

Aggregation Operators on Bounded Partially Ordered Sets, Aggregative Spaces and Their Duality

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Abstract. The present paper introduces aggregative spaces and their category **AGS**, and then establishes a dual adjunction between **AGS** and the category **Agop** of aggregation operators on bounded partially ordered sets. Spatial aggregation operators and sober aggregative spaces, enabling us to restrict the dual adjunction between **AGS** and **Agop** to a dual equivalence between the full subcategory of **Agop** consisting of spatial aggregation operators and the full subcategory of **AGS** consisting of sober aggregative spaces, will also be subjects of this paper.

Keywords: Aggregation operator, Aggregative space, Q-space, Generalized topological space, Category theory, Adjoint situation, Duality, Spatiality, Sobriety.

1 Introduction

There is a considerable interest in the studies on (n -ary) aggregation operators (agops for short) for replacing the particular bounded partially ordered set (poset for short) $([0, 1], \leq)$ by other reasonable bounded posets, e.g. agops on $([a, b], \leq)$ in [3,10,14], agops on $(I[0, 1], \preceq_w)$ (the so-called an interval-valued agops) in [12], agops on (L^*, \leq_{L^*}) in [8], triangular norms on a general bounded poset in [4], a general bounded lattice in [15,16], pseudo-uniforms on a general complete lattice in [17,18]. As is shown in [5,11,13], agops on general bounded posets and their category **Agop** provide a useful and an abstract framework for such studies.

The dualities between certain ordered algebraic structures and certain spaces have been an important issue in many branches of mathematics (see [6,7] and the references therein). The famous duality between the full subcategory **SobTop** of **Top** of all sober topological spaces and the full subcategory **SpatFrm** of **Frm** of all spatial frames [9] is one of such dualities. In an analogous manner to this duality, our aim in this paper is to find out an appropriate notion of space providing a categorical duality for agops. For this purpose, after the next preliminary section, we introduce the aggregative spaces and their category **AGS**, and establish a dual adjunction between **AGS** and **Agop** in Section 3. Section 4 provides the notions of spatial agops and sober aggregative spaces, and proves

a dual equivalence between the full subcategory **SobAGS** of **AGS** of all sober aggregative spaces and the full subcategory **SpatAgop** of **Agop** of all spatial agops. Furthermore, the presented dual adjunction and dual equivalence have been also discussed for some full subcategories of **Agop**.

2 Preliminaries

2.1 Categorical Tools

Adjoint situations and equivalences in the category theory are essential tools for formulating the main results of this paper. By definition, an adjoint situation $(\varrho, \phi) : F \dashv G : \mathbf{C} \rightarrow \mathbf{D}$ consists of functors $G : \mathbf{C} \rightarrow \mathbf{D}$, $F : \mathbf{D} \rightarrow \mathbf{C}$, and natural transformations $id_{\mathbf{D}} \xrightarrow{\varrho} GF$ (called the unit) and $FG \xrightarrow{\phi} id_{\mathbf{C}}$ (called the co-unit) satisfying the adjunction identities $G(\phi_A) \circ \varrho_{G(A)} = id_{G(A)}$ and $\phi_{F(B)} \circ F(\varrho_B) = id_{F(B)}$ for all A in \mathbf{C} and B in \mathbf{D} . If $(\varrho, \phi) : F \dashv G : \mathbf{C} \rightarrow \mathbf{D}$ is an adjoint situation for some ϱ and ϕ , then F is said to be a left adjoint to G , $F \dashv G$ in symbols. A functor $G : \mathbf{C} \rightarrow \mathbf{D}$ is called an equivalence if it is full, faithful and isomorphism-dense. In this case, \mathbf{C} and \mathbf{D} are called equivalent categories, denoted by $\mathbf{C} \sim \mathbf{D}$.

