

Power-Aware Throughput Control for Database Management Systems

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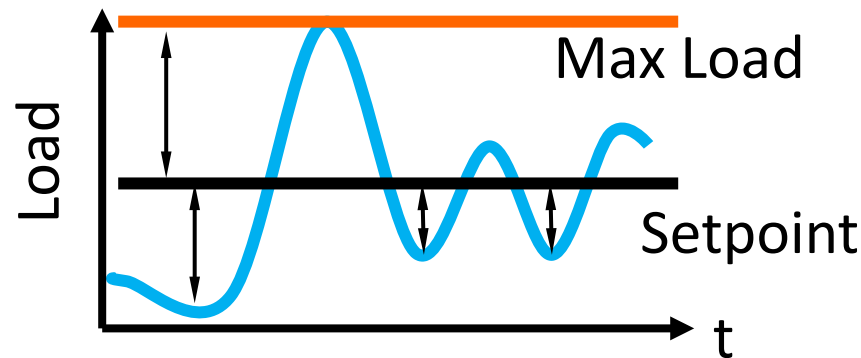
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Problem Overview and Motivation

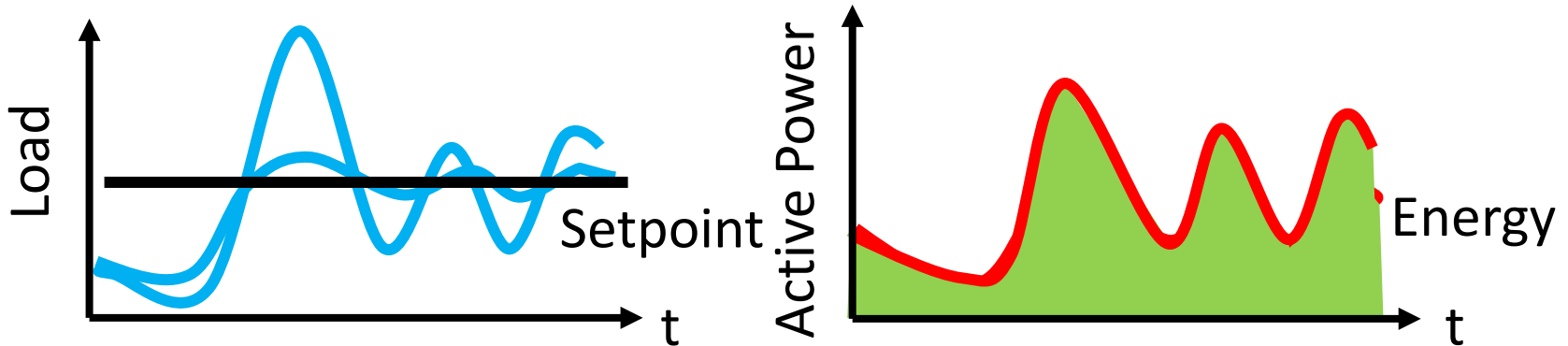
- Data centers are energy starving
- Database management system (DBMS) is one of the major services in data centers
- Control the DBMS throughput to save energy



- Power saving to the setpoint: low power hardware states
- Balanced performance: performance control

Power-Aware Throughput Control (PAT)

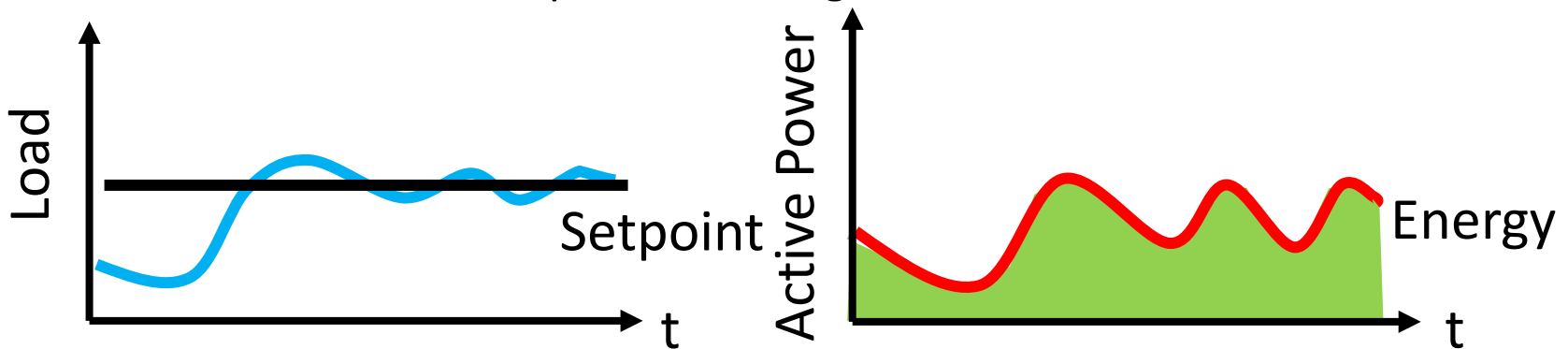
- **Goal:** minimize power consumption of a database system while maintaining its desired performance
- **Challenges**
 - Energy and DBMS: power consumption minimization
 - Control and DBMS: performance guarantee



- **Contributions:**
 - Energy profiling in DBMS
 - Feedback control design in DBMS

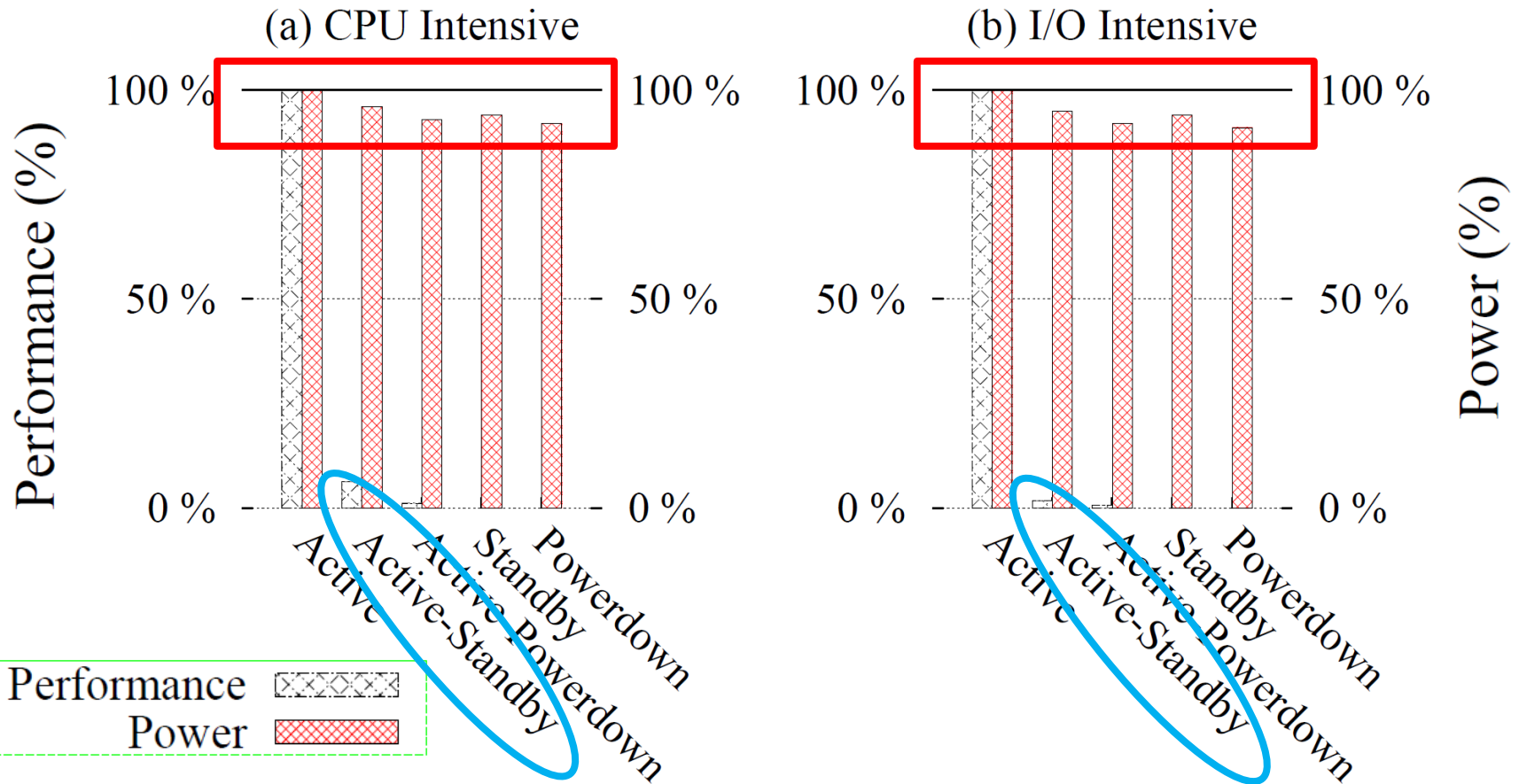
Power-Aware Throughput Control (PAT)

