

# A non linear Kohonen algorithm

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## Introduction and background material

It is well known (see [5],[6]) that, whether the stimuli distribution  $\mu$  on  $\mathbb{R}^d$  is continuous<sup>1</sup> or not, the  $d$ -dim Kohonen algorithm with 0 neighbour<sup>2</sup> and  $n$  points (or *units*) is a stochastic gradient that derives from a potential defined

$$\text{by } \forall x = (x_1, \dots, x_n) \in (\mathbb{R}^d)^n \quad E_n^{2,\mu}(x) := \int \min_{1 \leq i \leq n} \|x_i - \omega\|^2 \mu(d\omega) \quad (1)$$

provided that  $\mu$  has a compact support. In Neural Networks terminology  $x_i$  is for the weight of *unit*  $i$  and  $\omega$  is a generic stimulus. For obvious reasons, this potential may be considered as a quadratic one. In [6] was introduced a non quadratic generalization of this potential. For every  $\alpha > 0$ , the potential  $E_n^{\alpha,\mu}$  was simply defined on the space  $(\mathbb{R}^d)^n$  by

$$\forall x = (x_1, \dots, x_n) \in (\mathbb{R}^d)^n \quad E_n^{\alpha,\mu}(x_1, \dots, x_n) := \int \min_{1 \leq i \leq n} \|x_i - \omega\|^\alpha \mu(d\omega). \quad (2)$$

The Kohonen potential corresponds to  $\alpha = 2$ . Actually, if  $\mu$  is continuous,

$$E_n^{\alpha,\mu}(x) = \sum_{i=1}^n \int_{C_i(x)} \|x_i - \omega\|^\alpha \mu(d\omega) \quad (3)$$

where  $(C_i(x))_{1 \leq i \leq n}$  denotes the so-called Euclidean Voronoï tessellation of the space  $\mathbb{R}^d$  related to  $x$ . In fact, this tessellation is only defined on  $D_n := \{x \in (\mathbb{R}^d)^n / x_i \neq x_j \iff i \neq j\}$ , by

$$C_i(x) := \{u \in \mathbb{R}^d / \|x_i - u\| < \|x_k - u\|\}, \text{ if } k \neq i, 1 \leq i \leq n.$$

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<sup>1</sup>A measure is continuous iff no hyperplan is weighed by  $\mu$ . All the distributions that have a density are continuous.

<sup>2</sup>also known as a "space quantization algorithm".

Following decomposition (3),  $E_n^{\alpha,\mu}(x)$  was called the “ $(\alpha, \mu)$ -magnitude of the Voronoi tessellation at  $x = (x_1, \dots, x_n) \in (\mathbb{R}^d)^n$ ”.

When  $\alpha \in (0, 2]$ , the  $(\alpha, \mu)$ -magnitude function  $E_n^{\alpha,\mu}$  was originally introduced in [6] as an upper-bounding modulus in a new method for high dimensional numerical integration of  $\alpha$ -hölder<sup>3</sup> functions or functions having a  $(\alpha-1)$ -hölder first derivative. This new numerical integration technique is based on the Voronoi tessellation of any  $n$ -uplet  $x^*$  that minimizes  $E_n^{\alpha,\mu}$  (see [6]).

The proposed numerical method for the computation of  $x^*$  is of course the gradient descent algorithm related to the potential  $E_n^{\alpha,\mu}$

$$X^{t+1} = X^t - \varepsilon_{t+1} H_n^\alpha(X^t, \omega^{t+1}), \quad X^0 \in D_n \quad (4)$$

where, for every  $x \in D_n$  and  $\omega \in C$ ,  $H_n^\alpha(x, \omega) := \left( \frac{x_i - \omega}{\|x_i - \omega\|} \|x_i - \omega\|^{\alpha-1} \mathbf{1}_{C_i(x)}(\omega) \right)_{1 \leq i \leq n}$ ,  $(\varepsilon_t)_{t \geq 1}$  is  $]0, 1[$ -valued sequence of steps and  $\omega^t$  is an i.i.d. sequence of random variables with distribution  $\mu$ . This formula straightforwardly derives from the integral representation on  $D_n$  of  $\nabla E_n^{\alpha,\mu}(x) = \alpha \left( \int_{C_i(x)} \|x_i - \omega\|^{\alpha-1} \frac{x_i - \omega}{\|x_i - \omega\|} \mu(d\omega) \right)_{1 \leq i \leq n}$ .