Proposition 1. [2,7] *Given an adjoint situation $(\varrho, \phi) : F \dashv G : \mathbf{C}^{op} \rightarrow \mathbf{D}$, let $Fix(\phi)$ denote the full subcategory of \mathbf{C} of all \mathbf{C} -objects A such that $\phi_A^{op} : A \rightarrow FGA$ is a \mathbf{C} -isomorphism, and $Fix(\varrho)$ the full subcategory of \mathbf{D} of all \mathbf{D} -objects B such that $\varrho_B : B \rightarrow GFB$ is a \mathbf{D} -isomorphism. Then the following statements are true:*

- (i) *The restriction of $F \dashv G$ to $[Fix(\phi)]^{op}$ and $Fix(\varrho)$ induces an equivalence $[Fix(\phi)]^{op} \sim Fix(\varrho)$.*
- (ii) *If ϕ_A^{op} is an epimorphism in \mathbf{C} for each \mathbf{C} -object A , then both $Fix(\phi)$ and $Fix(\varrho)$ are reflective in their respective categories with the reflectors $F^{op}G^{op}$ and GF , and reflection arrows ϕ_A^{op} and ϱ_B , resp.*

For more information about adjoint situations and equivalences, we refer to [1].

2.2 Aggregation Operators and Their Categories

Let (L, \leq) be a bounded poset with the least element \perp and the greatest element \top . An aggregation operator on L is defined to be a function $A : \bigcup_{n \in \mathbb{N}^+} L^n \rightarrow L$

satisfying the following conditions:

(AG.1) If $\alpha_1 \leq \beta_1, \alpha_2 \leq \beta_2, \dots$ and $\alpha_n \leq \beta_n$ for all $\alpha_i, \beta_i \in L, i = 1, 2, \dots, n$ ($n \in \mathbb{N}^+$), then $A(\alpha_1, \dots, \alpha_n) \leq A(\beta_1, \dots, \beta_n)$.

(AG.2) $A(\alpha) = \alpha$ for all $\alpha \in L$.

(AG.3) $A(\perp, \dots, \perp) = \perp$ and $A(\top, \dots, \top) = \top$.

For $n \geq 2$, a function $B : L^n \rightarrow L$ is called an n -ary aggregation operator on L iff the conditions (AG.1) and (AG.3) are satisfied. A 1-ary aggregation operator $B : L \rightarrow L$ is the identity map id_L on L . Every aggregation operator A on L

uniquely determines a family of n -ary aggregation operators $\{A_n \mid n \in \mathbb{N}^+\}$ by $A_n(\alpha_1, \dots, \alpha_n) = A(\alpha_1, \dots, \alpha_n)$.

With regard to the special cases of (L, \leq) , an aggregation operator on L produces an aggregation process for fuzzy sets, interval-valued fuzzy sets, intuitionistic fuzzy sets, type 2 fuzzy sets and probabilistic metrics [5]. We further remark that whereas (AG.2) is proposed as a convention by some authors (e.g., see [3,11,13]), this condition is used to set up many interesting properties of aggregation operators such as their close connection with partially ordered groupoids in [5].

Definition 1. [5] *The category **Agop** of aggregation operators has as objects all triples (L, \leq, A) , where (L, \leq) is a bounded poset and A is an aggregation operator on L , and as morphisms all $(L, \leq, A) \xrightarrow{u} (M, \leq, B)$, where $u : (L, \leq) \rightarrow (M, \leq)$ is an order-preserving function such that $u(\perp) = \perp$, $u(\top) = \top$ and the following diagram commutes for all $n \in \mathbb{N}^+$:*

$$\begin{array}{ccc} L^n & \xrightarrow{u^n} & M^n \\ A_n \downarrow & & \downarrow B_n \\ L & \xrightarrow{u} & M, \end{array}$$

*i.e. $u(A(\alpha_1, \dots, \alpha_n)) = B(u(\alpha_1), \dots, u(\alpha_n))$ for all $\alpha_1, \dots, \alpha_n \in L$. Composition and identities in **Agop** are taken from the category **Set** of sets and functions.*

Definition 2. [5] (i) **Asagop** is the full subcategory of **Agop** of all (L, \leq, A) such that A is associative, i.e.

$$A(\alpha_1, \dots, \alpha_k, \dots, \alpha_n) = A_2(A_k(\alpha_1, \dots, \alpha_k), A_{n-k}(\alpha_{k+1}, \dots, \alpha_n))$$

for all $n \geq 2$, $k = 1, \dots, n - 1$ and $\alpha_1, \dots, \alpha_n \in L$.

