

Depth-Bound Heuristics and Iterative-Deepening Search Algorithms in Classical Planning

Bachelor's Thesis Presentation
Florian Spiess, 13 June 2017

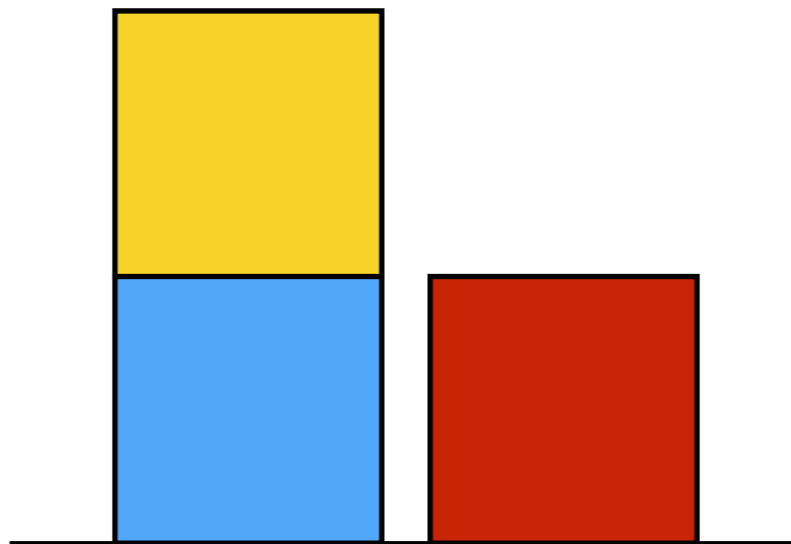
Departement of Mathematics and Computer Science
Artificial Intelligence

Classical Planning

- Goal: Find series of actions from initial to goal state
- Static, deterministic, fully observable, discrete, single-agent problems
- E.g.:
 - Shortest package delivery route
 - Stacking blocks