## 1 Design of the non linear Kohonen algorithm

Let us consider now a general unit set  $I \subset \mathbf{Z}^d$  endowed with a topological structure provided by a neighbourhood function  $\sigma$  defined on  $I \times I$ . In most practical cases  $\sigma(i, j) := v(i - j)$  with  $v(-x) = v(x)$  and we will often denote  $\sigma(i - j)$  instead of  $\sigma(i, j)$ . Then the algorithm displays as

(i) Computation of the winning unit  $i^{t+1} := i(\omega^{t+1}, X^t) = \operatorname{argmin}_k \|\omega^{t+1} - X_k^t\|$ . In case of conflict, one takes the lexicographic minimum,

$$(ii) \forall j \in I, X_j^{t+1} = X_j^t - \varepsilon_{t+1} \sigma(i^{t+1} - j) \frac{(X_j^t - \omega^{t+1})}{\|X_j^t - \omega^{t+1}\|} \|X_j^t - \omega^{t+1}\|^{\alpha-1}, \quad (5)$$

where  $(\varepsilon_t)_{t \geq 1}$  is still a sequence of  $(0, 1)$ -valued real numbers.

Where does this extension come from? Assume for some time that  $\mu$  has a discrete support, namely  $\operatorname{supp}(\mu) := \{\omega_1, \dots, \omega_p, \dots\}$  and set

$$\forall x \in (\mathbb{R}^d)^n \quad E_I^{\alpha,\mu,\sigma}(x) := \sum_{i,j \in I} \sigma(i - j) \int_{C_j(x)} \|x_i - \omega\|^\alpha \mu(d\omega).$$

Then, following [7] or [8], one checks that, whenever no  $\omega_p$  lies in the borders  $\partial C_i(x)$  of the tessels  $C_i(x)$  i.e.  $\operatorname{supp}(\mu) \cap (\cup_{i \in I} \partial C_i(x)) = \emptyset$ ,  $E_I^{\alpha,\mu,\sigma}$  is differentiable at  $x = (x_i)_{i \in I}$  with a gradient  $\nabla E_I^{\alpha,\mu,\sigma}$  given by

<sup>3</sup>  $f : E \mapsto F$  is  $\alpha$ -hölder iff  $\forall x, y \in E, \|f(x) - f(y)\|_F < C \|x - y\|_E^\alpha$  where  $E$  and  $F$  are normed vector space.

$$\nabla E_I^{\alpha, \mu, \sigma}(x) = \left( \sum_{j \in I} \sigma(i-j) \int_{C_j(x)} \frac{x_i - \omega}{\|x_i - \omega\|} \|x_i - \omega\|^{\alpha-1} \mu(d\omega) \right)_{i \in I} \quad (6)$$

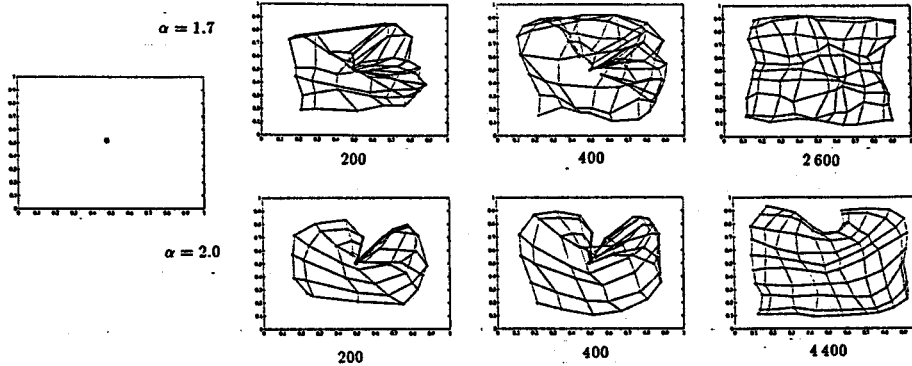
*A contrario* when  $\mu$  has a density  $f$  i.e.  $\mu = f(\omega)d\omega$ ,  $\frac{\partial}{\partial x_k} \mu(C_i(x))$  is generally not 0 whenever  $\overline{C_i(x)} \cap \overline{C_j(x)} \neq \emptyset$ . Actually, following [4], it reads

$$\frac{\partial}{\partial x_k} \mu(C_i(x)) = \int_{\overline{C_i(x)} \cap \overline{C_k(x)}} f(\omega) \left( \frac{1}{2} \bar{n}_x^{ik} + \frac{1}{\|x_i - x_k\|} \left( \frac{x_i + x_k}{2} - \omega \right) \right) \lambda_x^{ik}(d\omega)$$