(ii) **Smasagop** is the full subcategory of **Asagop** of all (L, \leq, A) such that A is symmetric, i.e.

$$A(\alpha_1, \dots, \alpha_n) = A(\alpha_{\pi(1)}, \dots, \alpha_{\pi(n)})$$

for all $n \in \mathbb{N}^+$, $\alpha_1, \dots, \alpha_n \in L$ and for all permutations $\pi(1), \dots, \pi(n)$ of $\{1, \dots, n\}$.

3 Aggregative Spaces and Their Relations with Aggregation Operators

3.1 Definition of Aggregative Spaces and Their Category

For a given set X , we call a subset τ of the power set $\mathcal{P}(X)$ of X an aggregative system on X if $\emptyset \in \tau$, $X \in \tau$, and $G_1, G_2 \in \tau$ implies $G_1 \cap G_2 \in \tau$ for all $G_1, G_2 \in \mathcal{P}(X)$. By an aggregative space, we mean a pair (X, τ) of a set X and an aggregative system τ on X . To formulate the category of aggregative spaces, we need to recall that every function $f : X \rightarrow Y$ determines a function $f^\leftarrow : \mathcal{P}(Y) \rightarrow \mathcal{P}(X)$, sending each subset G of Y to the preimage of G under f .

Definition 3. *The category of aggregative spaces, denoted by **AGS**, is a category whose objects are aggregative spaces, and whose morphisms $f : (X, \tau) \rightarrow (Y, \nu)$ are functions $f : X \rightarrow Y$ having the property that for every $G \in \nu$, $f^{\leftarrow}(G) \in \tau$. Composition and identities in **AGS** are the same as those in **Set**.*

It is remarkable to mention that **AGS** is a special kind of the category $(\mathcal{Z}_1, \mathcal{Z}_2)\mathbf{S}$ of $(\mathcal{Z}_1, \mathcal{Z}_2)$ -spaces, which has been developed in [6]. More clearly, **AGS** is exactly the same as the category $(\mathcal{V}_\perp, \mathcal{F})\mathbf{S}$, where \mathcal{V}_\perp and \mathcal{F} are the subset systems, defined by $\mathcal{V}_\perp(P) = \{\emptyset\}$ and $\mathcal{F}(P) = \{G \subseteq P \mid G \text{ is finite}\}$ for every poset P in [6].

3.2 Dual Adjunction between AGS and Agop

Our aim in this section is to show that there exists an adjoint situation $(\eta, \varepsilon) : \Omega_{AG} \dashv Pt_{AG} : \mathbf{Agop}^{op} \rightarrow \mathbf{AGS}$. For this purpose, we first establish the functors $\Omega_{AG} : \mathbf{AGS} \rightarrow \mathbf{Agop}^{op}$ and $Pt_{AG} : \mathbf{Agop}^{op} \rightarrow \mathbf{AGS}$.

Every aggregative space (X, τ) induces an **Agop**-object $(\tau, \subseteq, X(\cap)_\tau)$, where $X(\cap)$ is the aggregation operator on $\mathcal{P}(X)$, defined by $X(\cap)(G_1, \dots, G_n) = G_1 \cap \dots \cap G_n$ and $X(\cap)(G) = G$ for every $G, G_1, \dots, G_n \in \mathcal{P}(X)$ ($n \geq 2$), and $X(\cap)_\tau$ is the restriction of $X(\cap)$ to the set $\bigcup_{n \in \mathbb{N}^+} \tau^n$. This means that the assignment of

the **Agop**-object $(\tau, \subseteq, X(\cap)_\tau)$ to every aggregative space (X, τ) is a function Ω_{AG} from the objects of **AGS** to the objects of **Agop**. On the other hand, for a given **AGS**-morphism $f : (X, \tau) \rightarrow (Y, \nu)$, the restriction $f_{|\nu}^{\leftarrow}$ of f^{\leftarrow} to ν is an **Agop**-morphism $f_{|\nu}^{\leftarrow} : (\nu, \subseteq, Y(\cap)_\nu) \rightarrow (\tau, \subseteq, X(\cap)_\tau)$. Thus, Ω_{AG} can be extended to a functor from **AGS** to **Agop**^{op}:

Proposition 2. *The map $\Omega_{AG} : \mathbf{AGS} \rightarrow \mathbf{Agop}^{op}$, defined by*

$$\Omega_{AG}(X, \tau) = (\tau, \subseteq, X(\cap)_\tau) \text{ and } \Omega_{AG}(f) = \left(f_{|\nu}^{\leftarrow}\right)^{op},$$

is a functor.

