Safe Abstraction in Fast Downward

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Outline

Background

Planning and abstractions

Safe Abstraction

Safe abstractions in Fast Downward

Evaluation

Effect of safe abstraction on search

Conclusion

What have we learned? What is left to do?

Background Planning and abstractions

Planning – Trucks Problem



Planning – Variables

Variables in the trucks problem:

- Location of truck *truck* (*at-A, at-B*)
- Location of package p
- Occupancy of truck cargo

p (at-A, at-B, in truck) cargo (empty, contains-p) A state is a value assignment over all variables.

In this case: $truck \mapsto at -A$ $p \mapsto at -A$ $cargo \mapsto empty$



Planning – Operators

Operators in the trucks problem:

- Driving the truck (a to b, b to a)
- Drop the package (at a, at b)
- Pick-up the package (at a, at b)



Example: *pick-up-truck-a-p* Preconditions:

- $truck \mapsto at-A \checkmark$
- $p \mapsto at A \checkmark$
- *cargo* → *empty* ✓ Effects:
- $p \mapsto in$ -truck
- $cargo \mapsto contains-p$

Planning – Goal

Goal is a partial state

Goal in the trucks problem: $p \mapsto at-B$ A **plan** is the sequence of operators used to get to a goal

- pick-up-truck-a-p
- drive-truck-a-b
- drop-truck-b-p
- drive-truck-b-a



Abstraction

Problem:

• State space to be searched can grow exponential in problem description

Idea:

- Ignore some details
- Focus on bigger picture
- Guide search for a plan



Safe Abstraction

Safe abstractions in Fast Downward



How to find safe variables? How to refine abstract plan?

Safe (Variable) Abstraction

How to find safe variables?

Free Domain Transition Graph (Free DTG)

truck(location of truck)

- *at-A*
- *at-B*

drive-truck-a-b pre: *truck* → *at-A* eff: *truck* → *at-B drive-truck-b-a* pre: *truck* → *at-A* eff: *truck* → *at-B*

at-A Externally Required Externally Required Externally Caused

A variable can be abstracted safely if, in the free DTG:

- All ex. required values are strongly connected. ✓
- Every ex. required value can be reached from any ex. caused value. ✓
- The goal value (if present) can be reached from each ex. required value. \checkmark

drop-truck-b-p pre: *truck* → *at-B*, *p* → *in-truck*, *cargo* → *contains-p* eff: *p* → *at-B*, *cargo* → *empty drop-truck-a-p*

pick-up-truck-a-p
pick-up-truck-b-p

Safe (Variable) Abstraction

How to remove safe variables?

drive-truck-a-b pre: *truck* → *at-A* eff: *truck* → *at-B*

pick-up-truck-a-p

pre: *truck* \mapsto *at-A*, $p \mapsto$ *at-A*, *cargo* \mapsto *empty* eff: $p \mapsto$ *in-truck*, *cargo* \mapsto *contains-p*

Remove variable from goal condition (if present)

drive-truck-a-b pre: eff:

 $\begin{array}{l} pick-up-truck-a-p\\ pre: p\mapsto at-A, \, cargo\mapsto empty\\ eff: p\mapsto in-truck, \, cargo\mapsto contains-p \end{array}$





Safe (Variable) Abstraction

How to refine abstract plan?

```
\begin{array}{l} pick-up-truck-a-p\\ pre: \textit{truck} \mapsto \textit{at-A}, p \mapsto at-A, cargo \mapsto empty\\ eff: p \mapsto in-truck, cargo \mapsto contains-p\\ drop-truck-b-p\\ pre: \textit{truck} \mapsto \textit{at-B}, p \mapsto in-truck, cargo \mapsto contains-p\\ eff: p \mapsto at-B, cargo \mapsto empty \end{array}
```







Operator Composition

Can we abstract more variables?

Remaining operators: *drop-truck-a-p drop-truck-b-p pick-up-truck-a-p pick-up-truck-b-p*

All change variables *cargo* and *p*

Idea: Combine *pick-up* and *drop* operators

pkg1 a *Pick-up* operator always followed by *drop* operator

pick-up-truck-a-p pre: *..., cargo* → *empty* eff: *..., cargo* → *contains-p drop-truck-b-p* pre: *..., cargo* → *contains-p* eff: *..., cargo* → *empty*

Before and after sequence: $cargo \mapsto empty$



Operator Composition

Can we abstract more variables?