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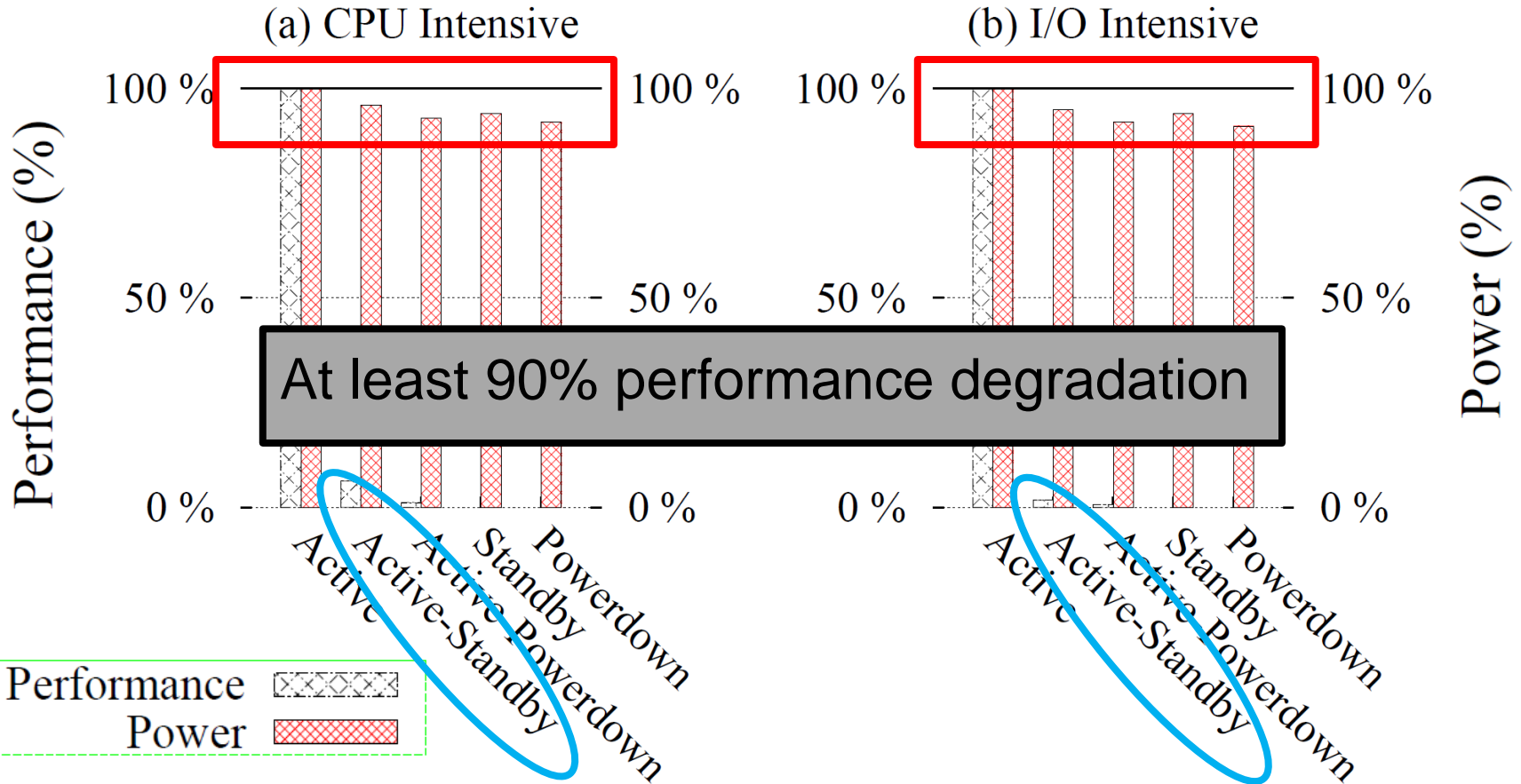
Characterization Study – Memory



- Low power state in memory may not be a good choice for power saving in DBMS workloads

Character

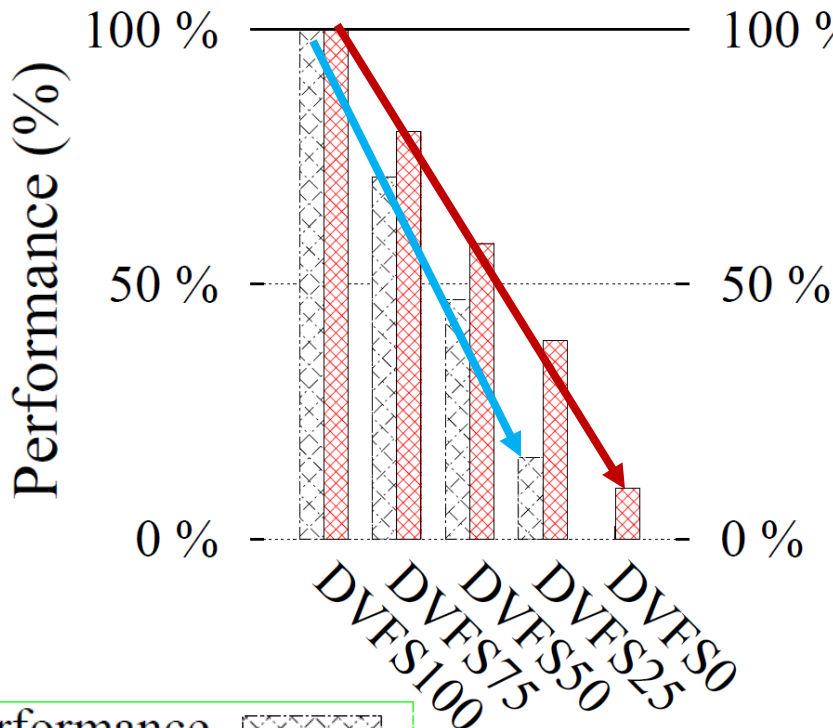
Less than 20% power savings



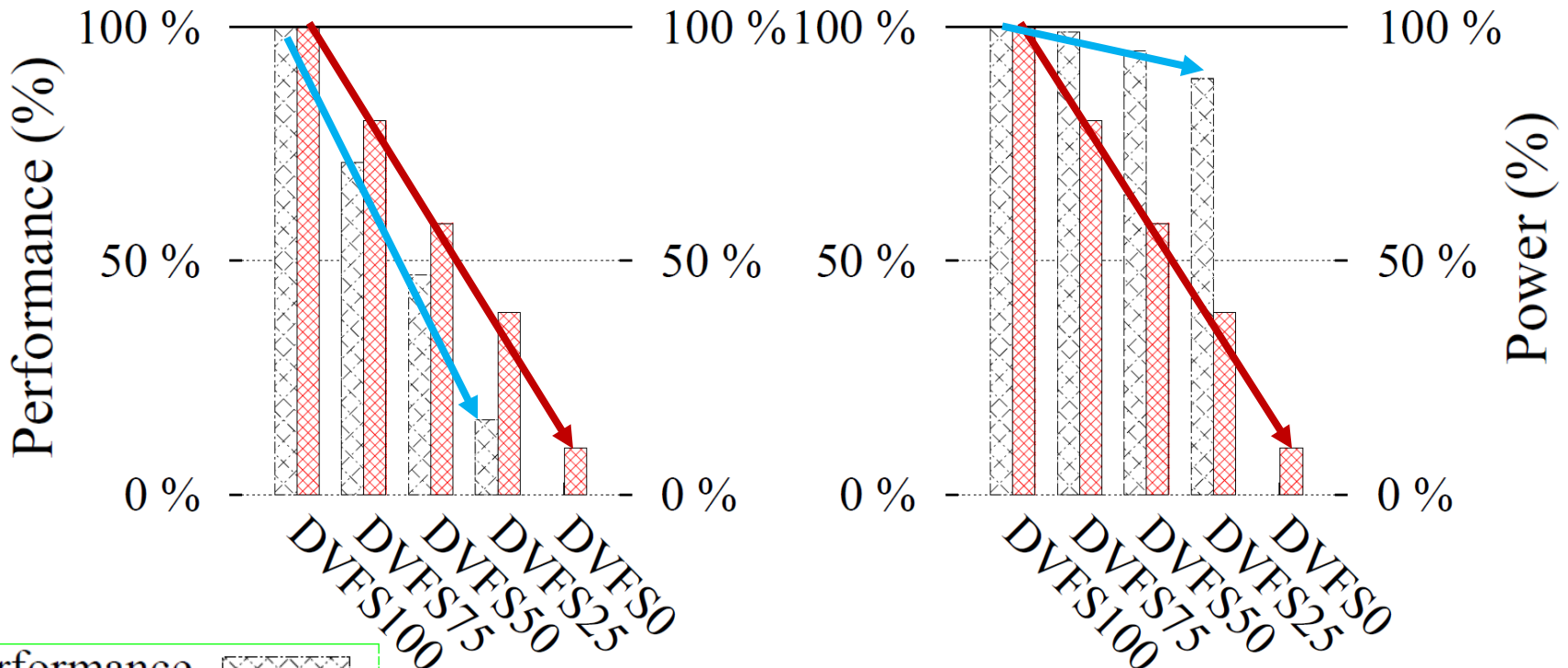
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Characterization Study – CPU

