



TEXAS A&M UNIVERSITY  
Engineering



# Does logic locking work with EDA tools?

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# Supply Chain Security of Hardware

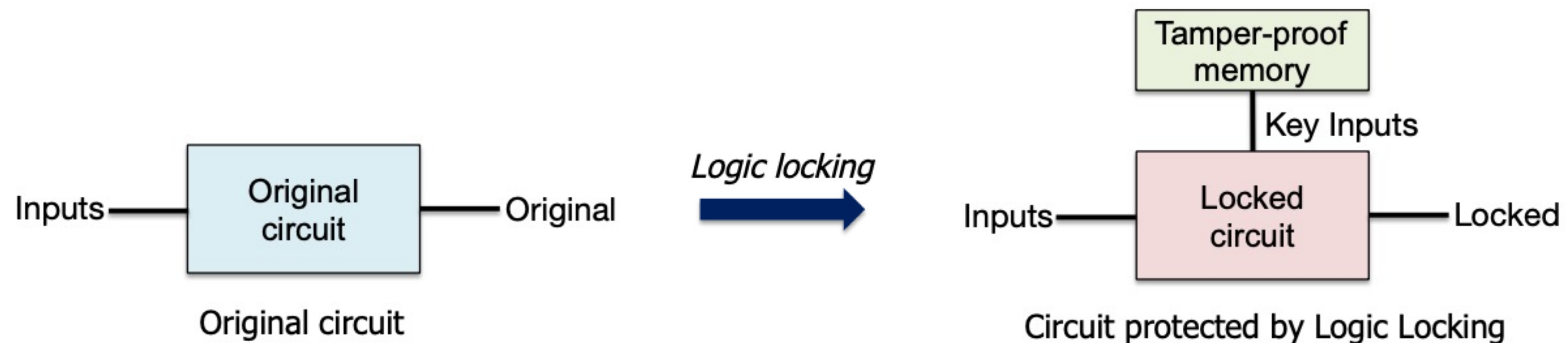
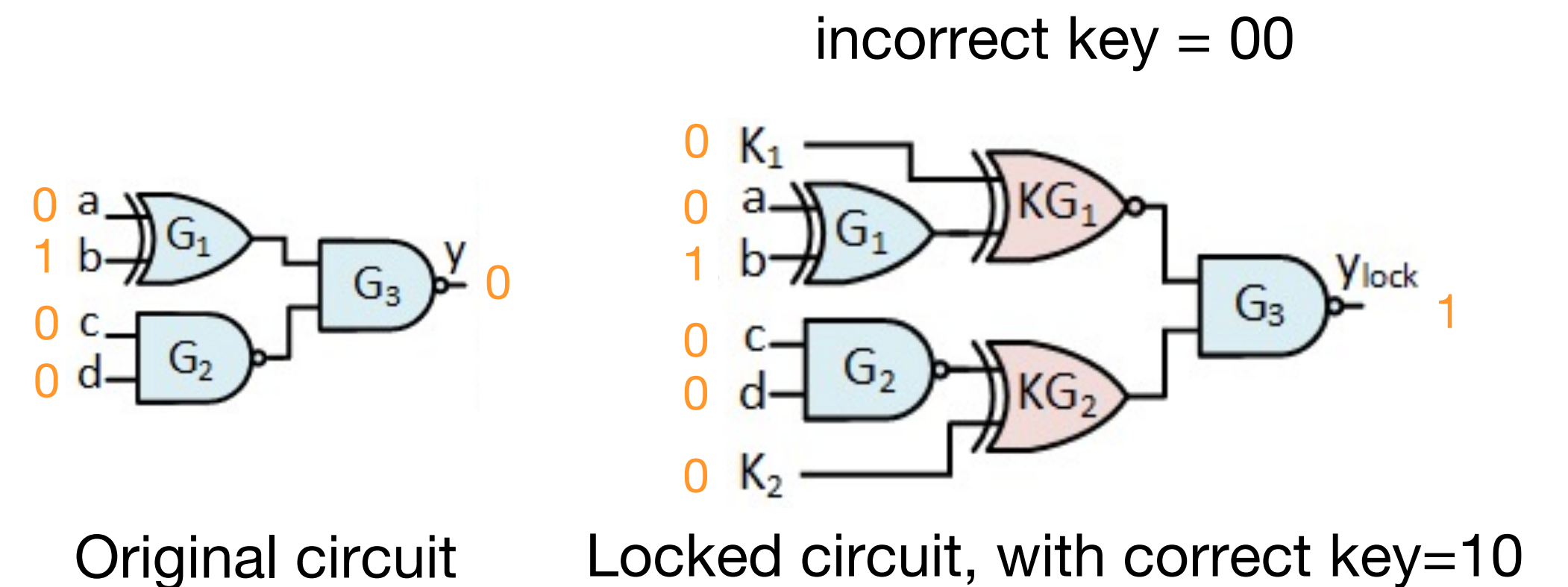
| Threat \ Countermeasure  | Countermeasure |                     |                 |               |
|--------------------------|----------------|---------------------|-----------------|---------------|
|                          | Watermarking   | Split manufacturing | IC camouflaging | Logic locking |
| IP piracy & Overbuilding | ✓              | ✓                   | ✓               | ✓             |
| Reverse engineering      | X              | ✓                   | ✓               | ✓             |
| IC counterfeiting        | ✓              | X                   | X               | ✓             |
| Hardware Trojan          | X              | ✓                   | X               | ✓             |