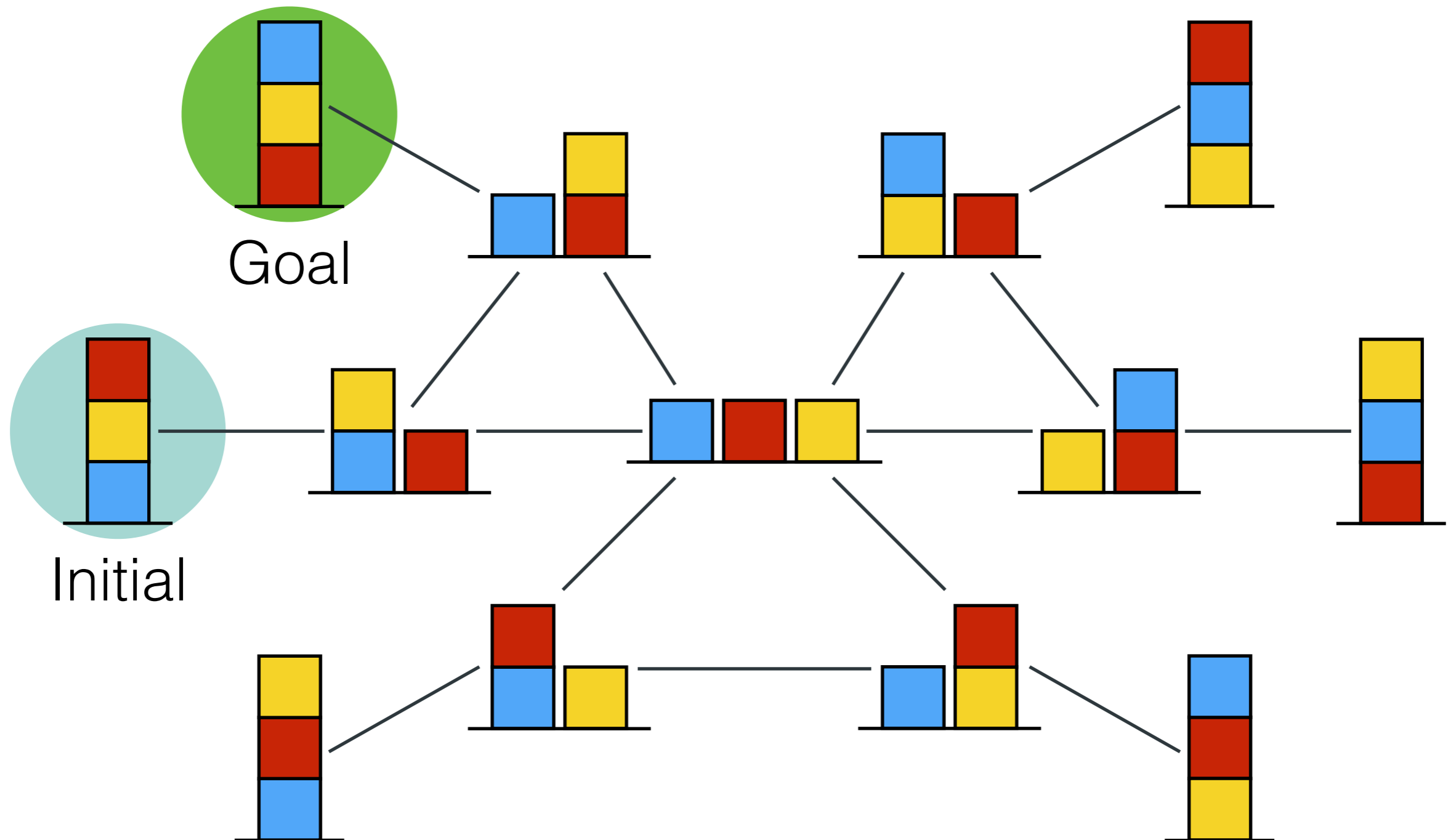
Blocks World

- Goal: Stack blocks in a certain order
- Only move one block at a time
- Only move blocks at the top of stacks



Blocks World

State Space



Heuristics

- Approximate goal distance
- Require time to construct / calculate

Goal

- Depth-bound heuristics
- Evaluate with iterative-deepening search algorithms
- Implementation in Fast Downward

Merge-and-Shrink

- Constructs abstract state space
- Calculates heuristic value in abstract state space

Merge-and-Shrink

State Space Representation

- States can be represented as lists of variables
- E.g. Logistics with one package, two trucks:
 - Package \rightarrow Left
 - Truck A \rightarrow Right
 - Truck B \rightarrow Left

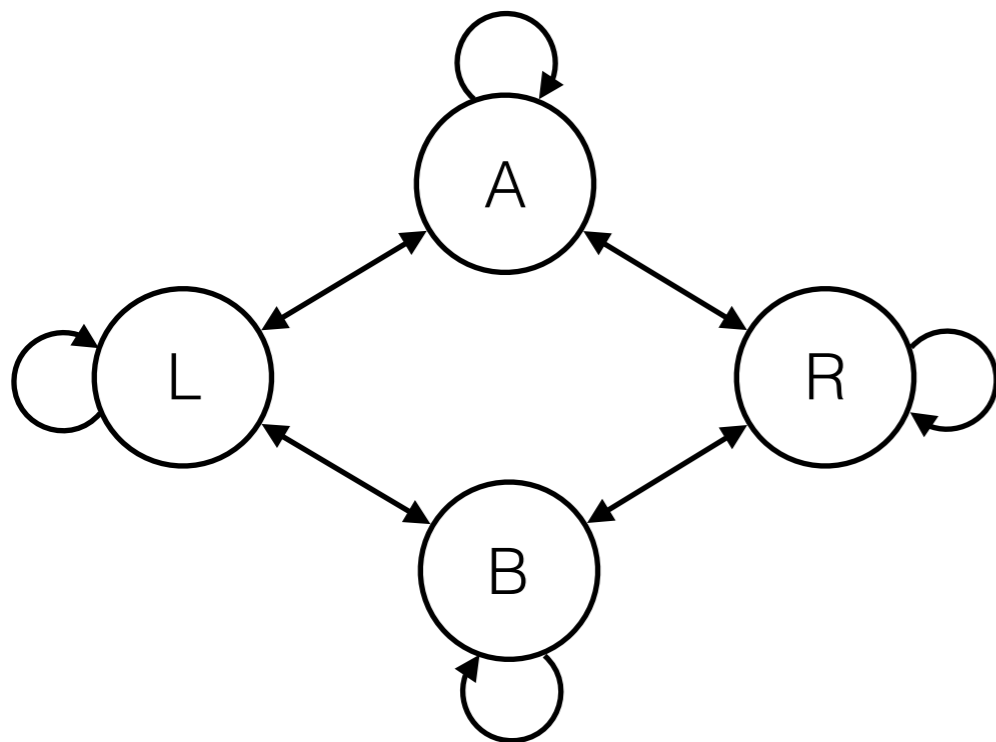


Merge-and-Shrink

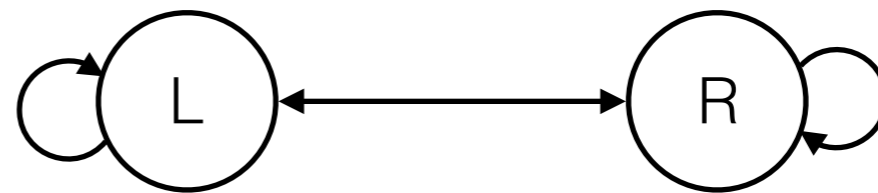
Projection

- Only considers state change of one variable
- E.g. projection onto:

Package



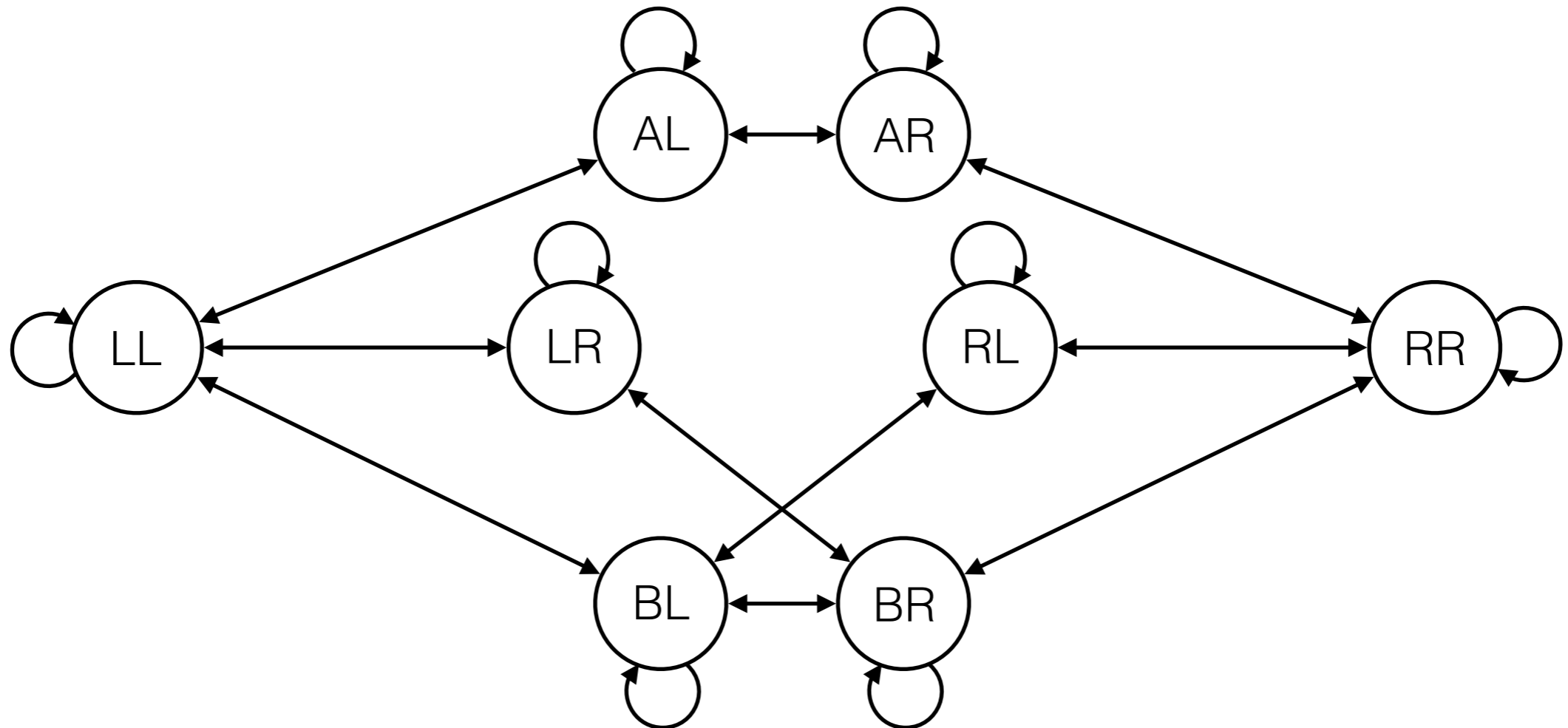
Truck A



Merge-and-Shrink

Merge

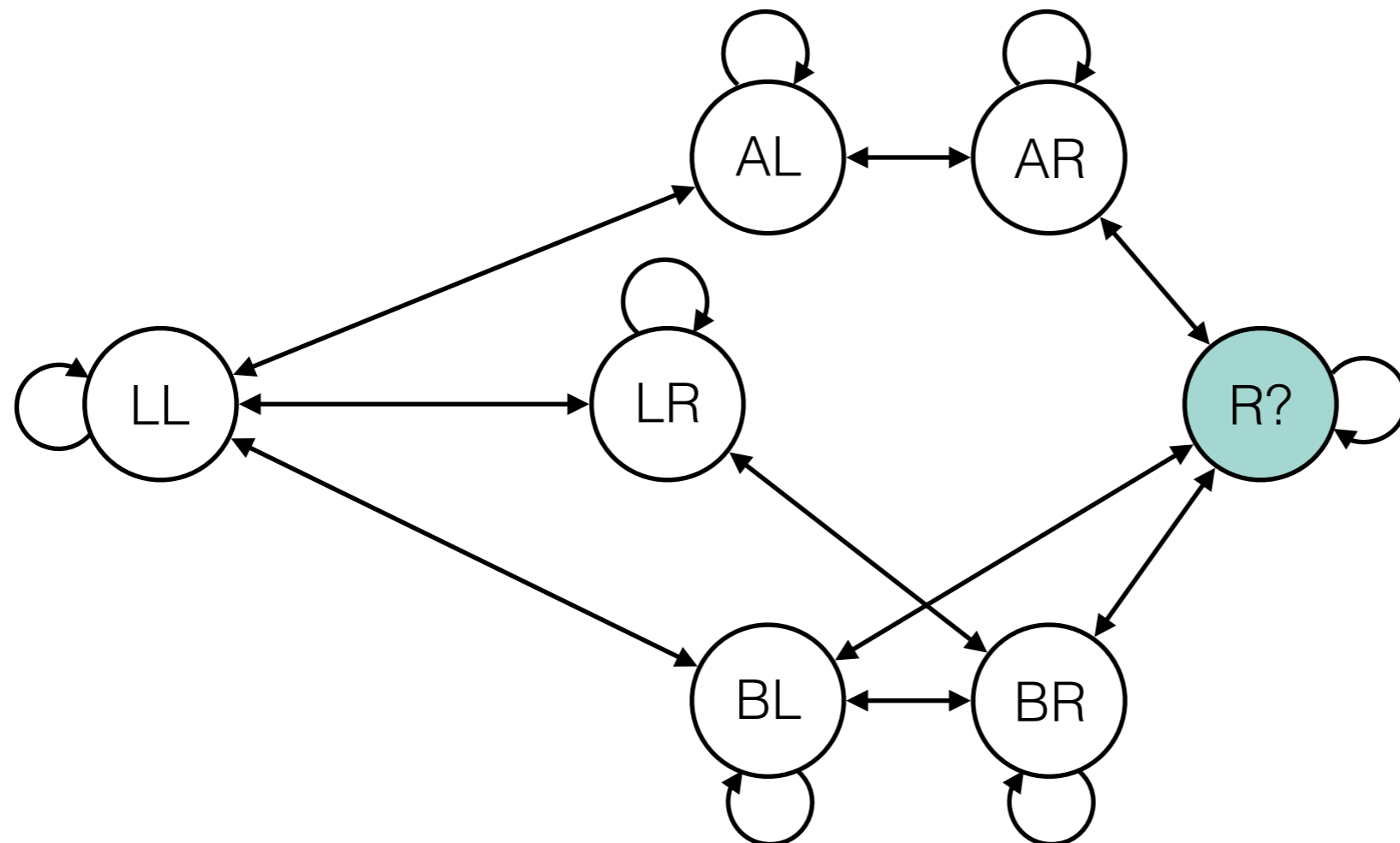
- Merge through synchronized product
- E.g. merge of projections on Package and Truck A:



Merge-and-Shrink

