

# Double Description Method in Cost Partitioning

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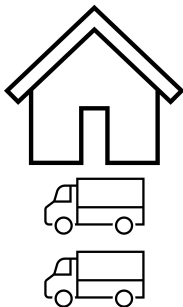
November 11, 2 × 3 × 337

# Planning

# Logistics Example



# Logistics Example

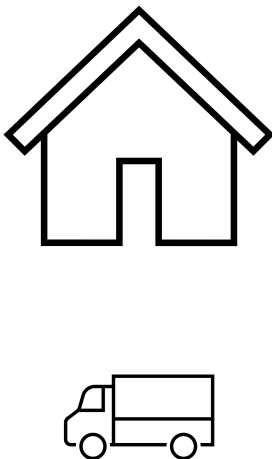
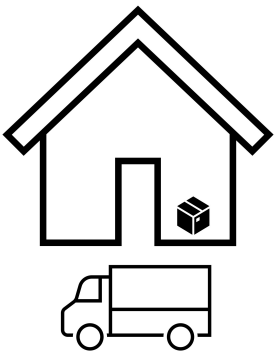


(informal) task description

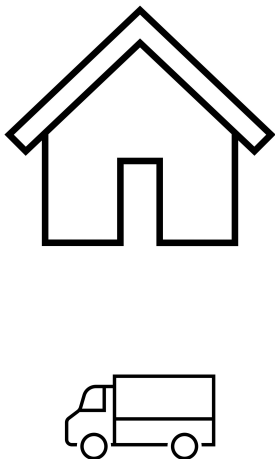
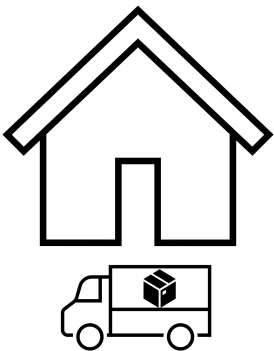
**actions:** trucks can drive from one location to the other and (un-)load package

**goal:** find sequence of actions such that package is at other location

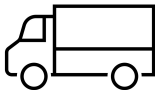
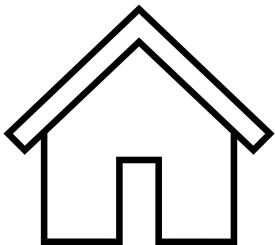
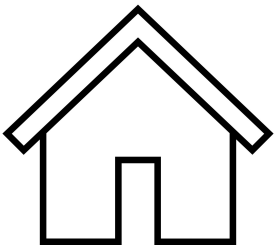
# Logistics Example



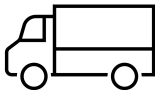
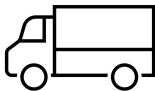
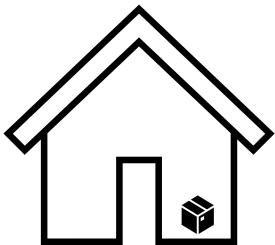
# Logistics Example



# Logistics Example



# Logistics Example





# Finding Plans

## Heuristic Search

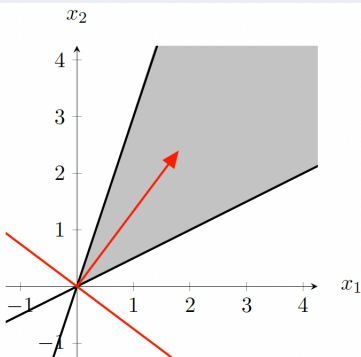
- **cost partitioning** over abstraction heuristics
- calculating (optimal) cost partitioning involves solving a large **linear program**  $\Rightarrow$  computationally expensive

# Linear Programs

## Example

$$\begin{aligned} \max \quad & \frac{3}{4}x_1 + x_2 \quad \text{s.t.} \\ & \frac{1}{2}x_1 - x_2 \leq 0 \\ & -3x_1 + x_2 \leq 0 \end{aligned}$$

## Solution Space



# Linear Programs

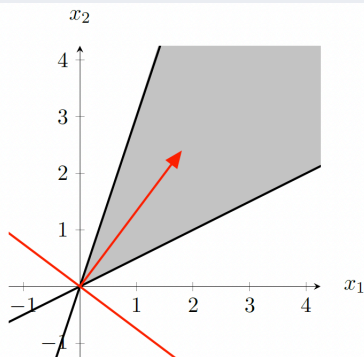
## Example

$$\begin{aligned} \max \quad & \frac{3}{4}x_1 + x_2 \quad \text{s.t.} \\ & \frac{1}{2}x_1 - x_2 \leq 0 \\ & -3x_1 + x_2 \leq 0 \end{aligned}$$

