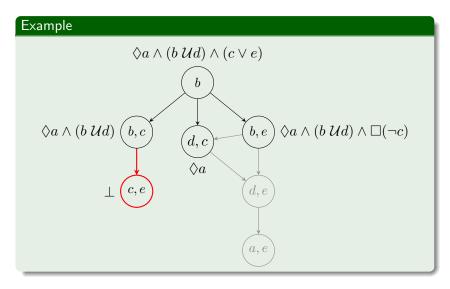
Let n and n' be two search nodes such that g(n) = g(n') and s(n) = s(n'). Let further  $\varphi_n$  and  $\varphi_{n'}$  be feasible for the respective node. Then  $\varphi_n \wedge \varphi_{n'}$  is feasible for both n and n'.

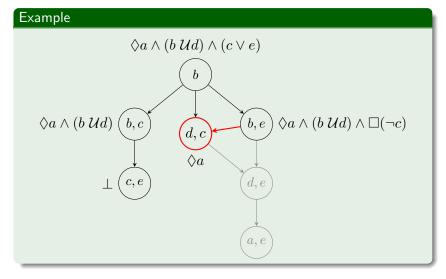


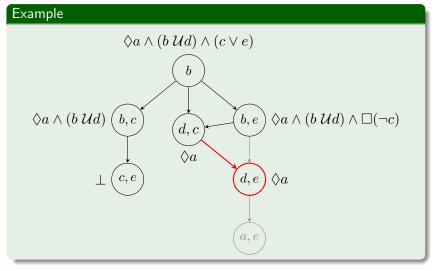


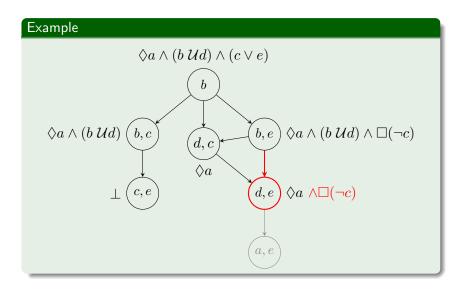


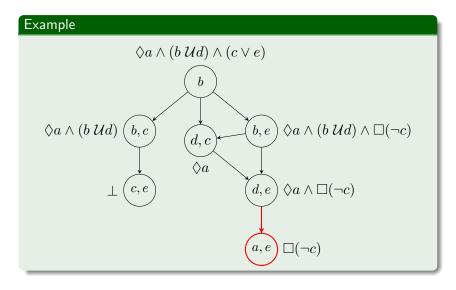












# Encoding Information in LTL<sub>f</sub> Formulas

### Possible sources of information:

- domain-specific knowledge
- temporally extended goals
- here: information used in specialized heuristics
  - Landmarks and their orderings (Hoffmann et al. 2004, Richter et al. 2008)
  - Unjustified Action Applications (Karpas and Domshlak 2012)

## Landmarks

Fact Landmark: A fact that must be true at some point in every plan (Hoffmann et al. 2004)

$$\rightarrow$$
 In LTL<sub>f</sub>:  $\Diamond l$ 

Further information about landmarks:

- First achievers:  $l \vee \bigvee_{a \in \mathit{FA}_l} \lozenge a$
- Required again:  $(\lozenge l)\mathcal{U}l'$
- Goal:  $\bigwedge_{g \in G} ((\lozenge g) \mathcal{U} \bigwedge_{g' \in G} g')$

# Unjustified Action Applications

If an action is applied, its effects must be of some use (Karpas and Domshlak 2012)