In the formulation of the functor $Pt_{AG} : \mathbf{Agop}^{op} \rightarrow \mathbf{AGS}$, we will use the notion of filter defined as follows.

Definition 4. *Let (L, \leq, A) be an object of **Agop**. A subset G of L is called a filter of (L, \leq, A) iff G satisfies the next conditions:*

- (F1) G is an upper set of (L, \leq) , i.e. for all $\alpha, \beta \in L$, $\alpha \in G$ and $\alpha \leq \beta$ imply $\beta \in G$,
- (F2) $\perp \notin G$,
- (F3) $\top \in G$,
- (F4) For all $\alpha_1, \dots, \alpha_n \in L$, $\alpha_1, \dots, \alpha_n \in G$ iff $A(\alpha_1, \dots, \alpha_n) \in G$.

Lemma 1. *Given an **Agop**-object (L, \leq, A) , let $\mathfrak{F}(L)$ denote the set of all filters of (L, \leq, A) . For each $a \in L$, let $\Psi_a = \{G \in \mathfrak{F}(L) \mid a \in G\}$ and $\Psi(L) = \{\Psi_a \mid a \in L\}$. Then, $Pt_{AG}(L, \leq, A) = (\mathfrak{F}(L), \Psi(L))$ is an aggregative space.*

Proof. For each $G \in \mathfrak{F}(L)$, by (F2) and (F3) in Definition 4, $\mathfrak{F}(L) = \Psi_\top \in \Psi(L)$ and $\emptyset = \Psi_\perp \in \Psi(L)$. Furthermore, we obtain from (F4) that for all $a, b \in L$, $\Psi_a \cap \Psi_b = \Psi_{A(a,b)}$. Therefore, $\Psi(L)$ is an aggregative system on $\mathfrak{F}(L)$.

Proposition 3. *The map $Pt_{AG} : \mathbf{Agop}^{op} \rightarrow \mathbf{AGS}$, defined by*

$$Pt_{AG} \left((L, \leq, A) \xrightarrow{u} (M, \leq, B) \right) = Pt_{AG}(L, \leq, A) \xrightarrow{Pt_{AG}(u)} Pt_{AG}(M, \leq, B),$$

where $[Pt_{AG}(u)](G) = (u^{op})^\leftarrow(G)$ for all $G \in \mathfrak{F}(L)$, is a functor.

Proof. Lemma 1 shows that Pt_{AG} maps the objects of \mathbf{Agop}^{op} to the objects of \mathbf{AGS} . Let $(L, \leq, A) \xrightarrow{u} (M, \leq, B)$ be an \mathbf{Agop}^{op} -morphism, i.e. $(M, \leq, B) \xrightarrow{u^{op}} (L, \leq, A)$ is an \mathbf{Agop} -morphism. For every $G \in \mathfrak{F}(L)$, since $(u^{op})^\leftarrow(G) \in \mathfrak{F}(M)$, $Pt_{AG}(u) : \mathfrak{F}(L) \rightarrow \mathfrak{F}(M)$ is a set map. In addition to this, we easily see that for every $b \in M$, $[Pt_{AG}(u)]^\leftarrow(\Psi_b) = \Psi_{u^{op}(b)}$, i.e. $[Pt_{AG}(u)]^\leftarrow(V) \in \Psi(L)$ for every $V \in \Psi(M)$. This proves that $Pt_{AG}(u) : Pt_{AG}(L, \leq, A) \rightarrow Pt_{AG}(M, \leq, B)$ is an \mathbf{AGS} -morphism. Hence, the assertion follows from the fact that Pt_{AG} preserves composition and identities.

To accomplish our task in this section, we now consider two natural transformations—the unit and co-unit of the asked adjunction—given in the next two lemmas.

