

Oxiflex - A Constraint Programming Solver for MiniZinc written in Rust

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15.07.2024

Constraint Programming Solver for MiniZinc written in Rust

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Rust



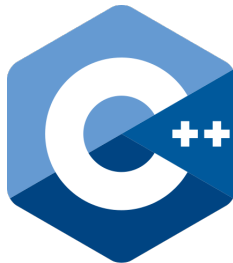
Why Rust?

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Possible Languages



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- Performance
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- Learn Rust

Constraint Programming Solver for MiniZinc written in Rust

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Constraint Network

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- Variables
 - Values to choose from
- Constraints
 - Rules for choosing values

Simple Example

Variables:

$$w \in \{1, 2, 3, 4\}$$

$$y \in \{1, 2, 3, 4\}$$

$$x \in \{1, 2, 3\}$$

$$z \in \{1, 2, 3\}$$

Constraints:

$$w = 2 \cdot x$$

$$w < z$$

$$y > z$$

Constraint Programming

```
var 1..4: w;  
var 1..4: y;  
var 1..3: x;  
var 1..3: z;  
  
constraint w = 2 * x;  
constraint w < z;  
constraint y > z;
```

Constraint Programming

```
var 1..4: w;  
var 1..4: y;  
var 1..3: x;  
var 1..3: z;  
  
constraint w = 2 · x;  
constraint w < z;  
constraint y > z;  
  
solve satisfy;
```

Constraint Programming

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var 1..4: y;  
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solve satisfy;  
  
→ MiniZinc!
```

MiniZinc





compilation



FlatZinc

```
array [1..2] of int: x_introduced_2_ = [1,-2];
```

```
array [1..2] of int: x_introduced_3_ = [1,-1];
```

```
array [1..2] of int: x_introduced_4_ = [-1,1];
```

```
var 2..4: w:: output_var;
```

```
var 1..4: y:: output_var;
```

```
var 1..3: x:: output_var;
```

```
var 1..3: z:: output_var;
```

```
constraint int_lin_eq(x_introduced_2_, [w,x], 0);
```

```
constraint int_lin_le(x_introduced_3_, [w,z], -1);
```

```
constraint int_lin_le(x_introduced_4_, [y,z], -1);
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```
solve satisfy;
```

FlatZinc

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FlatZinc constraint example

```
predicate int_lin_eq(array [int] of int: as,  
                    array [int] of var int: bs,  
                    int: c)
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FlatZinc constraint example

predicate `int_lin_eq`(array [int] of int: as,
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int: c)

$$c = \sum_i as[i] \cdot bs[i]$$

FlatZinc constraint example

predicate `int_lin_eq`(array [int] of int: as,
array [int] of var int: bs,
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$$c = \sum_i as[i] \cdot bs[i]$$

$$w = 2 \cdot x$$

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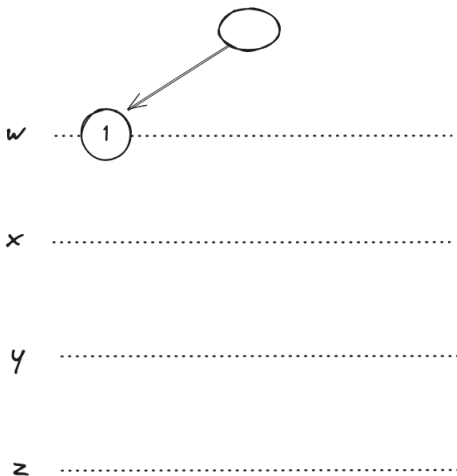
$$w = 2 \cdot x$$

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Backtracking

Backtracking Example



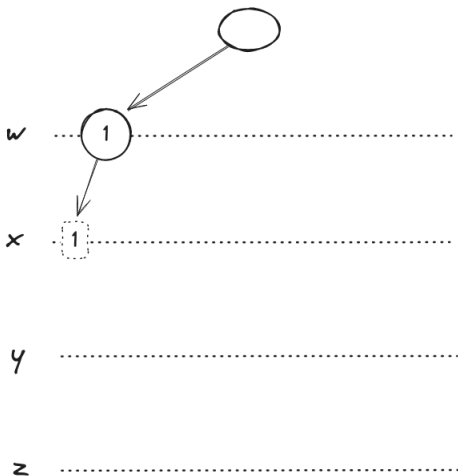
Constraints:

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$$w < z$$

$$y > z$$

Backtracking Example



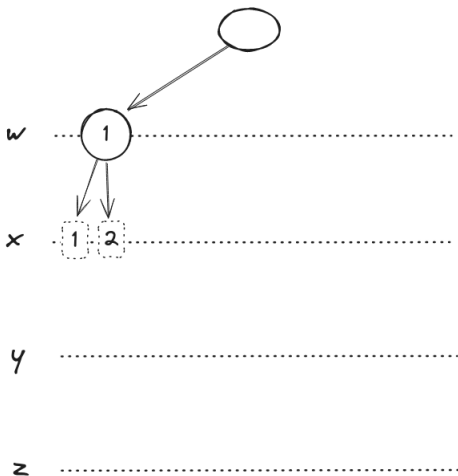
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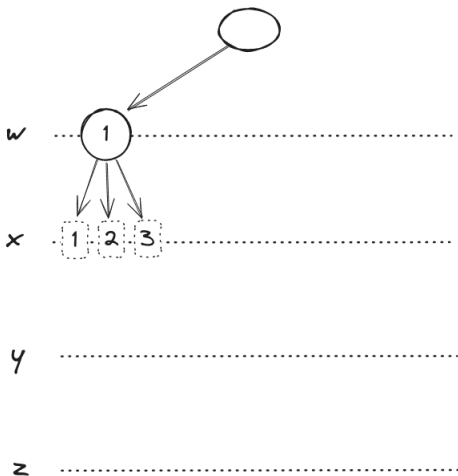
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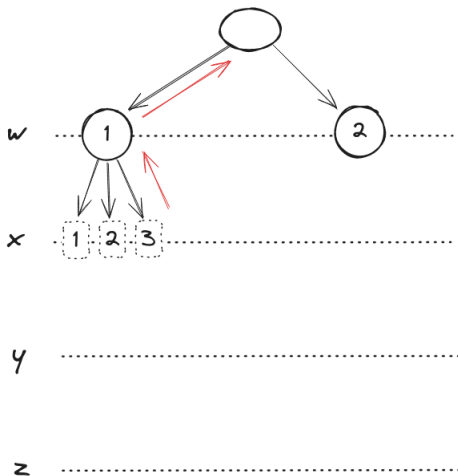
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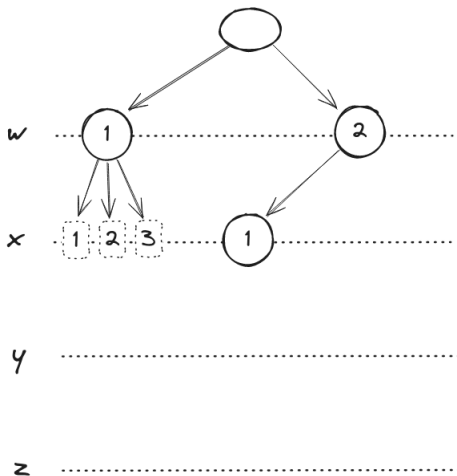
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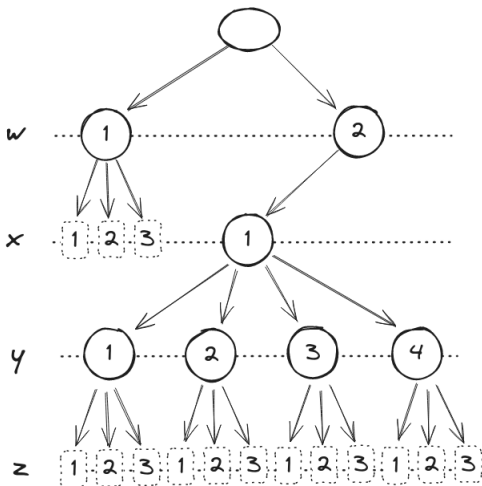
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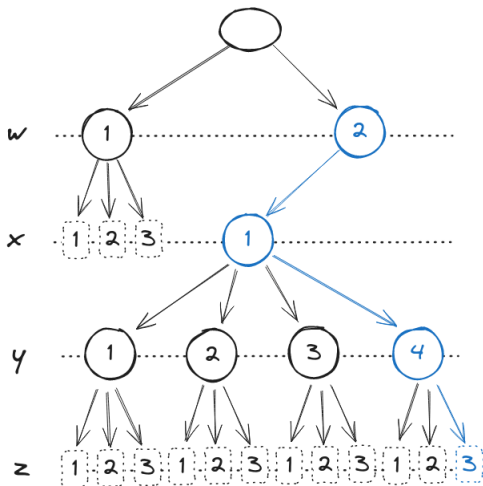
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Backtracking Example