(c) CPU Intensive



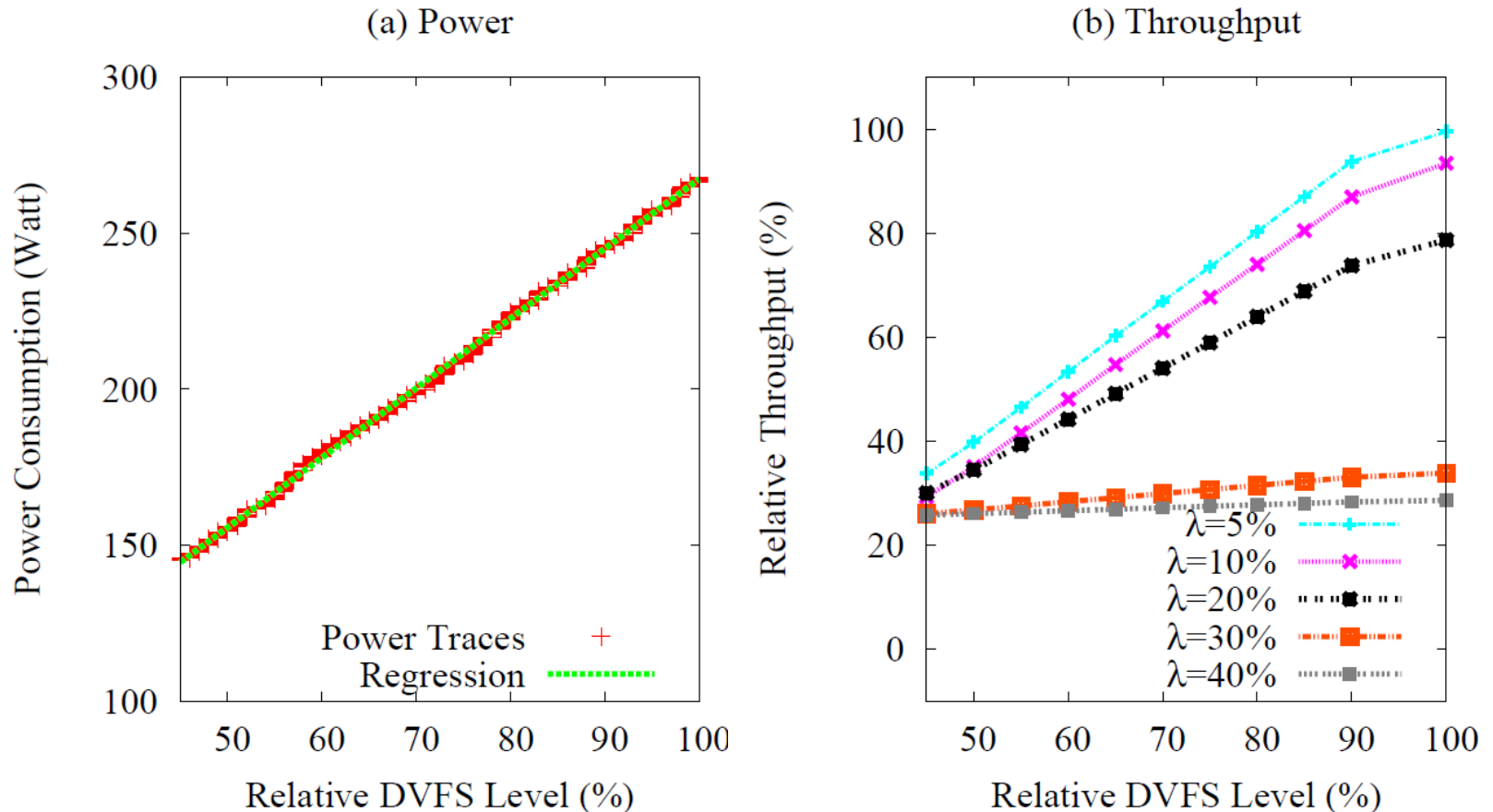
(d) I/O Intensive



Performance 
Power 

- CPU provides great power-saving and the relationship can be approximated as a linear model

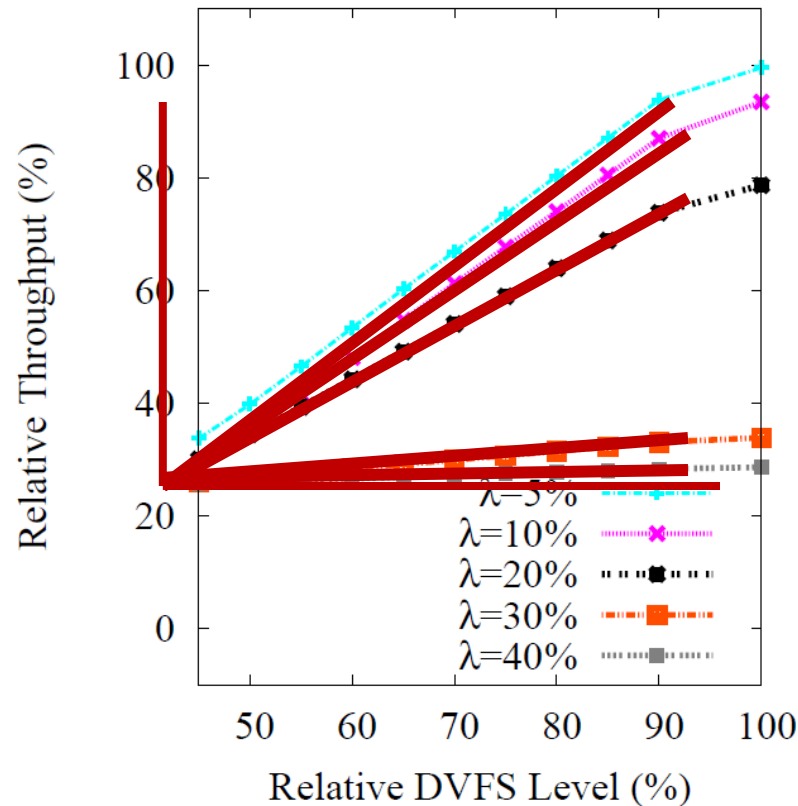
Characterization Study – DBMS Workloads



- The ratio of I/O intensive queries (λ) is the key in the workload statistics for power control

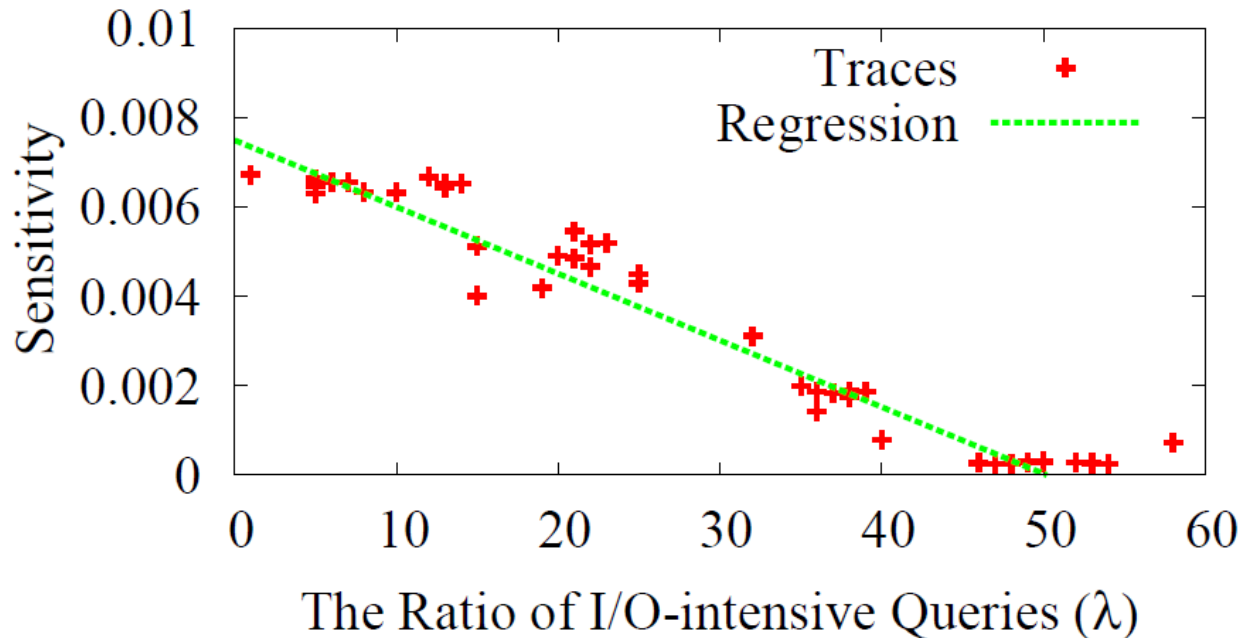
Characterization Study – DBMS Workloads

(b) Throughput



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Characterization Study – DBMS Workloads