Remaining operators: *pick-up-a-drop-a pick-up-a-drop-b pick-up-b-drop-a pick-up-b-drop-b*

All change variables only change *p*

 $\begin{array}{l} pick-up-a-drop-b\\ pre: p\mapsto a, \, cargo\mapsto empty\\ eff: p\mapsto b\end{array}$



A variable can be abstracted safely if, in the free DTG:

- All ex. required values are strongly connected. ✓
- Every ex. required value can be reached from any ex. caused value. ✓
- The goal value (if present) can be reached from each ex. required value. \checkmark



Evaluation

Effect of safe abstraction on search

Abstraction Results



	NONE	ABSTRACTION	ALL	ALL_SOFT
Atoms Abstracted	0.0%	6.58%	6.58%	6.58%
Abstraction Steps	0	570	570	570
# Abstracted Variables	0	3141	3141	3141
# Composite Operators	0	0	0	0

	Haslum $[7]$	Our results
gripper(20)	100%	1 - 8%
logistics(28+24)	100%	100%
movie(30)	100%	0%
mystery(27)	0%	0%
mprime(28)	0%	0%
$\operatorname{grid}(5)$	$\sim 50\%$	0%
freecell(80)	0%	0%
depot(21)	110%	112%
driverlog(19)	$0\!\!-\!\!25\%$	$0\!\!-\!\!35\%$
$\operatorname{rovers}(35)$	60–90%	$37 ext{-}78\%$
$\operatorname{satellite}(27)$	100%	3283%
$\operatorname{airport}(27)$	40 - 60%	0%

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Abstraction Results

	ABSTRACTION	ALL	ALL_SOFT
abstraction time	$\begin{array}{c} 0.05s \\ 0s \\ 0.05s \end{array}$	0.05s	0.05s
composition time		38.73s	240.52s
combined time		38.78s	240.57s

Search Results

Search Algorithms					
Lama-First	Changes behaviour		NONE	ABSTRACTION	ALL
FF Blind	of search	Blind	$0.68 \mathrm{s}$	0.41s	0.60s
Biilid		FF	$0.04 \mathrm{s}$	$0.04\mathrm{s}$	0.11s
		Lama-First	0.04s	0.04s	0.11s
			NONE	ABSTRACTION	ALL
		Blind	482	560	557
		FF	1219	1306	1255
		Lama-First	1624	1648	1553

Search Results



Conclusion

What have we learned? What is left to do?

Conclusion

- Unable to reproduce operator composition
- Safe abstraction can improve time and memory usage
- Can have counter-intuitive influence on sophisticated search algorithm
- Effectiveness of safe abstraction greatly depends on problem encoding

Future?

• Why are our results different from Haslum?

- Is the problem the encodings?
- Is the problem the implementation of composition?
- What do the other techniques Haslum mentioned do?
- Why do some domains (rovers, satellite) behave counter to our intuition in Lama-First?

Addendum_1 – Plan Length



Addendum_2 – Composition Time

composition-time



	ABSTRACTION	ALL	ALL_SOFT
abstraction time	$0.05 \mathrm{s}$	$0.05 \mathrm{s}$	$0.05 \mathrm{s}$
composition time	0s	$38.73\mathrm{s}$	240.52s
combined time	0.05s	38.78s	$240.57 \mathrm{s}$

Addendum_3 – Memory

	NONE	ABSTRACTION	ALL
Blind	$75.81 \mathrm{MB}$	60.01 MB	60.11 MB
\mathbf{FF}	$\mathbf{22.64~MB}$	$23.47 \mathrm{~MB}$	$23.49 \mathrm{MB}$
Lama-First	$\mathbf{22.56~MB}$	$23.18 \mathrm{MB}$	$23.19 \mathrm{MB}$

Addendum_4 – Free DTG example

Free DTG of *p*

Free Domain Transition Graph (Free DTG)

p (location of package)

- *at-A*
- *at-B*
- *in-truck*

drop-truck-a-p

pre: $truck \mapsto at$ -A, $p \mapsto in$ -truck, $cargo \mapsto contains$ -p eff: $p \mapsto at$ -A, $cargo \mapsto empty$

drop-truck-b-p

pick-up-truck-a-p pick-up-truck-b-p

A variable can be abstracted safely if, in the free DTG:

- All ex. required values are strongly connected. ×
- Every ex. required value can be reached from any ex. caused value. x
- The goal value (if present) can be reached from each ex. required value. x