## Generating Rays

Every solution can be written as a finite sum of generating rays

## Solution Space



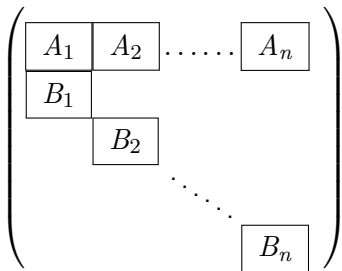
# Decomposition

# Dantzig-Wolfe Decomposition

$$\left( \begin{array}{cccc} A_1 & A_2 & \dots & A_n \\ B_1 & & & \\ & B_2 & & \\ & & \dots & \\ & & & B_n \end{array} \right)$$

- solves linear programs with special structure
- starts with linear program that uses less columns
- iteratively adds columns that improve solution
- to know which columns to add, it solves the **pricing problem**

# Dantzig-Wolfe in Cost Partitioning



- $B_i$ 's are the abstraction heuristics used in cost partitioning
- $A_i$ 's form the cost partitioning constraints

## Pricing Problem

Minimize  $c(y) - h$  subject to  
 $h \leq$  heuristic  $i$  under cost  $c$

⇒ one pricing problem per abstraction

# Pricing Problem in Cost Partitioning

## Constraints

$$d_0 = 0$$

$$d_t \leq d_s + c_\ell \quad \text{for all transitions from state } s \text{ to } t \text{ with cost } c_\ell$$

$$h \leq d_{s^*} \quad \text{for all goal states } s^*$$

## Example

$$d_0 = 0$$

$$d_0 \leq d_2 + c_0$$

$$d_1 \leq d_0 + c_0$$

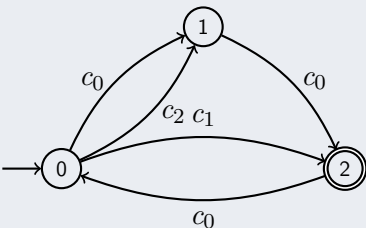
$$d_1 \leq d_0 + c_2$$

$$d_2 \leq d_0 + c_1$$

$$d_2 \leq d_1 + c_0$$

$$h \leq d_2$$

## Graph



# Finding Generating Rays



# Methods

## Double Description

Used to calculate the generating rays of linear constraints

## Fourier-Motzkin Elimination

Used to project out variables of linear constraints

# Strategies

**Pricing Problem  
(constraints over  
h,c,d)**

*project out distance  
variables (with FM)*

**constraints  
over h,c**

*compute  
gen. rays  
(with DD)*

*compute  
gen. rays  
(with DD)*

**generating rays  
over h,c,d**

*project out distance  
variables (ignore)*

**generating  
rays over h,c**

# Projection - A Closer Look

## Fourier-Motzkin Elimination

- 1 choose variable to project out

### Example

$$d_0 = 0$$

$$d_0 \leq d_2 + c_0$$

$$d_1 \leq d_0 + c_0$$

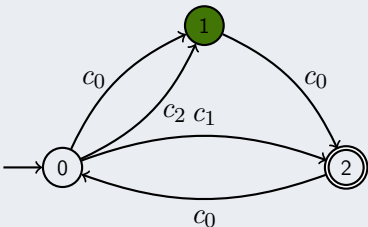
$$d_1 \leq d_0 + c_2$$

$$d_2 \leq d_0 + c_1$$

$$d_2 \leq d_1 + c_0$$

$$h \leq d_2$$

### Graph



# Projection - A Closer Look

## Fourier-Motzkin Elimination

- group constraints with respect to chosen variable

### Example

$$d_0 = 0$$

$$d_0 \leq d_2 + c_0$$

$$d_1 \leq d_0 + c_0$$

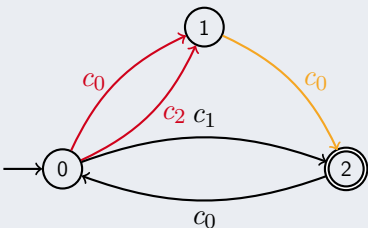
$$d_1 \leq d_0 + c_2$$

$$d_2 \leq d_0 + c_1$$

$$d_2 \leq d_1 + c_0$$

$$h \leq d_2$$

### Graph



# Projection - A Closer Look

## Fourier-Motzkin Elimination

- combine constraints from different groups  $\Rightarrow$  new constraints without chosen variables