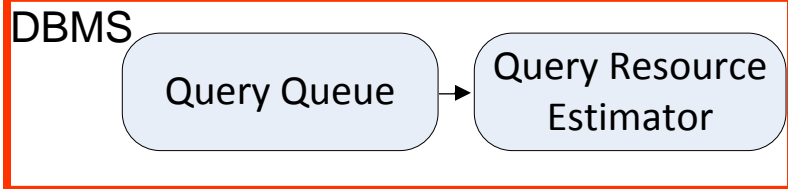
- The ratio of I/O intensive queries (λ) is the key in the workload statistics for power control

Characterization Study – Insights

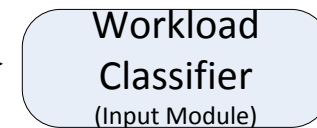
- Low power state in CPU provides great opportunities for power-saving
- DBMS throughput and CPU frequency can be approximated as a linear model
 - $r = A * f + B$
- The ratio of I/O intensive queries (λ) is the key in the workload statistics for power control
 - $r = \lambda A * f + B$

Framework

Query Workload

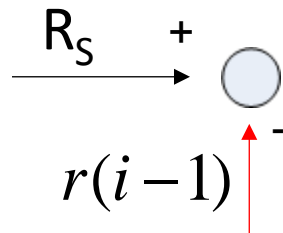


Workload Statistics



λ

PAT Control

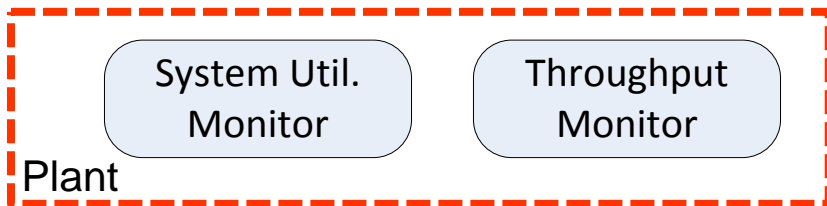


$\Delta r(i)$

PI Controller

Controller

$f(i)$



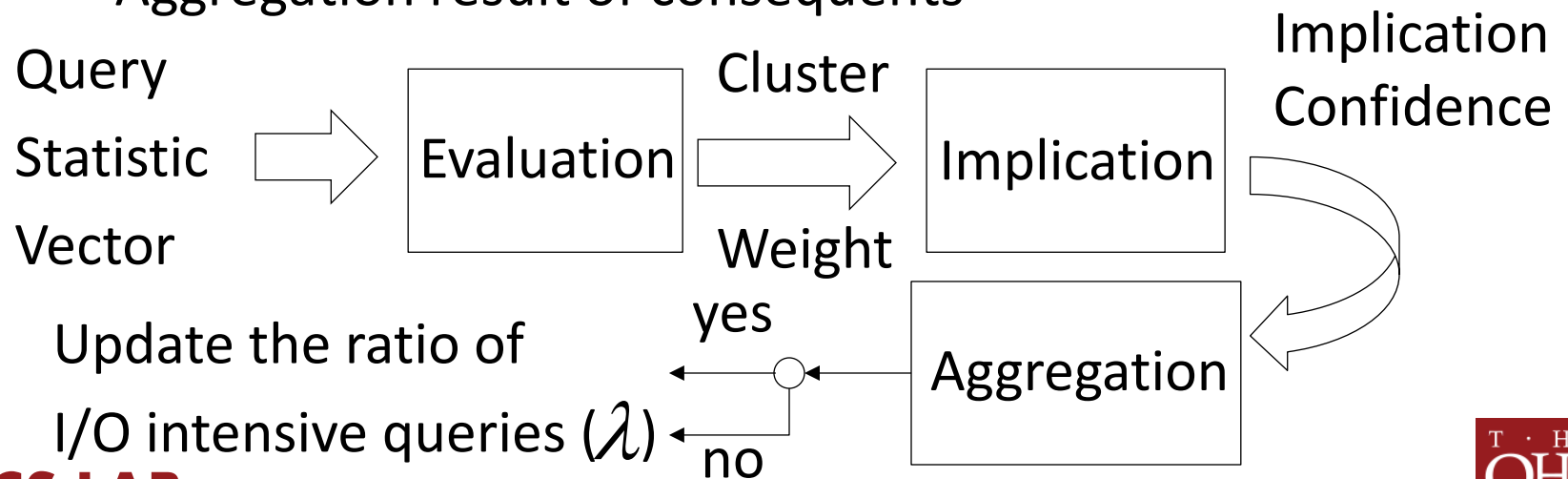
New CPU State

CPU Power State Modulator

Actuator

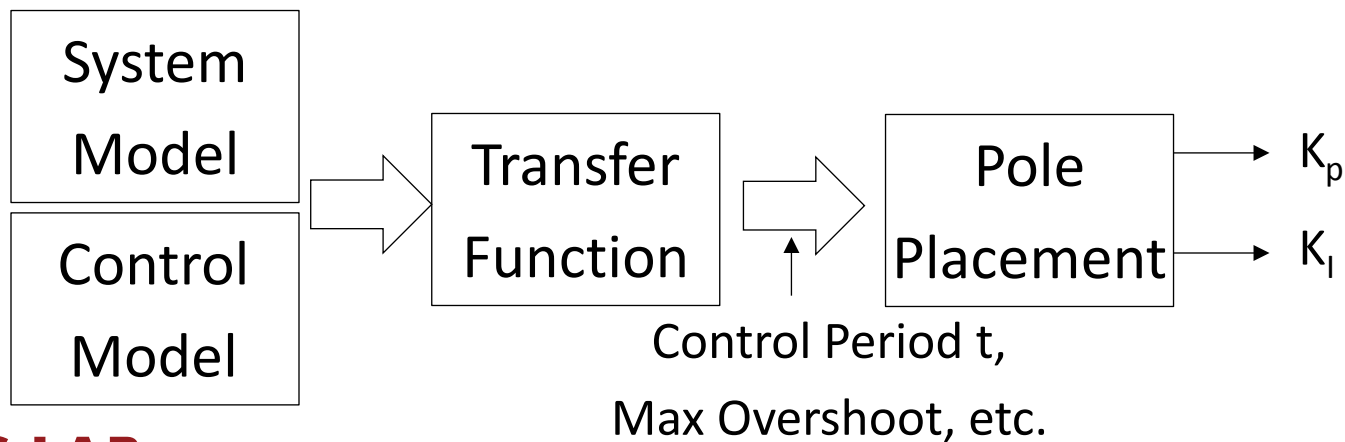
Fuzzy Workload Classifier

- One fuzzy rule base
 - Compute the membership of query to each learnt cluster
- The procedure of workload classification
 - Evaluation of antecedents
 - Implication calculation of consequents
 - Aggregation result of consequents



The Controller Design

- System model
 - $\Delta r(i) = \lambda A * f(i) + B$
- Control model (PI control)
 - Zero steady-state error
 - Short settling time
 - Stability