Constraints:

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$$y > z$$

Solution:

$$w = 2$$

$$x = 1$$

$$y = 4$$

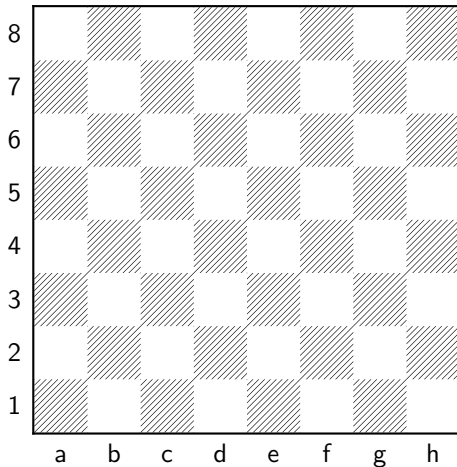
$$z = 3$$

Kinda like search...

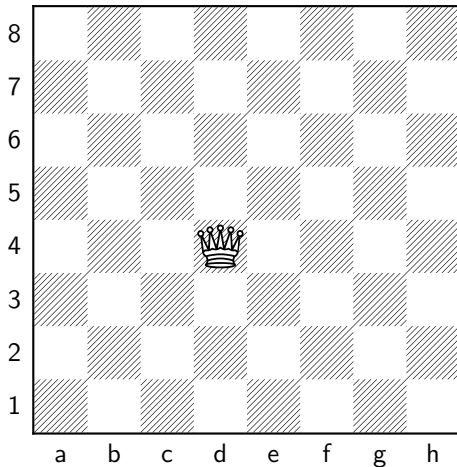
Can we do better?