Shrink

- Combine states to reduce size



Merge-and-Shrink

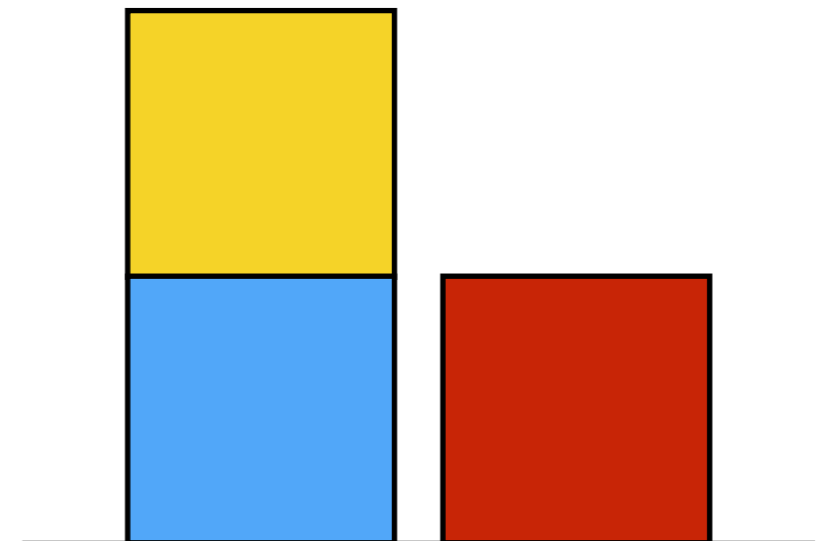
Modification

- Prune abstract states with cost $> f$ -bound
 - Reduce construction time
 - Increase heuristic accuracy