Methodology

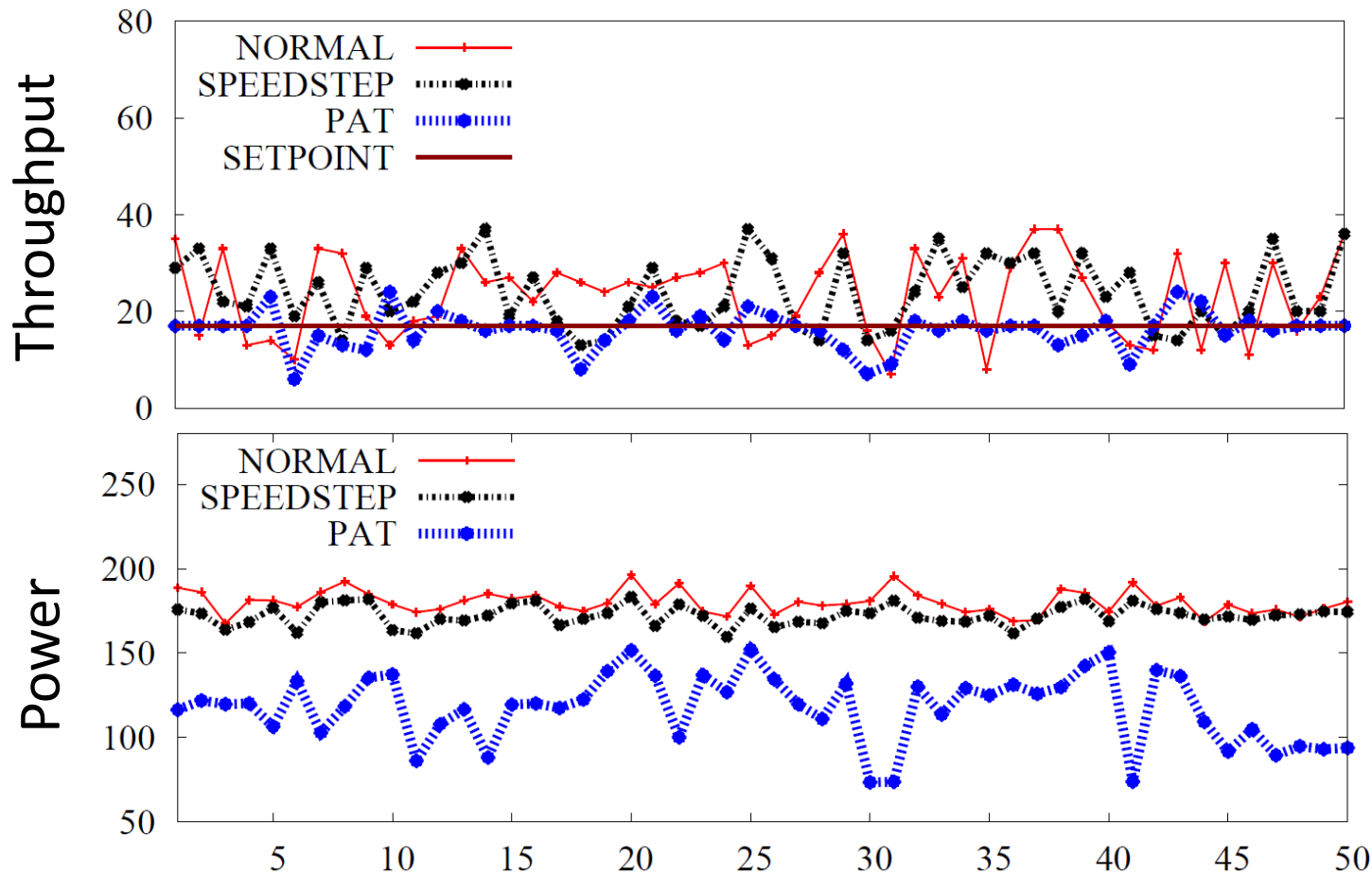
- Testbed
 - Hardware: a DELL PowerEdge R710 with 12-core Intel Xeon E5645, 16GB memory and 1TB storage
 - Software: Redhat 5 (V3.0.0) + PostgreSQL (V8.3.18)
 - Workload: TPC and SDSS
 - Power measurement: WattsUp power meter
- System Contention
 - Ideal: running DBMS alone
 - Competing: running DBMS with other processes with equal priority.
 - Preemptive: running DBMS with other processes with higher priority.

Baselines

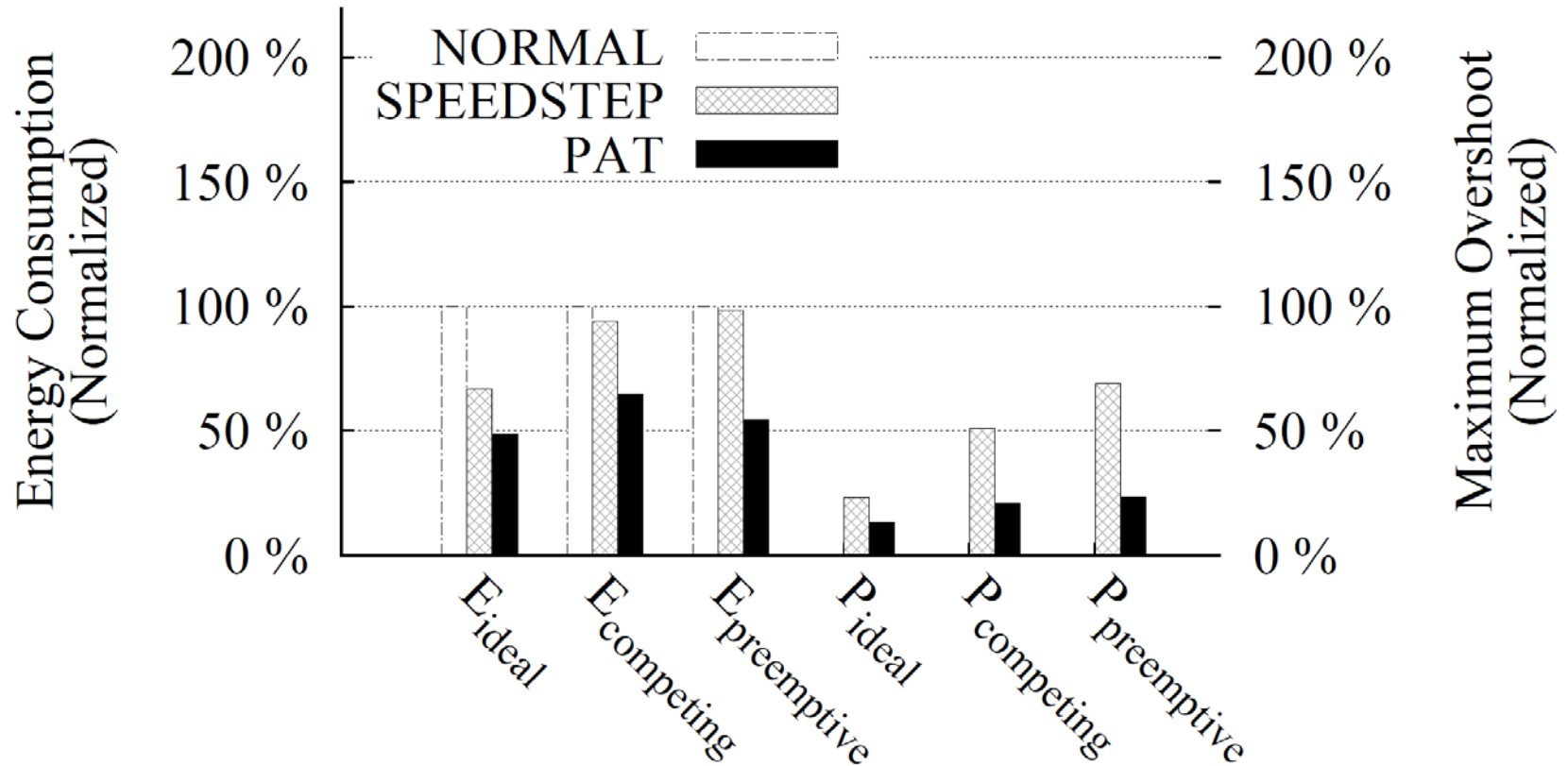
- State-of-the-Practice Baselines
 - NORMAL: no power management, performance first
 - SPEEDSTEP: Intel™ build-in power management

- State-of-the-Art Baselines
 - TRADITION: Open-loop control based on known workload statistics
 - HEURISTIC: Ad-hoc control based on performance-power model
 - SCTRL: System-level power control for DBMS performance