Lemma 2. *For every \mathbf{AGS} -object (X, τ) , the map $\eta_{(X, \tau)} : X \rightarrow \mathfrak{F}(\tau)$, defined by $\eta_{(X, \tau)}(x) = \tau(x) = \{G \in \tau \mid x \in G\}$, is an \mathbf{AGS} -morphism $(X, \tau) \rightarrow Pt_{AG}\Omega_{AG}(X, \tau)$. Moreover, $\eta = (\eta_{(X, \tau)})_{(X, \tau) \in Ob(\mathbf{AGS})} : id_{\mathbf{AGS}} \rightarrow Pt_{AG}\Omega_{AG}$ is a natural transformation.*

Proof. It is obvious that for every $x \in X$, $\tau(x) \in \mathfrak{F}(\tau)$, and so $\eta_{(X, \tau)} : X \rightarrow \mathfrak{F}(\tau)$ is indeed a map. To see that $\eta_{(X, \tau)} : (X, \tau) \rightarrow Pt_{AG}\Omega_{AG}(X, \tau)$ is an \mathbf{AGS} -morphism, note first that

$$Pt_{AG}\Omega_{AG}(X, \tau) = (\mathfrak{F}(\tau), \Psi(\tau)).$$

Then, since $\eta_{(X, \tau)}^\leftarrow(\Psi_G) = G$ for every $G \in \tau$, we obtain that $\eta_{(X, \tau)}^\leftarrow(V) \in \tau$ for every $V \in \Psi(\tau)$, i.e. $\eta_{(X, \tau)} : (X, \tau) \rightarrow Pt_{AG}\Omega_{AG}(X, \tau)$ is an \mathbf{AGS} -morphism. The proof of the second part requires only the naturality of η which means the commutativity of the rectangle

$$\begin{array}{ccc} (X, \tau) & \xrightarrow{\eta_{(X, \tau)}} & Pt_{AG}\Omega_{AG}(X, \tau) \\ f \downarrow & & \downarrow & Pt_{AG}\Omega_{AG}(f) \\ (Y, \nu) & \xrightarrow{\eta_{(Y, \nu)}} & Pt_{AG}\Omega_{AG}(Y, \nu) \end{array}$$

for every \mathbf{AGS} -morphism $(X, \tau) \xrightarrow{f} (Y, \nu)$. Since

$$[Pt_{AG}\Omega_{AG}(f)](\eta_{(X, \tau)}(x)) = \left(f|_\nu^\leftarrow\right)^\leftarrow(\eta_{(X, \tau)}(x)),$$

the commutativity of the rectangle above follows from the following observation: For every $x \in X$ and every $H \in \nu$,

$$\begin{aligned} H \in \eta_{(Y,\nu)}(f(x)) &\Leftrightarrow H \in \nu \text{ and } f(x) \in H \Leftrightarrow f_{|\nu}^{\leftarrow}(H) \in \tau \text{ and } x \in f_{|\nu}^{\leftarrow}(H) \\ &\Leftrightarrow f_{|\nu}^{\leftarrow}(H) \in \eta_{(X,\tau)}(x) \Leftrightarrow H \in \left(f_{|\nu}^{\leftarrow}\right)^{\leftarrow}(\eta_{(X,\tau)}(x)) \\ &\Leftrightarrow H \in [Pt_{AG}\Omega_{AG}(f)](\eta_{(X,\tau)}(x)). \end{aligned}$$

Lemma 3. For every **Agop**-object (L, \leq, A) , the map $e_{(L, \leq, A)} : L \rightarrow \Psi(L)$, defined by $e_{(L, \leq, A)}(a) = \Psi_a$, is an **Agop**-morphism $(L, \leq, A) \rightarrow \Omega_{AG}Pt_{AG}(L, \leq, A)$. Moreover, $\varepsilon = \left(e_{(L, \leq, A)}^{op}\right)_{(L, \leq, A) \in Ob(\mathbf{Agop})} : \Omega_{AG}Pt_{AG} \rightarrow id_{\mathbf{Agop}^{op}}$ is a natural transformation.