Landmark Cut

State Space Representation

- States can be represented as set of propositions
- E.g. Blocks world:
 - state = {Y-on-B, B-on-F, R-on-F}



Landmark Cut

Delete Free Planning Task

- Acquired proposition cannot be lost

- E.g.:
{Y-on-F, B-on-F, R-on-F}

— move Yellow onto Blue —>

{Y-on-F, B-on-F, R-on-F, Y-on-B}

Landmark Cut

- Estimates the minimum cost of a delete free plan
- Iteratively sums costs of required actions

Landmark Cut

Modification

- Stop calculation once sum of costs $>$ f -bound
 - Reduce calculation time

IDA* Search

- Iterative-deepening A*
- Tree search
- Explores paths until $f > f\text{-bound}$
- Restarts with increased $f\text{-bound}$
- No open list
- No closed list \rightarrow low memory usage

IDA* Search

Implementation

- Successor generation requires closed list in Fast Downward
- With closed list
- With duplicate detection

IDBFA* Search

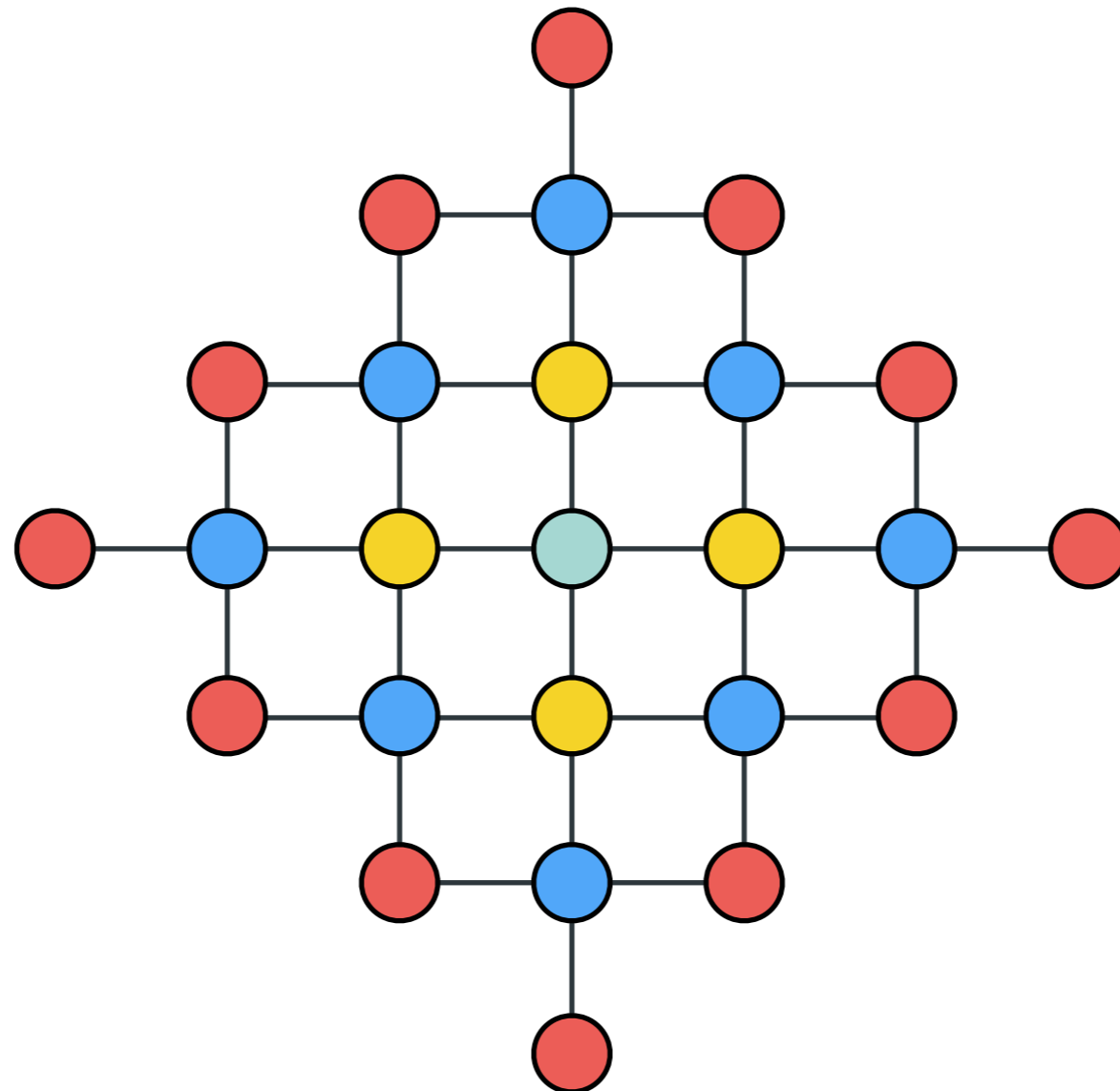
- Iterative-deepening breadth-first A*
- A* search but prunes nodes with $f > f\text{-bound}$
- No solution \longrightarrow increase $f\text{-bound}$