A Snapshot of PAT in Competing Setup

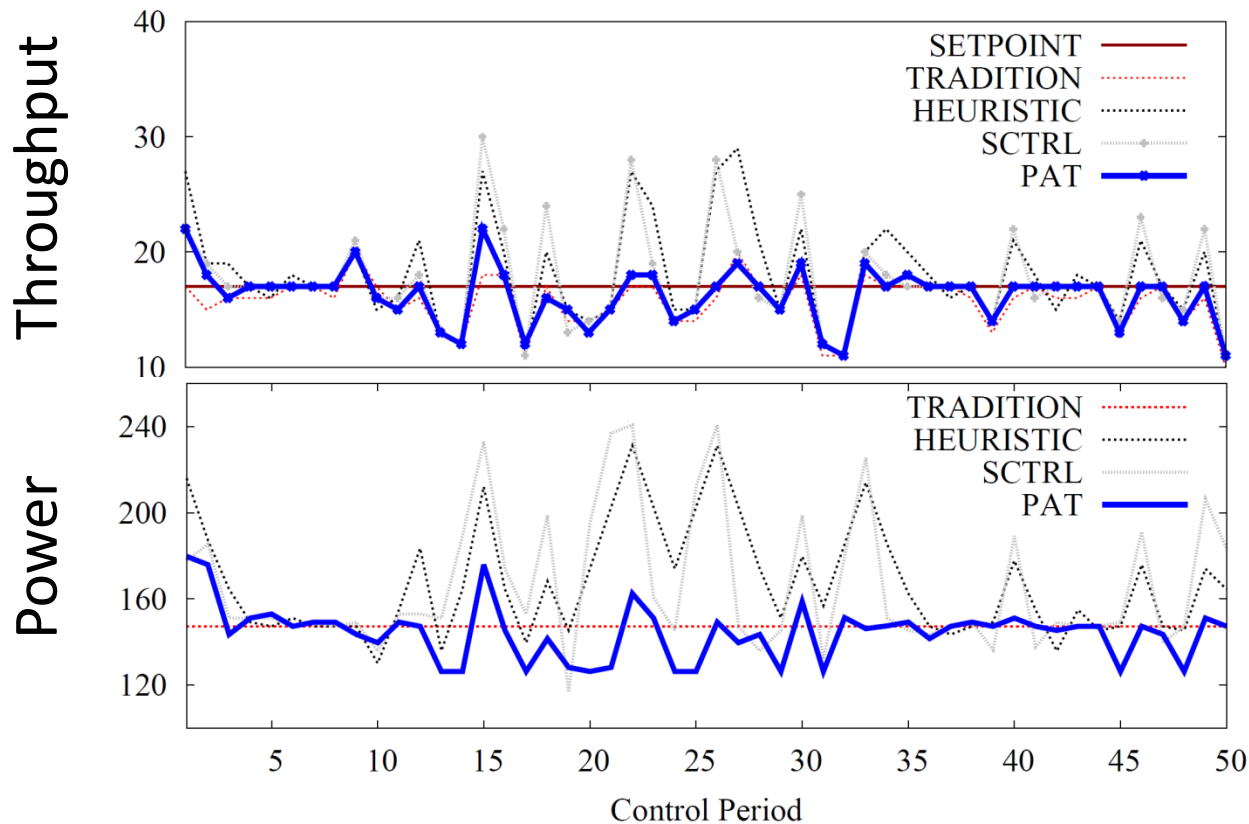


How PAT Performs in All Scenarios

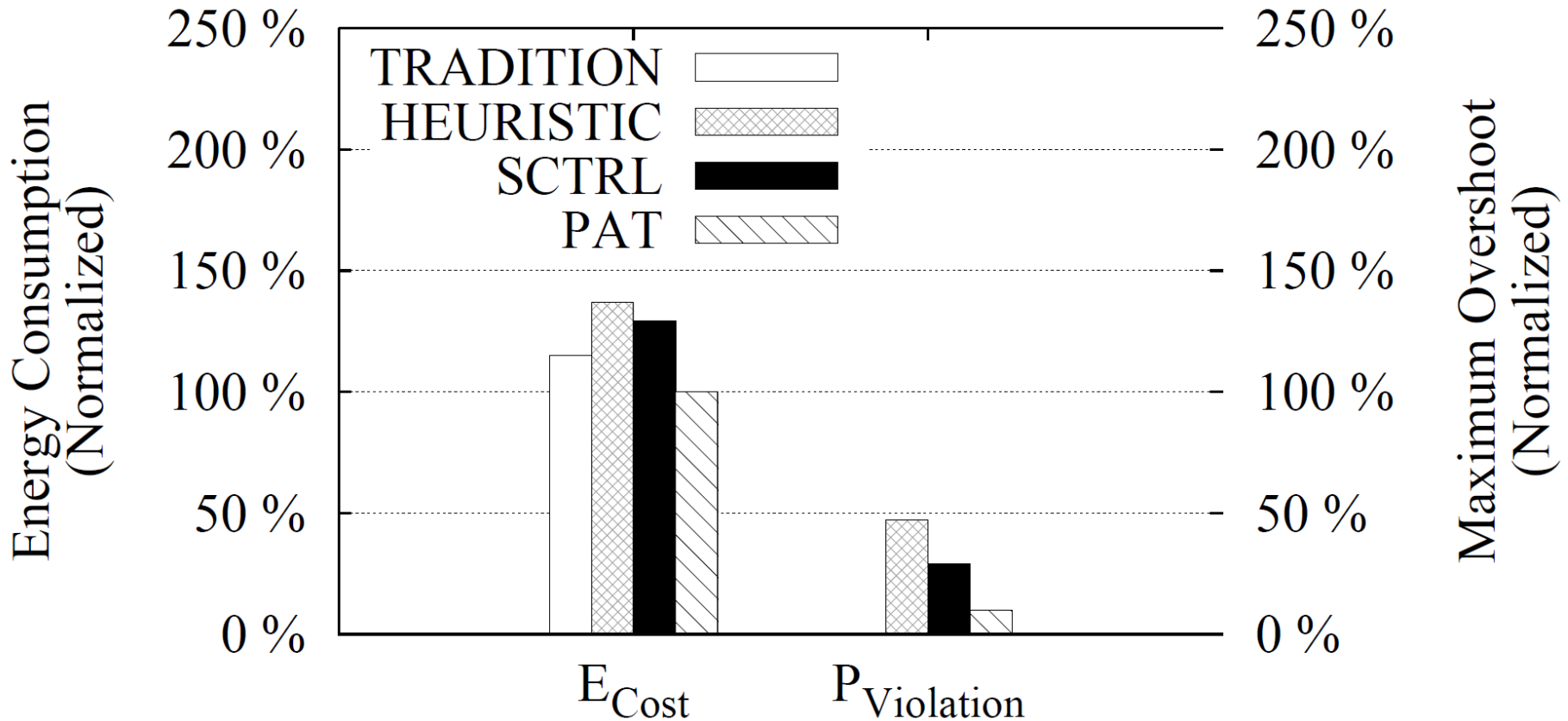


- PAT saves energy cost (15% more than the SPEEDSTEP, 51.3% more than NORMAL)

A Snapshot of Comparison



Comparison with Other Control Methods



- PAT has the smallest overshoot and the best energy saving

Conclusion

- Minimizing power consumption of DBMS under a user-specified performance bound
- Controller design on the system characteristics is the solution
- Empirically study the relationship between energy and DBMS processing
- PAT: a feedback control framework with system characteristics

- Thanks

Acknowledge

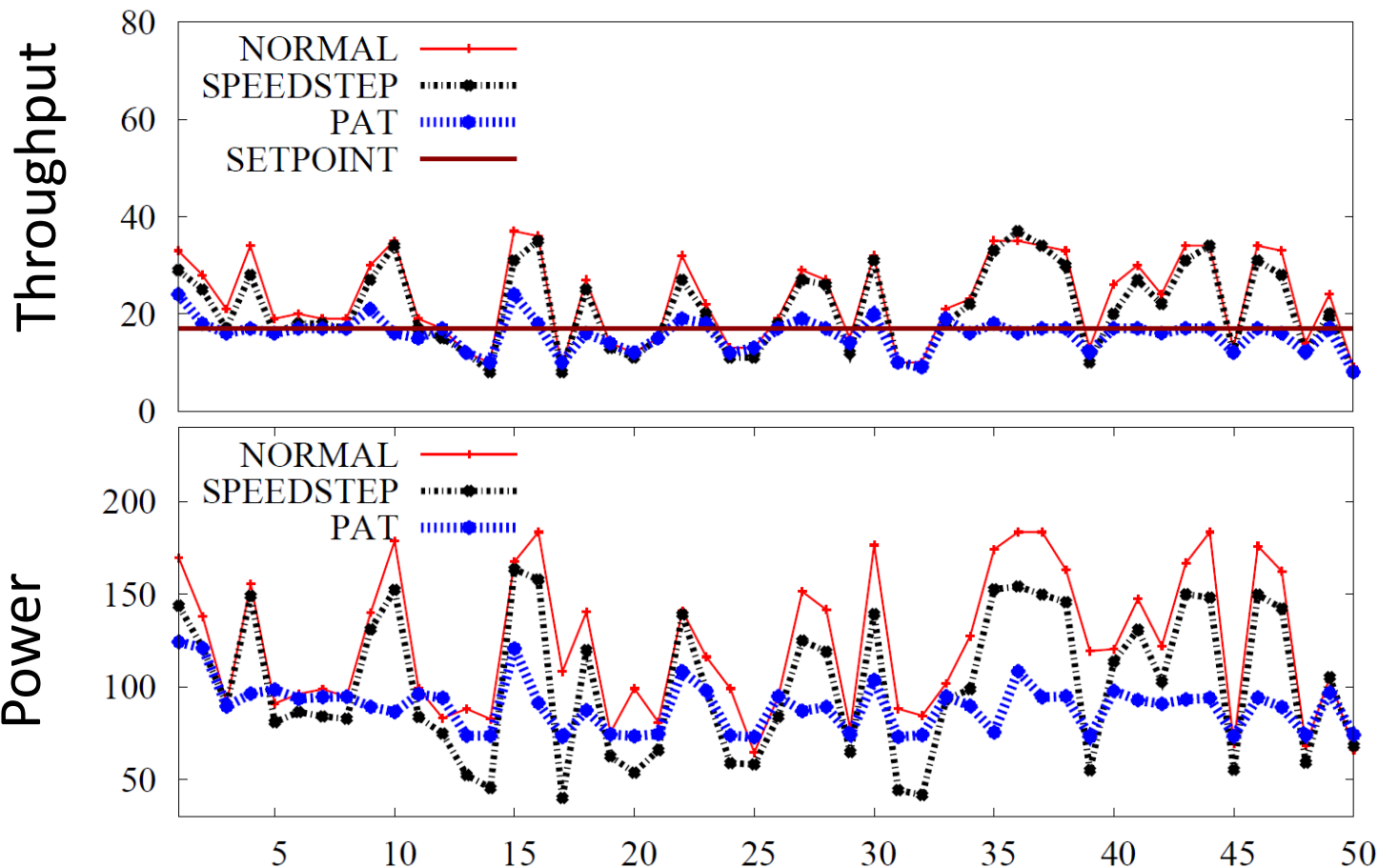
- This research was supported by the US National Science Foundation (NSF) grants IIS-1117699, IIS-1156435, and CSR-1143607