8-Queens Problem

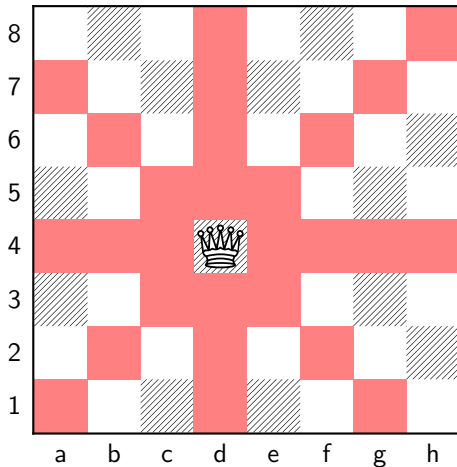
Chessboard



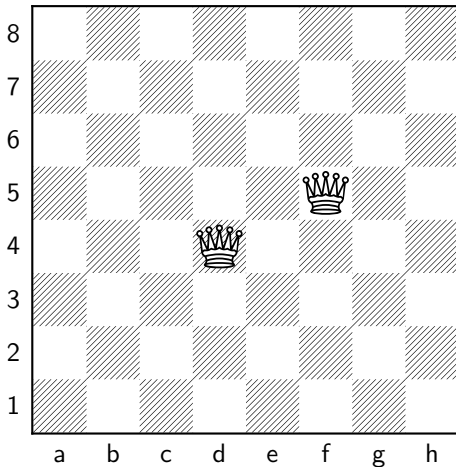
One Queen



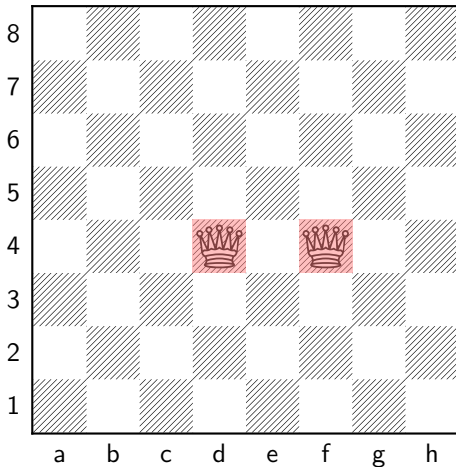
One Queen



Two Queens



Two Queens



N-Queens Problem

- Scalable

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 - 8-Queens → N-Queens

N-Queens Problem

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 - $n \times n$ chessboard

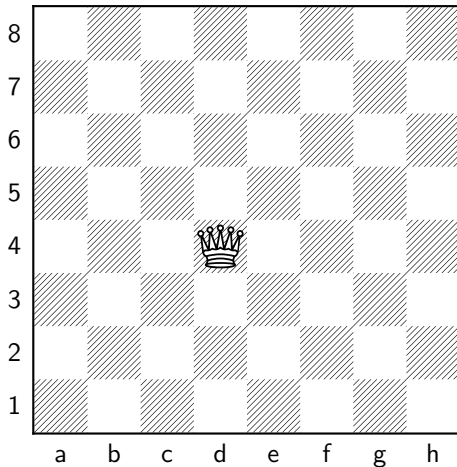
N-Queens Problem

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 - $n \times n$ chessboard
 - n Queens

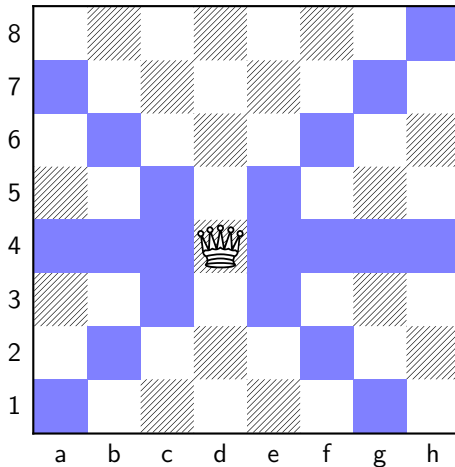
Inference

Forward Checking

Forward Checking Example

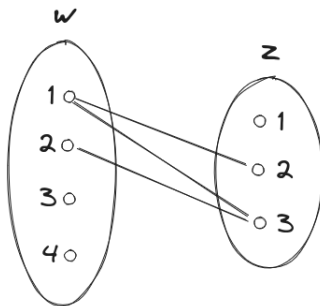


Forward Checking Example

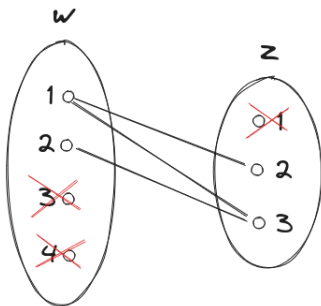


Arc Consistency

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Improvements

- Inference

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- Inference
 - Forward Checking
 - Arc consistency
 - AC-1

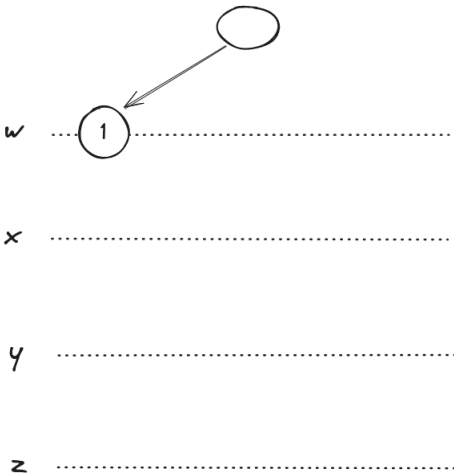
Improvements

- Inference
 - Forward Checking
 - Arc consistency
 - AC-1
 - AC-3

Improvements

- Inference
 - Forward Checking
 - Arc consistency
 - AC-1
 - AC-3
- Dynamic Variable Ordering

Dynamic Variable Ordering



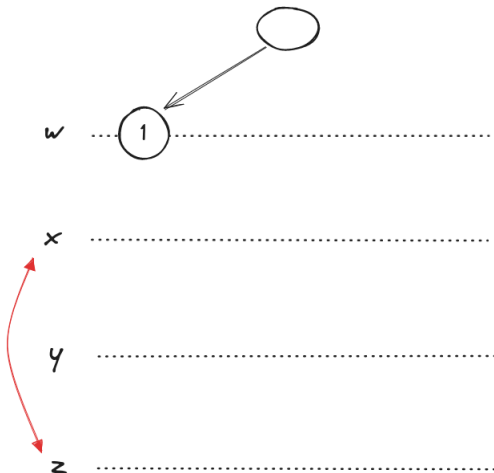
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Dynamic Variable Ordering