Breadth-First Heuristic Search

- Store explored nodes —> High space complexity
- Only search frontier required to find goal
 - > Delete visited nodes
 - No duplicate detection!
 - No solution path!

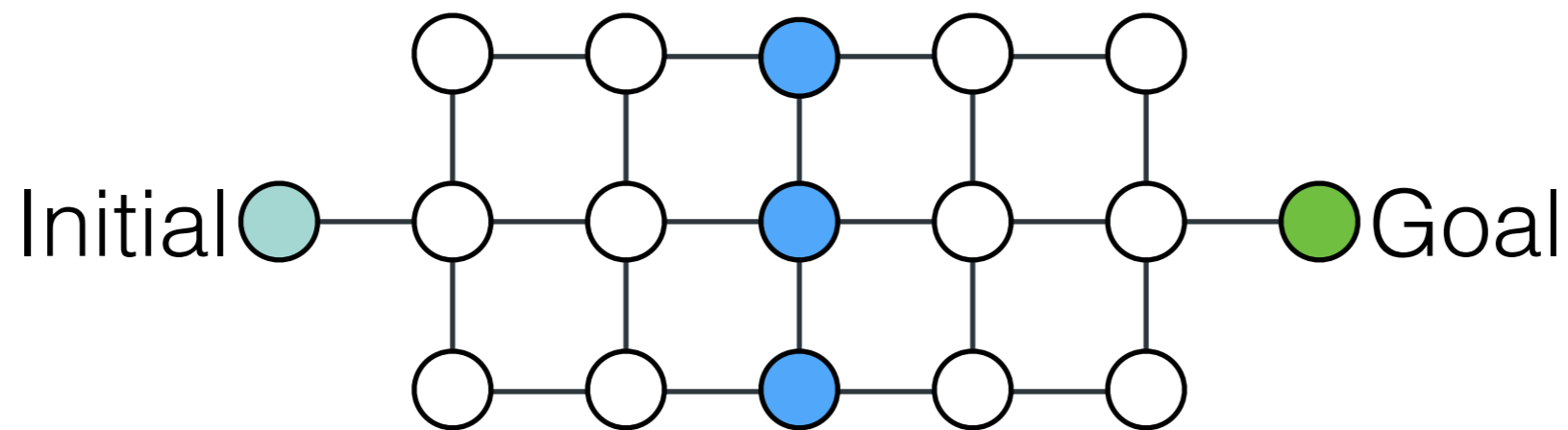
Breadth-First Heuristic Search

- Breadth-first search explores nodes in 'depth-layers'



Breadth-First Heuristic Search

- Save one intermediate layer
- Recursively solve problems



Evaluation

- Experiments on 1667 Tasks
(from 57 domains)
- IDBFHS on subset of 160 Tasks
(from 6 unit-cost, undirected graph domains)