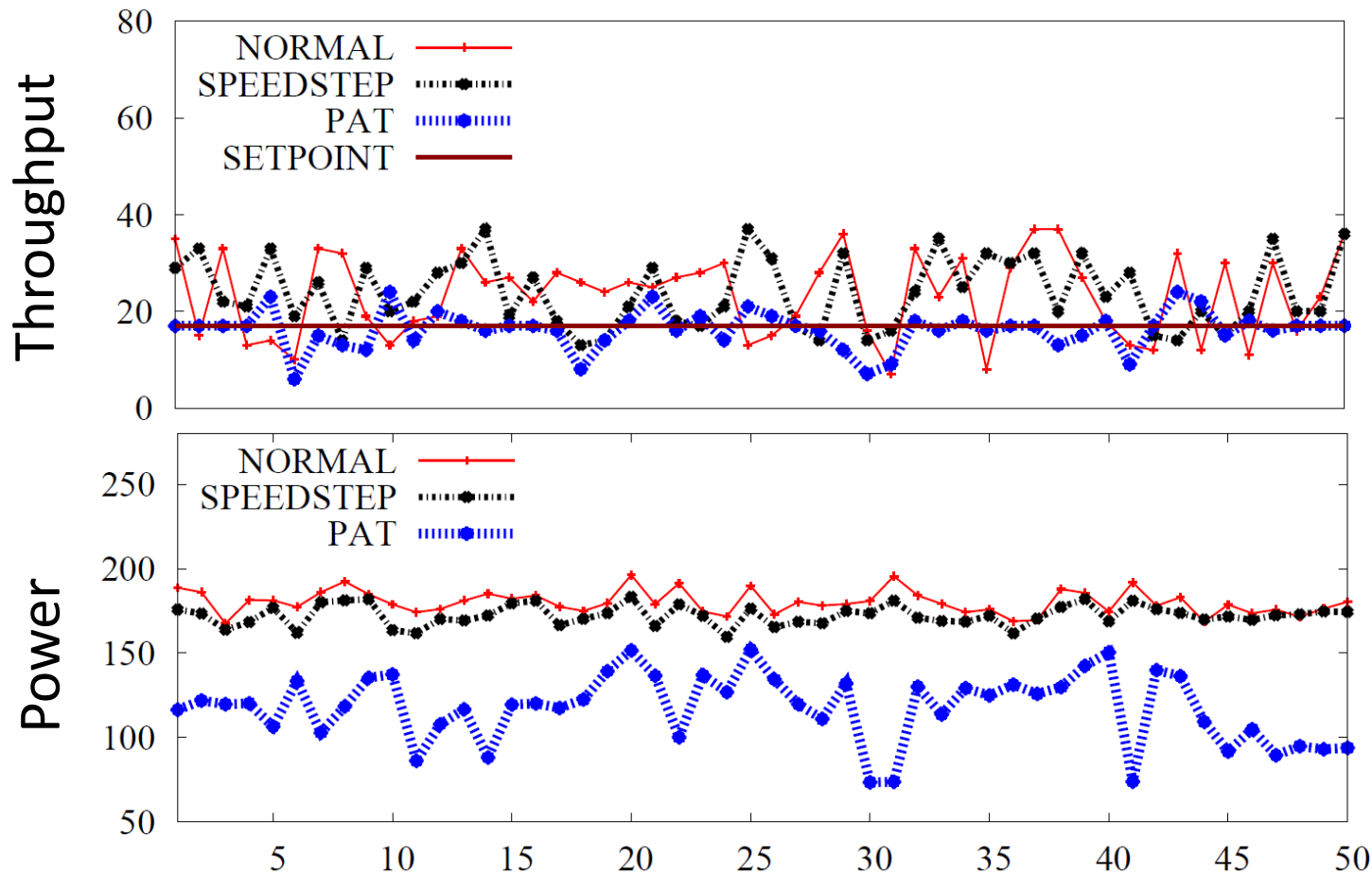
Related Work

- Control in DBMS
 - Feedback Control in DBMS [Tu et al. VLDB 2007]
- Other Applications
 - Power Shaving
 - Power Over-commitment
 - Proportional Energy Consumption