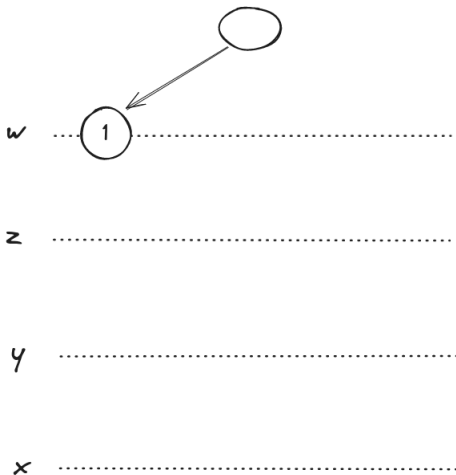
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Dynamic Variable Ordering



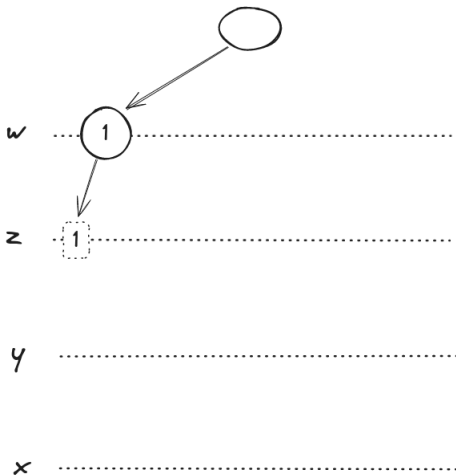
Constraints:

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Dynamic Variable Ordering



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Oxiflex

Demo

Limitations

- FlatZinc builtins

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 - IntLinEq
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- FlatZinc builtins
 - IntLinEq
 - IntLinLe
 - IntLinNe
- No floating points
- No minimize / maximize
- Only one solution