Proof. Consider first that $\Omega_{AG}Pt_{AG}(L, \leq, A) = \left(\Psi(L), \subseteq, \mathfrak{F}(L) (\cap)_{\Psi(L)}\right)$. In order to prove that $e_{(L, \leq, A)} : (L, \leq, A) \rightarrow \Omega_{AG}Pt_{AG}(L, \leq, A)$ is an **Agop**-morphism, we proceed as follows. For all $a, b \in L$ with $a \leq b$, since $\Psi_a \subseteq \Psi_b$ (by (F1) in Definition 4), $e_{(L, \leq, A)}$ is order-preserving. We also obtain from (F2) and (F3) in Definition 4 that $e_{(L, \leq, A)}$ preserves \perp and \top . Furthermore, by making use of (F4) in Definition 4, we see that $\Psi_{a_1} \cap \dots \cap \Psi_{a_n} = \Psi_{A(a_1, \dots, a_n)}$ for all $a_1, \dots, a_n \in L$, and so

$$\begin{aligned} e_{(L, \leq, A)}(A(a_1, \dots, a_n)) &= \Psi_{A(a_1, \dots, a_n)} = \Psi_{a_1} \cap \dots \cap \Psi_{a_n} \\ &= e_{(L, \leq, A)}(a_1) \cap \dots \cap e_{(L, \leq, A)}(a_n) \\ &= \mathfrak{F}(L) (\cap)_{\Psi(L)}(e_{(L, \leq, A)}(a_1), \dots, e_{(L, \leq, A)}(a_n)). \end{aligned}$$

This completes the proof of the first part of the assertion. For the second part, the only property of ε that we have to verify is its naturality, i.e. the commutativity of the diagram

$$\begin{array}{ccc} \Omega_{AG}Pt_{AG}(L, \leq, A) & \xrightarrow{\varepsilon_{(L, \leq, A)}} & (L, \leq, A) \\ \Omega_{AG}Pt_{AG}(u) \downarrow & & \downarrow u \\ \Omega_{AG}Pt_{AG}(M, \leq, B) & \xrightarrow{\varepsilon_{(M, \leq, B)}} & (M, \leq, B) \end{array} \quad (1)$$

for each **Agop**^{op}-morphism $(L, \leq, A) \xrightarrow{u} (M, \leq, B)$. Since all arrows and all compositions in (1) are taken in **Agop**^{op}, it can be simplified to be a rectangle

$$\begin{array}{ccc} (M, \leq, B) & \xrightarrow{e_{(M, \leq, B)}} & \left(\Psi(M), \subseteq, \mathfrak{F}(M) (\cap)_{\Psi(M)}\right) \\ u^{op} \downarrow & & \downarrow Pt_{AG}(u)_{|\Psi(M)}^{\leftarrow} \\ (L, \leq, A) & \xrightarrow{e_{(L, \leq, A)}} & \left(\Psi(L), \subseteq, \mathfrak{F}(L) (\cap)_{\Psi(L)}\right), \end{array} \quad (2)$$

where all arrows and all compositions are performed in **Agop**. The commutativity of (2) is obtained as

$$\begin{aligned} \left[Pt_{AG}(u)_{|\Psi(M)}^{\leftarrow} \circ e_{(M, \leq, B)}\right](b) &= [Pt_{AG}(u)]^{\leftarrow}(\Psi_b) = \Psi_{u^{op}(b)} \\ &= [e_{(L, \leq, A)} \circ u^{op}](b) \end{aligned}$$

for all $b \in M$.

Theorem 1. $(\eta, \varepsilon) : \Omega_{AG} \dashv Pt_{AG} : \mathbf{Agop}^{op} \rightarrow \mathbf{AGS}$ is an adjoint situation.

Proof. It is not difficult to check the adjunction identities

$$Pt_{AG} (\varepsilon_{(L, \leq, A)}) \circ \eta_{Pt_{AG}(L, \leq, A)} = id_{Pt_{AG}(L, \leq, A)},$$

$$\varepsilon_{\Omega_{AG}(X, \tau)} \circ \Omega_{AG} (\eta_{(X, \tau)}) = id_{\Omega_{AG}(X, \tau)}$$

for every **Agop**-object (L, \leq, A) and every **AGS**-object (X, τ) . Then, the required result follows immediately from Lemma 2 and Lemma 3.