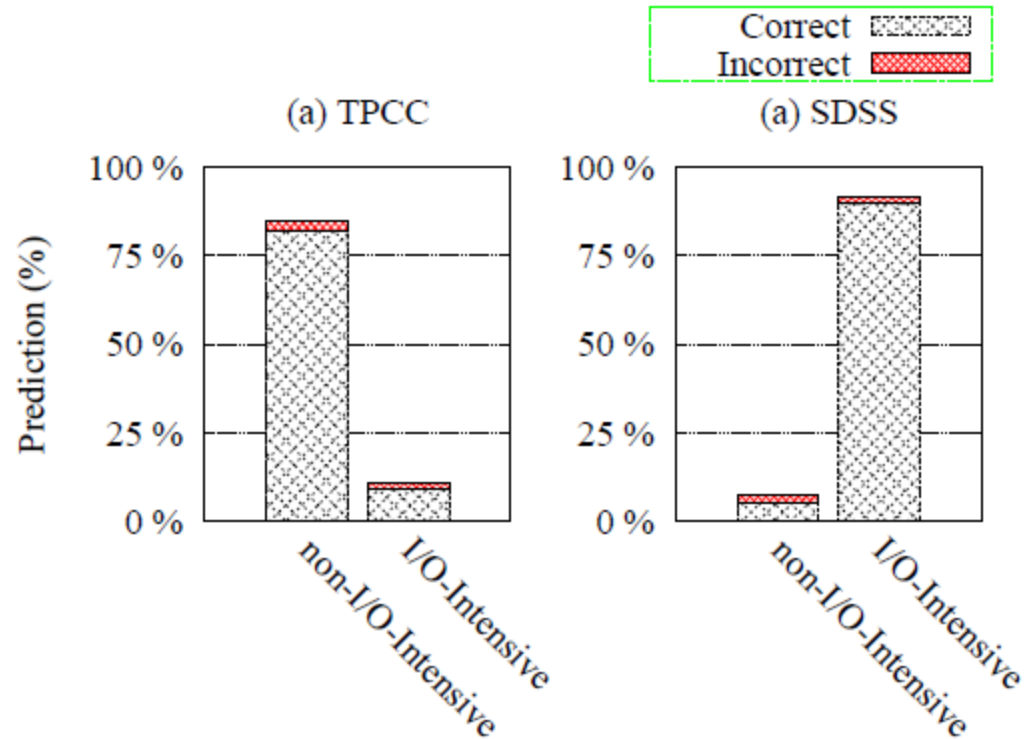
A Snapshot of PAT in Ideal Setup



A Snapshot of PAT in Competing Setup

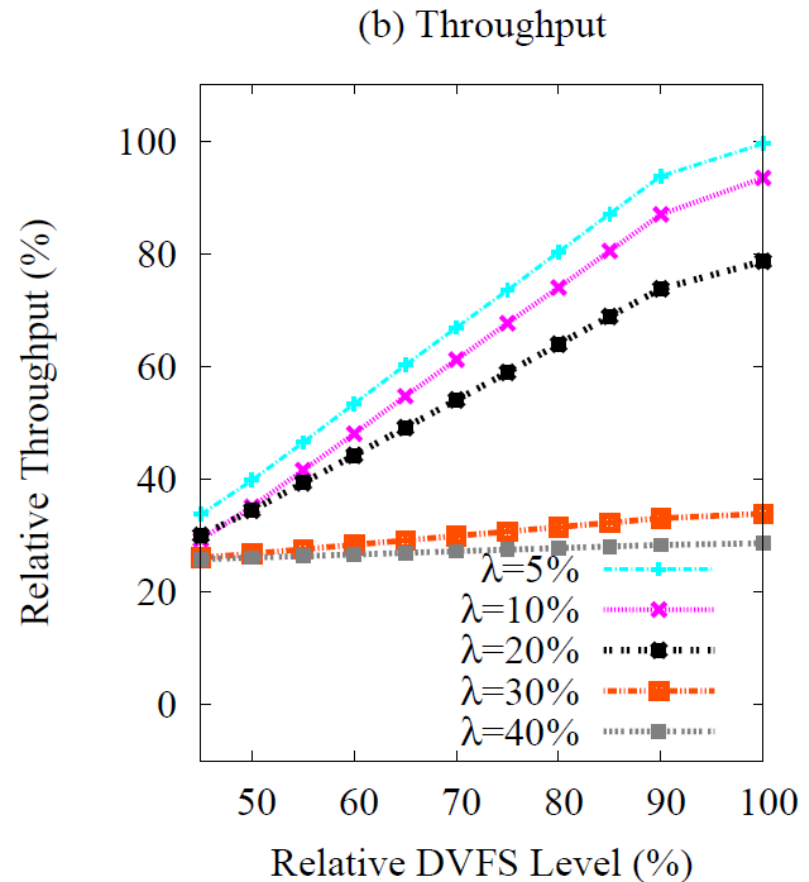
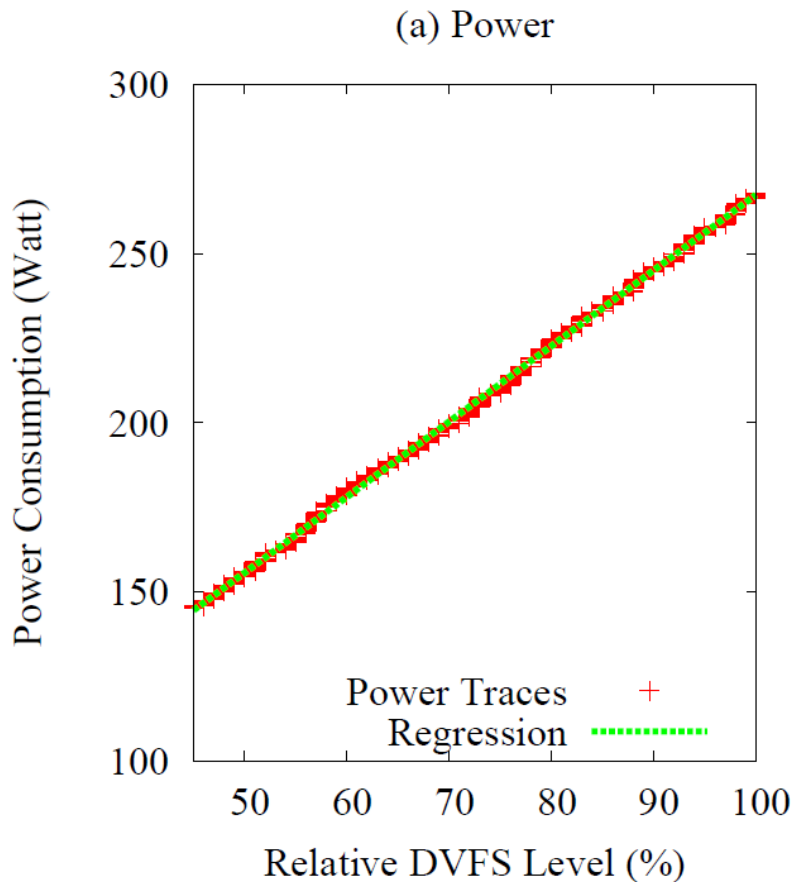


Backup: Errors From Modeling



- FWC could produce almost 10% error in the system, The designed maximum overshoot is 40% in PAT.

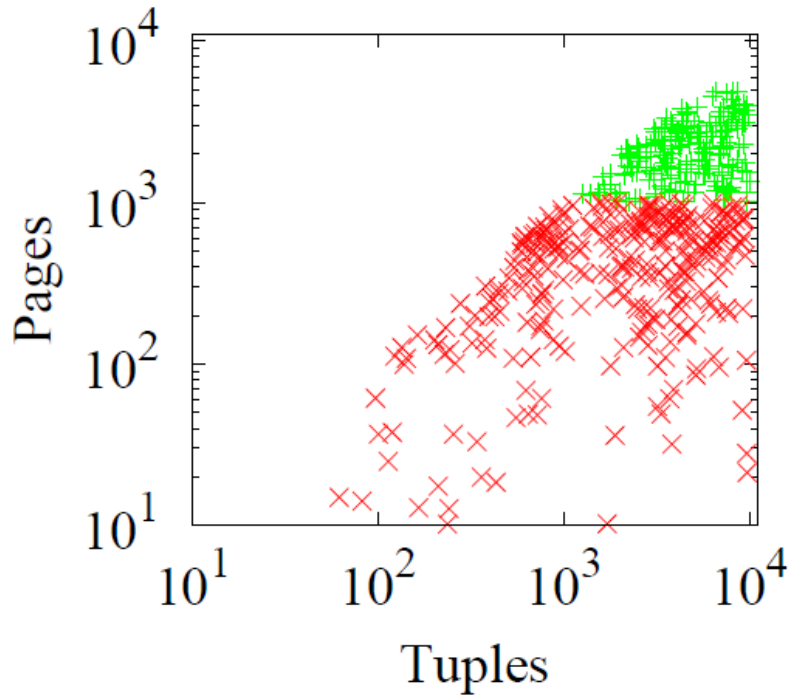
Backup Characterization Study – Energy and



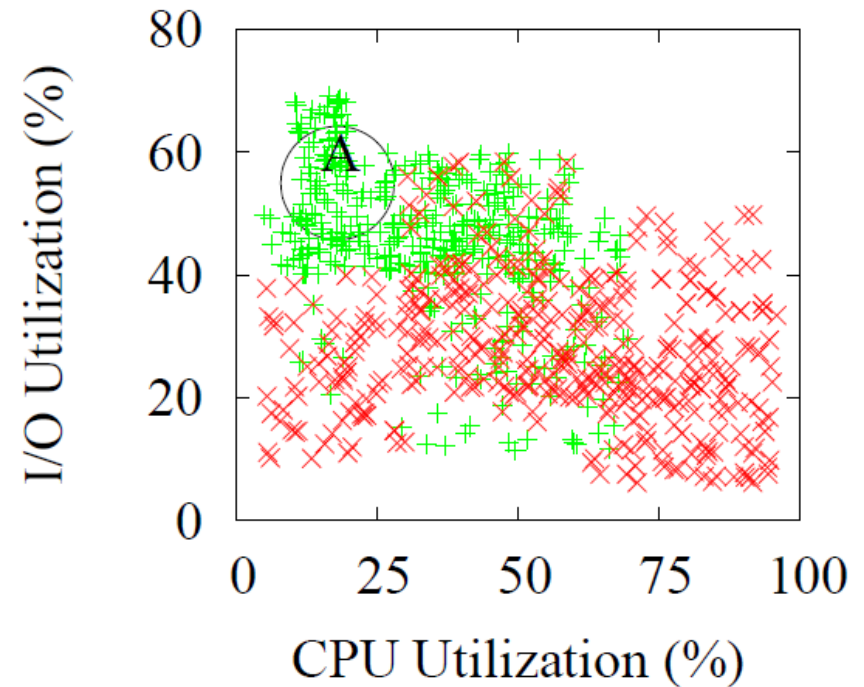
- DBMS throughput and CPU frequency can be approximated as a linear model

Backup: Rule-based Workload Classifier

(a) Estimated



(b) Runtime



I/O Queries	+
non-I/O Queries	x

Outline

- Problem Overview & Motivation
- Characterization Study: DBMS and Energy
- Overall Framework
- Fuzzy Workload Classifier
- The Controller Design
- Evaluation
- Conclusion

Fuzzy Workload Classifier

- One fuzzy rule base
 - If $[d_{CPU}^i, d_{I/O}^i] \in \text{cluster } X_j$,
then $[u_{CPU}^i, u_{I/O}^i] = M_j [d_{CPU}^i, d_{I/O}^i] + N_j$
- The procedure of workload classification
 - Evaluation: $[d_{CPU}^i, d_{I/O}^i] \rightarrow [u_{CPU}^i, u_{I/O}^i]$ per X_j
 - Implication calculation: $\frac{\sum(p_j) - p_i}{\sum(p_j)} \rightarrow t_j^i$ of X_j
 - Aggregation result: $[\sum t_j^i u_{CPU}^i, \sum t_j^i u_{I/O}^i]^T$