where  $\bar{n}_x^{ik} := \frac{x_k - x_i}{\|x_k - x_i\|}$  and  $\lambda_x^{ik}$  denotes the  $(d-1)$ -dimensional Lebesgue measure on  $\overline{C_i(x)} \cap \overline{C_k(x)}$  (when not reduced to a  $(d-2)$ -dimensional affine subspace). So, in this case  $\nabla E_I^{\alpha, \mu, \sigma}$  is still differentiable at any point of  $D_n$ , but has no longer any integral representation with respect to  $\mu$ . More generally, the equation (6) holds as soon as  $\mu(C_i(x+h)) - \mu(C_i(x)) = o(h)$  as  $h \rightarrow 0$ .

## 2 Application to the accelerated self-organization

Rather unexpectedly, this generalization of the algorithm turned out to have some interesting self-organizing feature. Actually, this observation is the main motivation for writing this contribution. Many 2-dimensional simulations implemented with various distributions showed that self-organization is carried much faster with  $\alpha$  less than 2, at least when  $\varepsilon$  is small. For instance with a unit set  $I := \{1, \dots, 7\} \times \{1, \dots, 7\}$ , an 8-neighbour  $\sigma$  function<sup>4</sup>,  $\mu := U([0, 1]^2)$  and  $\varepsilon := 0.1$ , we observed an obvious self-organization after 2000 trials with  $\alpha = 1.7$  instead of 4400 trials with the usual  $\alpha = 2$  parameter. Fig. 1 below shows the main self-organizing steps with these two values of  $\alpha$ .



Nevertheless, if definitely "obvious" on simulations, self-organization is not

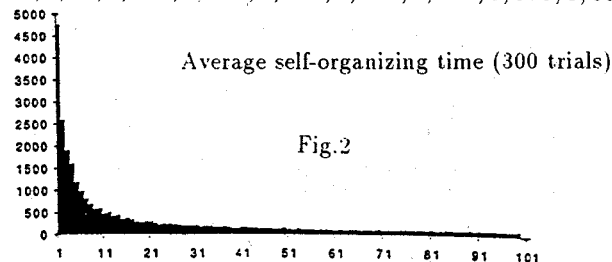
<sup>4</sup> $\sigma(i-j) := \mathbf{1}_{\{|i_1 - j_1| \leq 1\}} \mathbf{1}_{\{|i_2 - j_2| \leq 1\}}$  where  $i = (i_1, i_2)$ ,  $j = (j_1, j_2)$ .

a rigorously defined notion in multi-dimensional settings so far. So, in order to carry out a relevant study of the self-organization time as a function of the parameter  $\alpha$ , we had to restrict to the one dimensional setting where self-organization amounts to monotonicity of the  $i \mapsto x_i$ . We based our comparison on a former empirical study mentioned in [2] and recalled below in a few words.

## 2.1 Self-organization of the linear 1-dim Kohonen with two neighbours

According to the theoretical results the self-organizing time of the 2 neighbour 1-dimensional Kohonen algorithm is *a.s.* finite and actually has an exponential moment ([3], [1]) for example when the stimuli distribution is locally continuous.

The simulations already discussed in [2] emphasize that the self-organizing time is quite "reasonable" when  $\varepsilon$  is not too small. These simulations, implemented with  $n = 10$  points and  $\mu = U([0, 1])$ , were carried out on 300 independant trials of (independant) stimuli for each of the 99 selected values of the step  $\varepsilon \in \{k/100, 1 \leq k \leq 99\}$ . The same initial value was chosen at random for all the simulations. It contains 7 breaks of monotonicity :  $x := (0, 102; 0, 901; 0, 49; 0, 700; 0, 049; 0, 895; 0, 251; 0, 884; 0, 875; 0, 692)$ .



These results obviously show that the self-organization time, at least in the 1-dimensional setting, is a steep decreasing function of the step  $\varepsilon$ . If it looks natural for small values of  $\varepsilon$  ( $\varepsilon \approx 0$ ), this is much more unexpected for the great values of  $\varepsilon$  ( $\varepsilon \approx 1$ ). This phenomenon is strongly related to the number of neighbours, actually 2. Thus, if one considers a (linear) Kohonen with 4 or 6 neighbours, the self-organizing time has a minimum value and then increases again. If  $\alpha < 2$  we will see below that the same phenomenon occurs.