Results

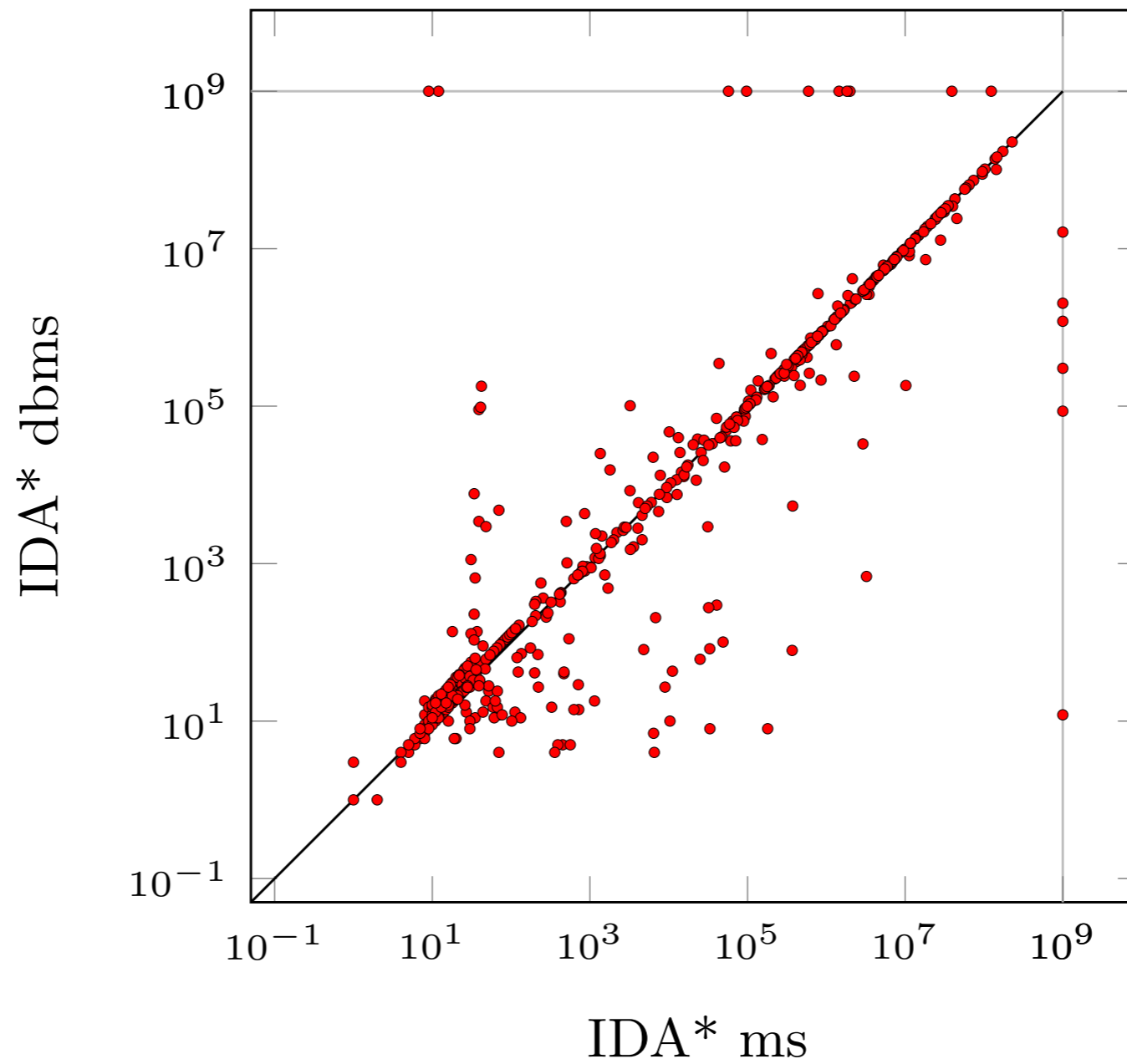
IDA* Comparison

	Standard	Merge-and-Shrink Depth-bound	Difference	Standard	Landmark Cut Depth-bound	Difference
Coverage	725	721	-4	848	833	-15
Expansions	4252.10	2790.90	-1461.2	3259.94	3286.78	26.84
Memory	62302616	61688396	-614220	12920584	12326636	-593948
Real search time	0.05	0.03	-0.01	0.68	0.72	0.04
Search time	0.24	4.62	4.38	1.20	1.37	0.17
Total time	2.79	4.69	1.9	1.30	1.49	0.19

- Depth-bound heuristics have lower coverage
- Depth-bound heuristics are slower
- Depth-bound M&S requires fewer expansions