Remark 1. Since $\Omega_{AG}(X, \tau) = (\tau, \subseteq, X(\cap)\tau)$ is an object of **Smasagop** for every **AGS**-object (X, τ) , the adjoint situation in Theorem 1 can be restricted to an adjoint situation $(\eta, \varepsilon^r) : \Omega_{AG}^r \dashv Pt_{AG}^r : \mathbf{Smasagop}^{op} \rightarrow \mathbf{AGS}$, where $\Omega_{AG}^r (Pt_{AG}^r)$ is the co-domain (the domain) restriction of $\Omega_{AG} (Pt_{AG})$ and $\varepsilon_{(L, \leq, A)}^r = \varepsilon_{(L, \leq, A)}$ for every **Smasagop**-object (L, \leq, A) . An analogous adjoint situation can also be written for the category **Asagop** instead of **Smasagop**.

4 Spatial Aggregation Operators, Sober Aggregative Spaces and Their Duality

Spatiality and sobriety are two important notions that enable us to restrict the adjunction $\Omega_{AG} \dashv Pt_{AG}$ to an equivalence. To clarify this fact, we first start with their definitions:

Definition 5. (i) An **Agop**-object (L, \leq, A) is called *spatial* iff for all $a, b \in L$ with $a \not\leq b$, there exists a $G \in \mathfrak{F}(L)$ such that $a \in G$ and $b \notin G$.

(ii) An **AGS**-object (X, τ) is called *sober* iff for all $\mathcal{U} \in \mathfrak{F}(\tau)$, there exists a unique $x \in X$ such that $\mathcal{U} = \tau(x)$.

Proposition 4. Let (L, \leq, A) be an **Agop**-object, and (X, τ) an **AGS**-object.

(i) (L, \leq, A) is spatial iff $e_{(L, \leq, A)} : (L, \leq, A) \rightarrow \Omega_{AG}Pt_{AG}(L, \leq, A)$ is an **Agop**-isomorphism.

(ii) (L, \leq, A) is spatial iff (L, \leq, A) is isomorphic to $(\nu, \subseteq, Y(\cap)\nu)$ for some aggregative space (Y, ν) .

(iii) (X, τ) is sober iff $\eta_{(X, \tau)} : (X, \tau) \rightarrow Pt_{AG}\Omega_{AG}(X, \tau)$ is an **AGS**-isomorphism.

Proof. (i) Note first that (L, \leq, A) is spatial iff for all $a, b \in L$, $\Psi_a \subseteq \Psi_b$ implies $a \leq b$. Now, by assuming spatiality of (L, \leq, A) , this equivalence directly gives the injectivity of the underlying set map of $e_{(L, \leq, A)}$, and so does its bijectivity. It is easy to check that $e_{(L, \leq, A)}^{-1} : \Omega_{AG}Pt_{AG}(L, \leq, A) \rightarrow (L, \leq, A)$ is an **Agop**-morphism, and so $e_{(L, \leq, A)} : (L, \leq, A) \rightarrow \Omega_{AG}Pt_{AG}(L, \leq, A)$ is an **Agop**-isomorphism. Conversely, if $e_{(L, \leq, A)} : (L, \leq, A) \rightarrow \Omega_{AG}Pt_{AG}(L, \leq, A)$ is an **Agop**-isomorphism, then since $e_{(L, \leq, A)}^{-1} : \Omega_{AG}Pt_{AG}(L, \leq, A) \rightarrow (L, \leq, A)$ is an **Agop**-morphism, $\Psi_a \subseteq \Psi_b$ implies $a = e_{(L, \leq, A)}^{-1}(\Psi_a) \leq e_{(L, \leq, A)}^{-1}(\Psi_b) = b$ for all $a, b \in L$, so (L, \leq, A) is spatial.