- one of its effects is necessary for applying another action
- ② one of its effects is a goal variable (that is not made false again)

$$\begin{split} \varphi_a &= \bigvee_{e \in add(a) \backslash G} \left( (e \land \bigwedge_{\substack{a' \in A \text{ with } \\ e \in add(a')}} \neg a') \mathcal{U} \bigvee_{\substack{a' \in A \text{ with } \\ e \in pre(a')}} a' \right) \lor \\ \bigvee_{\substack{e \in add(a) \cap G}} \left( (e \land \bigwedge_{\substack{a' \in A \text{ with } \\ e \in add(a')}} \neg a') \mathcal{U} \big( \text{last} \lor \bigvee_{\substack{a' \in A \text{ with } \\ e \in pre(a')}} a' \big) \big) \end{split}$$

## Heuristics

- Very rich temporal information possible
  - → heuristics can use different levels of relaxation
- Proof of concept heuristic extracts landmarks from node-associated formulas
  - → looses temporal information between landmarks

# Extracting Landmarks from the Formula

- Convert formula into positive normal form ("¬" only before atoms)
  - can be computed efficiently
  - progression preserves positive normal form
- ② Transform formula into implied formula where ◊ in front of every literal, no other temporal operators
- Transform formula into CNF
- Oismiss clauses which are true already in current state
- 5 Extract disjunctive action landmarks from individual clauses

# Experiment Setup

### Configurations:

- $h_{\rm LA}$ : standard admissible landmark heuristic (Karpas and Domshlak 2009)
- $b_{AL}^{LM}$ : LTL landmark extraction heuristic with sources:
  - Landmarks (First achievers, Required again, Goal)
- f b  $h_{
  m AL}^{
  m LM+UAA}$ : LTL landmark extraction heuristic with sources:
  - Landmarks (First achievers, Required again, Goal)
  - Unjustified Action Applications
  - all heuristics use BJOLP landmark extraction and optimal cost partitioning
  - ullet search algorithm:  $h_{
    m LA}$  uses  ${
    m LM-A^*},$  the others a slight variant we call  ${
    m LTL-A^*}$

# Coverage

	$h_{ m LA}$	$h_{ m AL}^{ m LM}$	$h_{ m AL}^{ m LM+UAA}$
airport (50)	31	28	26
elevators-08 (30)	14	14	13
floortile (20)	2	2	4
freecell (80)	52	51	50
mprime (35)	19	19	20
nomystery (20)	18	17	16
openstacks-08 (30)	14	12	12
openstacks-11 (20)	9	7	7
parcprinter-08 (30)	15	14	14
parcprinter-11 (20)	11	10	10
pipesworld-tan (50)	9	10	10
scanalyzer-08 (30)	10	9	9
sokoban-08 (30)	22	21	22
tidybot (20)	14	14	13
other domains (931)	483	483	483
Sum (1396)	723	711	709

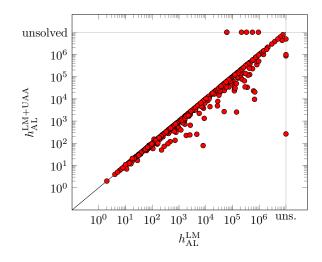
# Memory Consumption

 $h_{
m LA}$  looses no task due to memory limit, but  $h_{
m AL}^{
m LM}$  11 in total

- ullet airport: over 300% of memory usage compared to  $h_{\mathrm{LA}}$
- average: 120%
- approx. half the domains < 100%

# Impact of Unjustified Action Applications

Comparison of expansions between  $h_{
m AL}^{
m LM}$  and  $h_{
m AL}^{
m LM+UAA}$ :



## Conclusion

- ullet associate nodes in the search space with LTL $_f$  formulas
- ightarrow conditions for optimal plan
- separation between finding information and exploiting information
- allows to easily combine information from different sources
- concrete examples in this paper:
  - finding information: landmarks and unjustified action applications
  - exploiting information: extracting landmarks

## Future Work

- better informed heuristics (less relaxation)
- describe other kinds of information
  - PDDL 3 trajectory constraints
  - flow-based heuristics (van den Briel et al. 2007; Bonet 2013; Pommerening et al. 2014)
  - mutex information
- strengthening other heuristics with the information of LTL<sub>f</sub> trajectory constraints