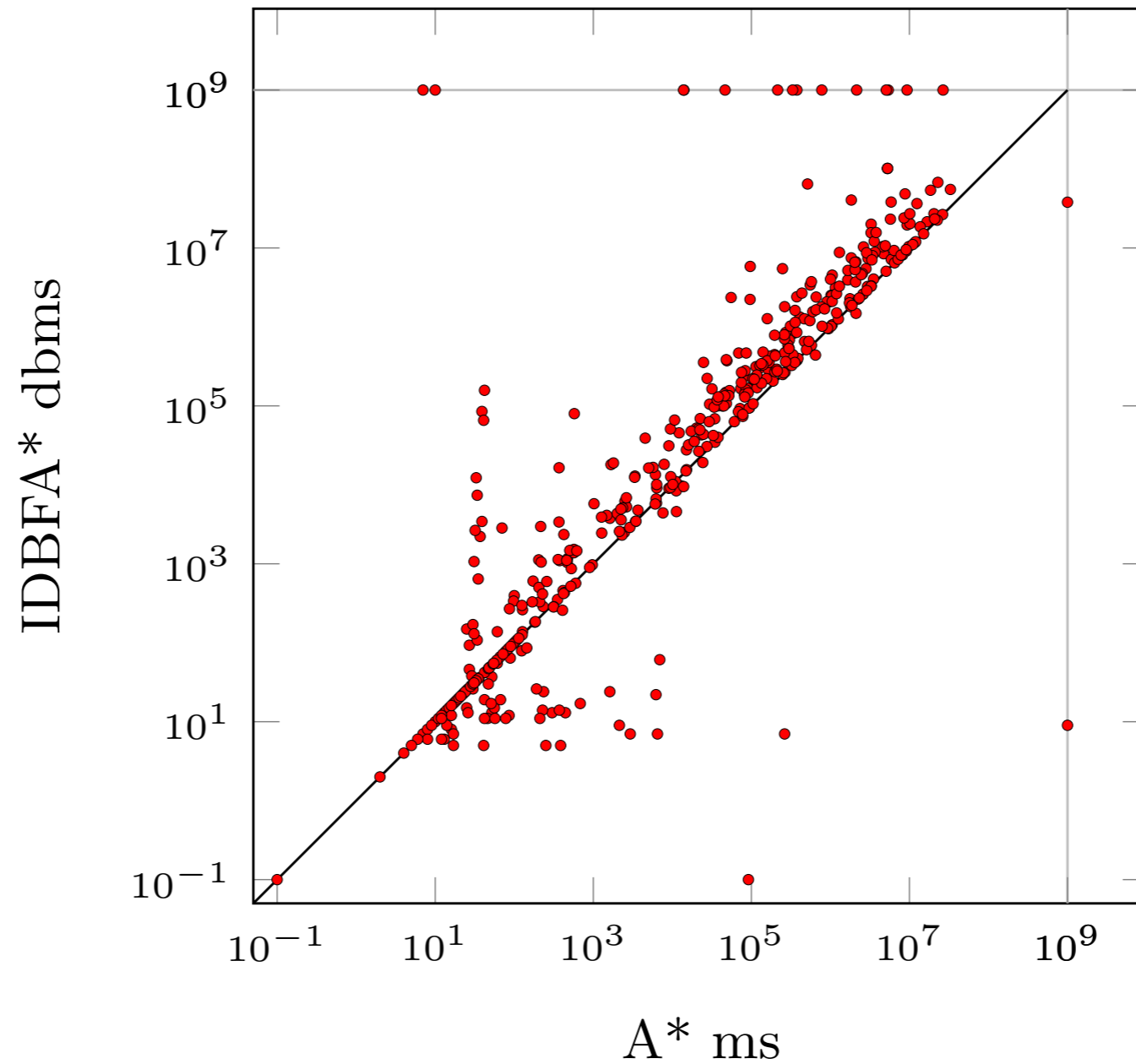
Results

Expansions



Results

Expansions



Results

A* and IDBFHS

	A*		IDBFHS	
	Merge-and-Shrink	Landmark Cut	Merge-and-Shrink	Landmark Cut
Coverage	88	82	75	80
Expansions	2184.86	2020.73	23929.42	11388.37
Memory	6320500	1032548	9927308	1518924
Search time	0.14	0.50	4.39	1.79
Total time	1.30	0.52	4.43	1.81

- IDBFHS completed fewer tasks than A*
- IDBFHS had higher peak memory

Conclusion

- Depth-bound LM-cut not enough time gain
- Depth-bound M&S slower because of construction
- Depth-bound M&S more accurate for easy tasks

Future Work

- Algorithm determines task complexity:
 - Simple: use depth-bound M&S
 - Complex: use unbound M&S
- Increase M&S depth-bound in greater steps

Thank you for your attention!

Results

Summary

	A*		IDA*				IDBFA*	
	ms	lmcut	ms	lmcut	dbms	dblmcut	dbms	dblmcut
Coverage	745	882	725	848	721	833	728	840
Expansions	1822.21	1301.20	3939.90	3088.52	2587.65	3113.72	2389.86	3079.64
Memory	63368336	21006000	53595072	9802372	52926128	9409960	60730232	20403740
Search time	0.13	0.60	0.22	1.12	4.46	1.28	4.76	1.33
Total time	2.01	0.65	2.68	1.22	4.53	1.40	5.07	1.45

Results

Total Time