(ii) If (L, \leq, A) is spatial, then we have from (i) that (L, \leq, A) is isomorphic to $(\nu, \subseteq, Y(\cap)_\nu)$ for the aggregative space $(Y, \nu) = Pt_{AG}(L, \leq, A)$. Conversely, suppose (L, \leq, A) is isomorphic to $(\nu, \subseteq, Y(\cap)_\nu)$ for some aggregative space (Y, ν) , i.e. there exists an **Agop**-isomorphism $u : (L, \leq, A) \rightarrow (\nu, \subseteq, Y(\cap)_\nu)$. Then, for $a, b \in L$ with $a \not\leq b$, since $u(a) \not\subseteq u(b)$, there exists at least one $z \in Y$ such that $z \in u(a)$ and $z \notin u(b)$. It is clear that $u(a) \in \nu(z)$ and $u(b) \notin \nu(z)$, and so $a \in u^\leftarrow(\nu(z))$ and $b \notin u^\leftarrow(\nu(z))$. Hence, we obtain the spatiality of (L, \leq, A) from the fact that $u^\leftarrow(\nu(z))$ is a filter of (L, \leq, A) .

(iii) follows from that for a given **AGS**-object (X, τ) , (X, τ) is sober iff the underlying set map of $\eta_{(X, \tau)}$ is a bijection iff $\eta_{(X, \tau)} : (X, \tau) \rightarrow Pt_{AG}\Omega_{AG}(X, \tau)$ is an **AGS**-isomorphism.

Corollary 1. *The full subcategory **SpatAgop** of **Agop** of all spatial objects is dually equivalent to the full subcategory **SobAGS** of **AGS** of all sober objects.*

Proof. Since Proposition 4 (i) and (iii) verify that $Fix(\varepsilon) = \mathbf{SpatAgop}$ and $Fix(\eta) = \mathbf{SobAGS}$, the assertion follows from Theorem 1 and Proposition 1 (i).

Proposition 5. ***SpatAgop** and **SobAGS** are reflective subcategories of **Agop** and of **AGS** with reflectors $\Omega_{AG}^{op}Pt_{AG}^{op}$ and $Pt_{AG}\Omega_{AG}$, and the reflection arrows $e_{(L, \leq, A)}$ and $\eta_{(X, \tau)}$, respectively.*

Proof. Since $e_{(L, \leq, A)} = \varepsilon_{(L, \leq, A)}^{op}$ is obviously an epimorphism in **Agop**, and $Fix(\varepsilon) = \mathbf{SpatAgop}$ and $Fix(\eta) = \mathbf{SobAGS}$, Proposition 1 (ii) directly yields the claimed result.

Proposition 6. *Let (L, \leq, A) be an **Agop**-object, and (X, τ) an **AGS**-object.*

(i) $\Omega_{AG}(X, \tau)$ is spatial, (ii) $Pt_{AG}(L, \leq, A)$ is sober.

Proof. (i) is immediate from Proposition 4 (ii). To see (ii), let us first consider the fact that $Pt_{AG}(L, \leq, A) = (\mathfrak{F}(L), \Psi(L))$. Then, the sobriety of $Pt_{AG}(L, \leq, A)$ follows from the observation that for all $\mathcal{U} \in \mathfrak{F}(\Psi(L))$, $G = \{a \in L \mid \Psi_a \in \mathcal{U}\}$ is the unique element of $\mathfrak{F}(L)$ with the property that $\mathcal{U} = [\Psi(L)](G)$.

Proposition 7. *The full subcategory **SpatAsagop** of **Asagop** of all spatial objects, the full subcategory **SpatSmasagop** of **Smasagop** of all spatial objects and **SpatAgop** are equivalent to each others.*

Proof. Since **SpatSmasagop** is a full subcategory of **SpatAgop**, the inclusion functor **SpatSmasagop** \hookrightarrow **SpatAgop** is a full and faithful functor. For every **SpatAgop**-object (L, \leq, A) , by Proposition 4 (i) and Proposition 6 (i), $\Omega_{AG}Pt_{AG}(L, \leq, A)$ is a **SpatSmasagop**-object, and $e_{(L, \leq, A)} : (L, \leq, A) \rightarrow \Omega_{AG}Pt_{AG}(L, \leq, A)$ is a **SpatAgop**-isomorphism. This proves that the inclusion functor **SpatSmasagop** \hookrightarrow **SpatAgop** is isomorphism-dense, and hence an equivalence. Similarly, the inclusion functor **SpatAsagop** \hookrightarrow **SpatAgop** is an equivalence, which completes the proof.

Corollary 2. $SpatAgop^{op} \sim SpatAsagop^{op} \sim SpatSmasagop^{op} \sim SobAGS.$

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