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Benchmarks

N-Queens Problem

Queens Time

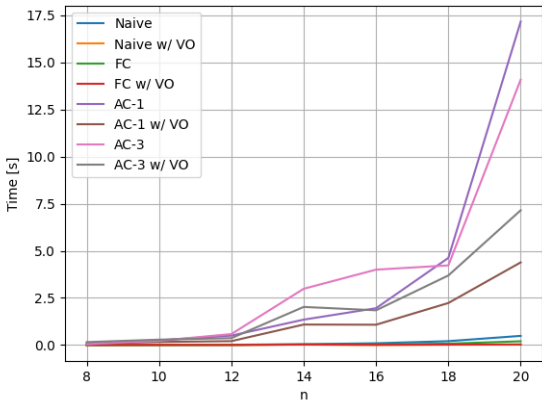


Figure: Averaged over > 10 runs

Queens Time

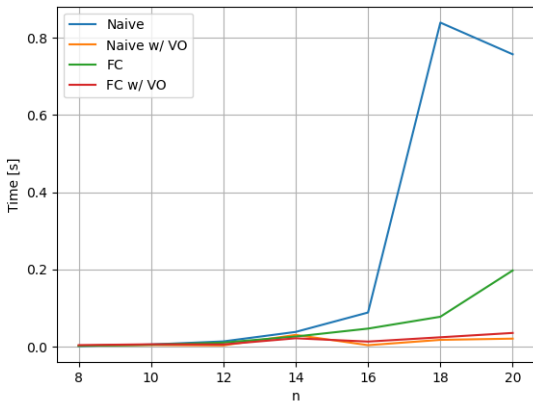


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Queens Iterations

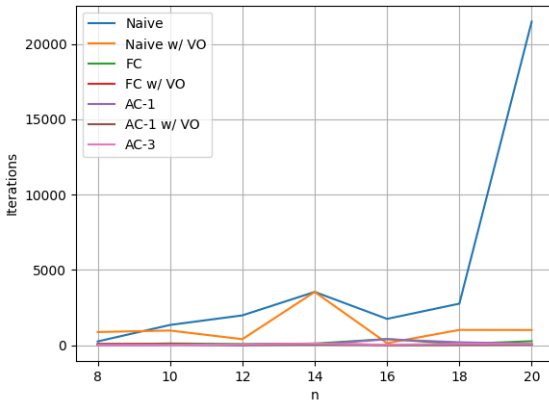


Figure: Averaged over 5 runs

Queens Iterations

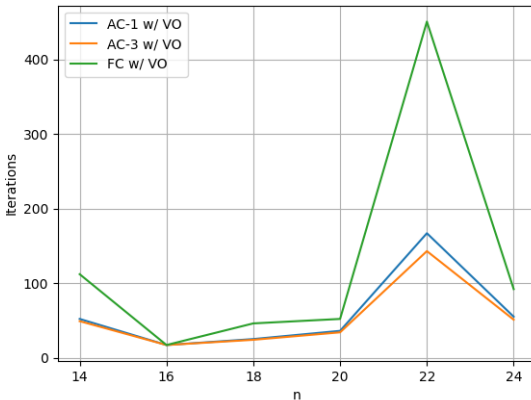


Figure: Averaged over 5 runs

Slow Convergence

Slow Convergence for small n

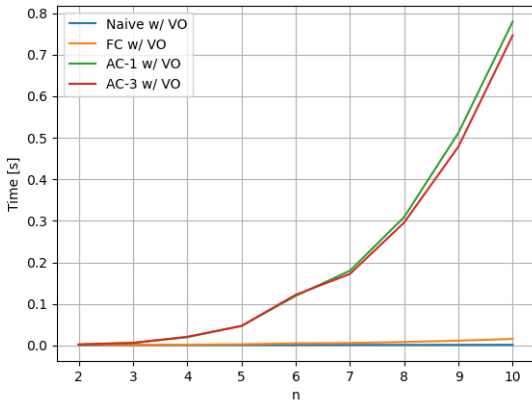


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Slow Convergence for small n

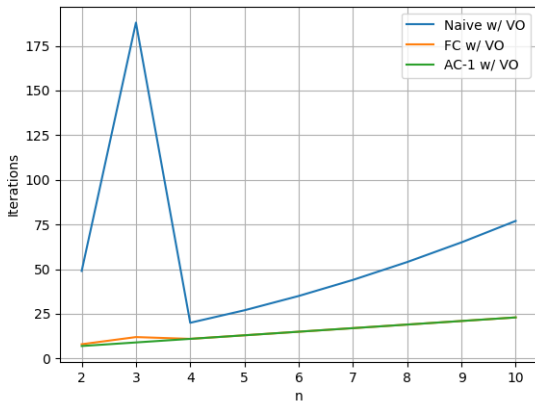


Figure: Averaged over 5 runs

Slow Convergence Comparison

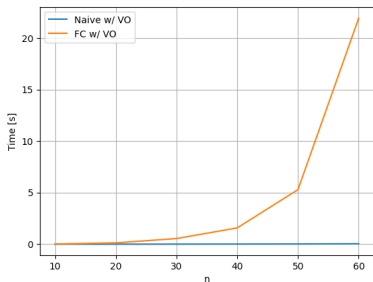


Figure: Averaged over > 10 runs

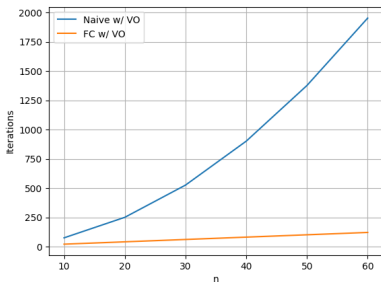


Figure: Averaged over 5 runs

Oxiflex vs Gecode

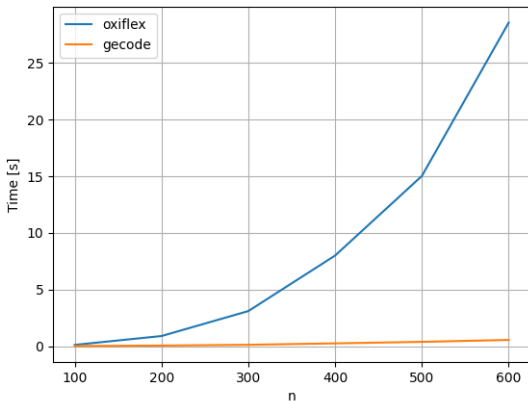


Figure: Averaged over > 10 runs

Conclusion

The end