## 2.2 Simulations with the non linear Kohonen algorithm

The most striking feature of the above study is that, at least with 2 neighbours, the closer to 1  $\varepsilon$  is, the faster self-organization occurs. However,  $\varepsilon \approx 1$  is a totally unrealistic choice when taking into account the quantization phase of the process which requires to let  $\varepsilon$  go to 0. So, it seems quite valuable to cut down the self-organization time for some  $\varepsilon$  close to 0.

Notice that *a contrario*, whenever  $\alpha < 2$ , self-organization is no longer stable,

at least on a theoretical point of view <sup>5</sup>. However, as far as actual simulations are concerned, self-organization is quite robust for small enough  $\epsilon$ . Of course some important degradations is observed when  $\epsilon$  increases. That is why we processed our testing bench only for small  $\epsilon \in (0, 0.45]$ , namely  $\epsilon \in \{\frac{k}{100}, 1 \leq k \leq 10, 0.12, 0.14, 0.16, 0.18, 0.20, 0.25, 0.30, 0.35, 0.40, 0.45\}$ .

The non linear Kohonen algorithm was tested with 4 values of the parameter  $\alpha$ , namely  $\alpha \in \{1.6, 1.7, 1.8, 1.9\}$ , to which the linear case  $\alpha = 2$  was added in the below tables. The initial value and the number of trials - 300 - for each couple  $(\alpha, \epsilon)$  is the same as for  $\alpha = 2$ . The results are displayed in the below table (rounded up to the nearest integer) and frequency histograms.

$\epsilon$	$\alpha$	1.6	1.7	1.8	1.9	2.0
0.01		956*	1 589	2 132	2 756	4 791
0.02		530*	743	1 178	1 727	2 735
0.03		410*	552	882	1 215	1 764
0.04		320*	420	587	872	1 349
0.05		252*	377	474	781	1 134
0.06		252*	314	477	638	983
0.07		224*	286	394	547	908
0.08		218*	280	330	486	724
0.09		190*	216	293	486	601
0.10		193*	193	267	461	523
0.12		171*	175	240	378	499
0.14		173*	175	245	312	463
0.16		177	150*	199	286	385
0.18		185	141*	170	240	349
0.20		209	142*	173	222	322
0.25		407	131*	149	211	284
0.30		719	216	134*	180	224
0.35		2 610	369	132*	148	198
0.40		7 114	660	157	143*	178
0.45		21 170	1 554	1 942	126*	174

The \* denotes the minimum self-organizing time observed for a given  $\epsilon$ .

### 2.3 Conclusions

In our opinion, the following remarks are quite interesting for future simulations and tests:

- If  $\epsilon < 0.15$  then the self-organizing time is always increasing with  $\alpha$ .
- If  $\epsilon > 0.15$  then the self-organizing time goes through a minimum as a function of  $\alpha$ . For example, if  $\epsilon = 0.16$ , it is achieved at  $\alpha = 1.7$  ; for  $\epsilon = 0.30$ , it is achieved at  $\alpha = 1.8$ . We think that, for every  $\epsilon \in (0.15, 1)$ , there is an optimal value  $\alpha_{\min}(\epsilon)$  of the parameter  $\alpha$  that probably grows up with  $\epsilon$ .

<sup>5</sup>It means that  $i \mapsto X_i^t$  may loose its monotonicity after self-organization occurred which is impossible if  $\alpha = 2$ .

- Whatsoever, the most interesting feature of these simulations is that when  $\varepsilon$  is close to 0, the self-organization goes faster with smaller  $\alpha$ . Thus, if  $\varepsilon = 0.01$ , the self-organization time is more than 5 times shorter with  $\alpha = 1.6$  than with  $\alpha = 2$ . Consequently, using  $\alpha = 1.6$  and small values of  $\varepsilon$  will both achieve an accelerated self-organization and a good space quantization.

## Conclusion

In this paper, we give some evidences about the superiority of this “non linear” Kohonen algorithm to achieve both self-organization and space quantization, provided that one works with some small constant gain parameter  $\varepsilon$ . Although these first empirical results are quite promising, some further investigations must be carried out in two directions:

- A wide testing bench to confirm the efficiency of the method and to develop some “know how” concerning the optimal parameters  $(\alpha, \varepsilon)$ .
- Some theoretical study in order to provide some analytic knowledge of the function  $\varepsilon \mapsto \alpha_{\min}(\varepsilon)$ .

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