### Example

$$d_0 = 0$$

$$d_0 \leq d_2 + c_0$$

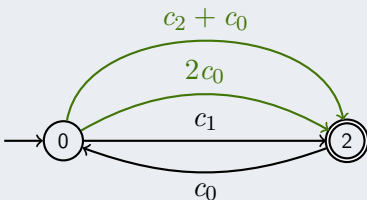
$$d_2 \leq d_0 + c_0 + c_0$$

$$d_2 \leq d_0 + 2c_0$$

$$d_2 \leq d_0 + c_1$$

$$h \leq d_2$$

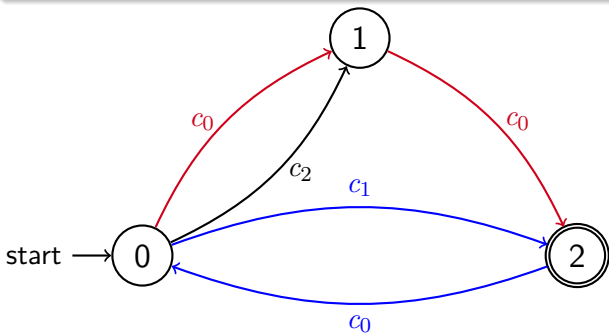
### Graph



# Potential Shortcut for Projection

## Observation

- nodes get eliminated  $\Rightarrow$  new edges
- edges represent **open** or **closed walks** in the original graph



# Potential Shortcut for Projection

Applying Fourier-Motzkin elimination to all distance variables  $\Rightarrow$   
constraints represent open walks from start to goal or closed walks  
in the original graph...

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... **but it gets even better**



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Applying Fourier-Motzkin elimination to all distance variables  $\Rightarrow$  constraints represent open walks from start to goal or closed walks in the original graph...

... **but it gets even better**

## Theorem

The constraint system representing all **simple paths** and **simple cycles** has the same solution space as the original pricing problem.

# Strategies

**Pricing Problem  
(constraints over  
h,c,d)**

*compute  
gen. rays  
(with DD)*



**generating rays  
over h,c,d**

*project out distance  
variables (with FM)*



*project out distance  
variables (calc. simple  
cycles and simple paths)*

**constraints  
over h,c**

*compute  
gen. rays  
(with DD)*



**generating  
rays over h,c**



*project out distance  
variables (ignore)*

# Experiments

# Setup

- 1590 tasks from International Planning Competition
- considered projections onto one or two variables for every task
- run involved calculating all generating rays for all projections onto one or two variables of one task
- time limit of 5 min per run
- memory limit of 2 GiB per run
- experiment run on Intel Xeon Silver 4114

# Number Of Solved Tasks

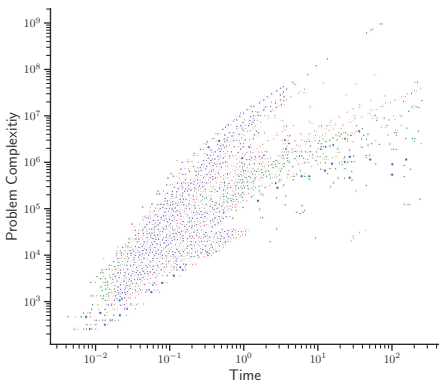
SYS1	Solved	Time Limit Reached	Memory Limit Reached
DDPROJ	1398	<b>192</b>	—
FMDD	1560	9	<b>21</b>
SCDD	<b>1590</b>	—	—

SYS2	Solved	Time Limit Reached	Memory Limit Reached
DDPROJ	163	<b>1427</b>	—
FMDD	196	327	<b>1107</b>
SCDD	351	<b>1013</b>	226

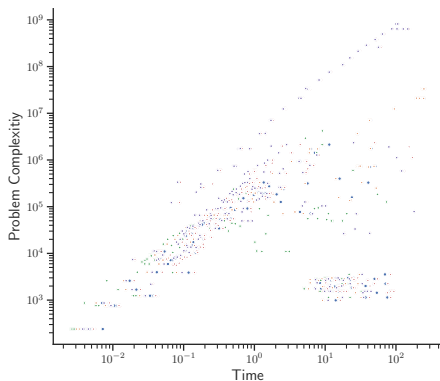
# Complexity vs. Runtime

**Idea:** complexity can be measured by multiplying number of constraints (at the beginning) by the number of variables

## SYS1



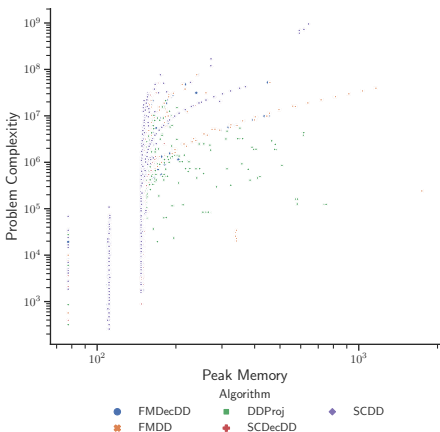
## SYS2



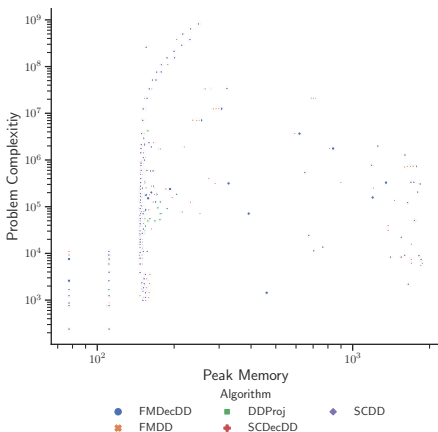
# Complexity vs. Peak Memory Consumption

**Idea:** complexity can be measured by multiplying number of constraints (at the beginning) by the number of variables

## SYS1

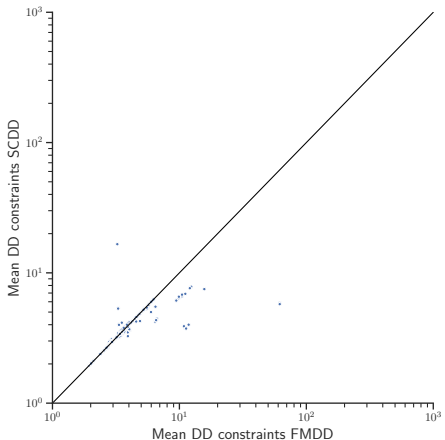


## SYS2

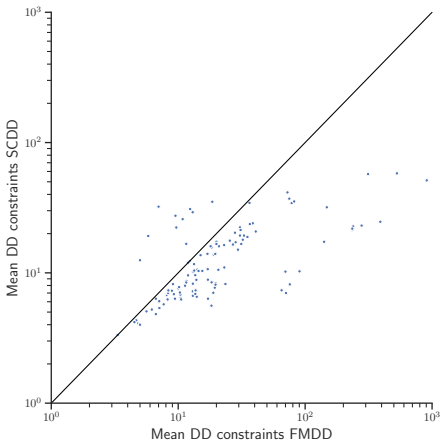


# Redundant Constraints

## SYS1



## SYS2





# Conclusion and Outlook

# Conclusion

- projecting first seems to be the superior strategy
- projecting by calculating simple paths and simple cycles seems to outperform (naive) Fourier-Motzkin
- our measure of complexity seems to be a good indicator for runtime but not for peak memory consumption

# Outlook

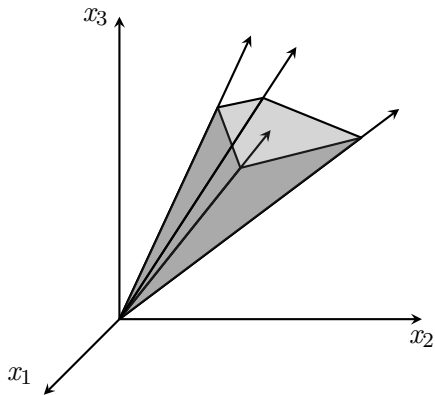
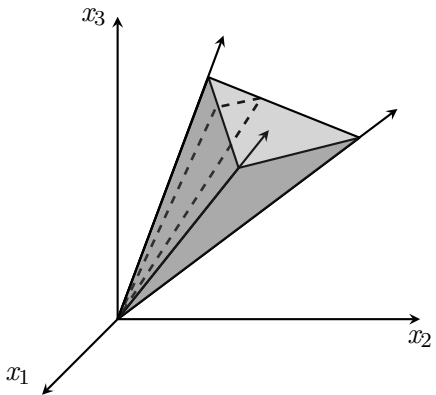
- Does precomputing the generating rays improve performance of cost partitioning?
- detecting and removing redundant constraints
- interesting generating rays

## Pricing Problem

Minimize  $c(y) - h$  subject to  
 $h \leq$  heuristic  $i$  under cost  $c$

- decomposing solution space

# Double Description Method



# Decomposition Polyhedral Cones