✓ denotes a successful defense, X denotes an unsuccessful defense

Logic locking can defend against all these threats

# Logic Locking

- Modifying logic and adding key inputs
- Only with correct key, the output is correct
- Secret key is stored in tamper-proof memory



# Existing Logic Locking Techniques

| Attack \ Defense                                  | Query-based attack |          |               | Structural attack |      |      |
|---|--------------------|----------|---------------|-------------------|------|------|
|   | Sensitization      | SAT, SMT | AppSAT, 2-DIP | SPS, ATR          | FALL | SAIL |
| XOR-based (random, strong, fault-based LUT-based) | X                  | X        | X             | ✓                 | ✓    | X    |
| Point-function (AND-tree, SARLock, Anti-SAT)      | ✓                  | ✓        | ✓             | X                 | ✓    | ✓    |
| CAC-HD, CAC-flex                                  | ✓                  | ✓        | ✓             | ✓                 | X    | ✓    |
| CAC-rem   | ✓                  | ✓        | ✓             | ✓                 | ✓    | ✓    |

✓ denotes a successful defense, X denotes an unsuccessful defense

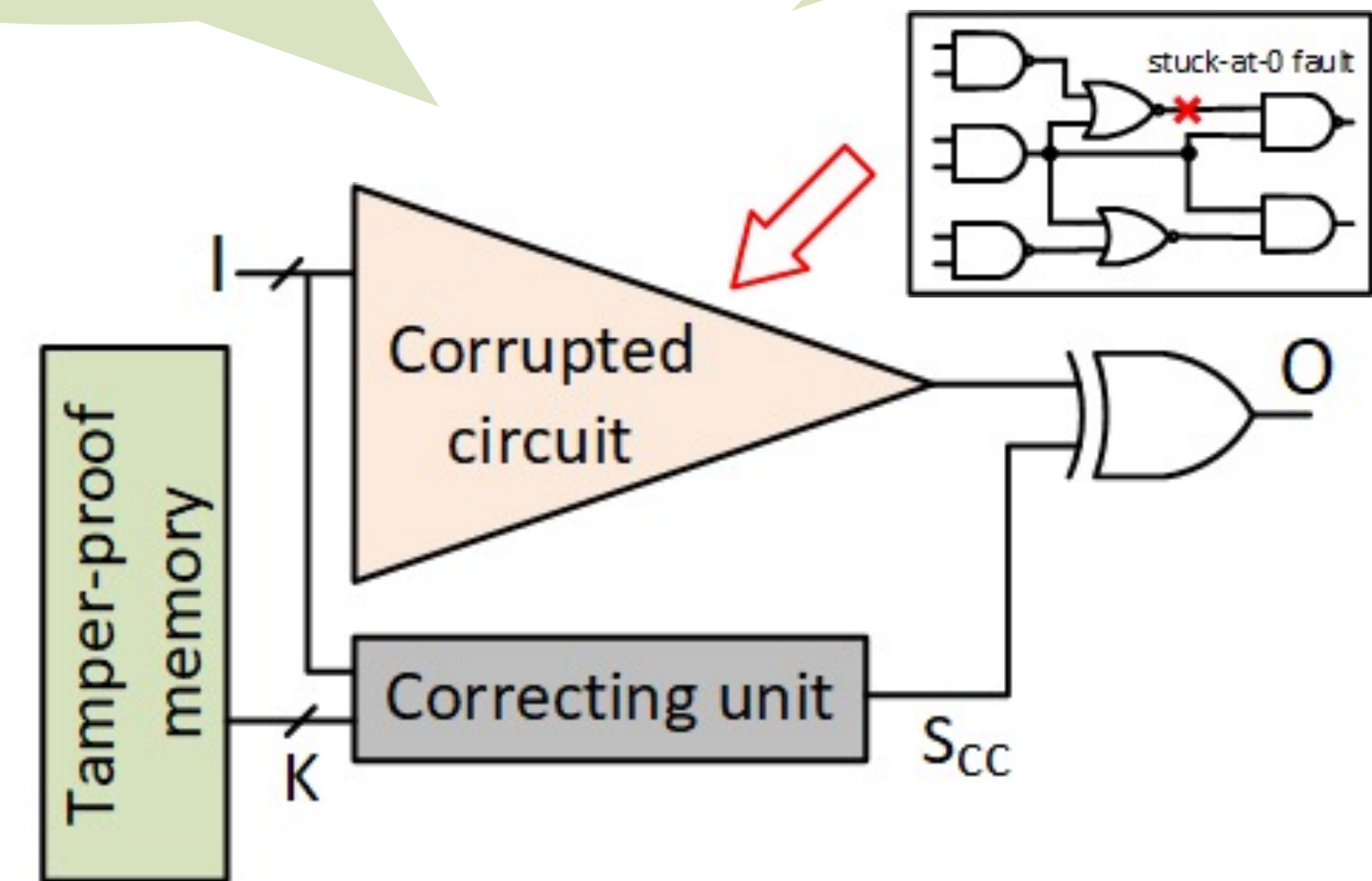
- Point-function and CAC techniques are resilient to query-based attacks.
- CAC-rem is the only technique that defends against all existing attacks.

# Corrupted and Corrected (CAC)

- Aka stripped-functionality logic locking
- Components in CAC-locked circuit
  - Corrupted circuit
    - Protected input pattern (PIP)
    - $in \notin PIP \Leftrightarrow f_{orig}(in) \neq f_{cp}(in)$
  - Correcting unit
    - $key = key_c \Leftrightarrow \forall in, f_{orig}(in) = f_{lock}(in, key)$
- CAC-rem is unbroken since 2020

Resilient against query-based attacks

Resilient against structural attacks



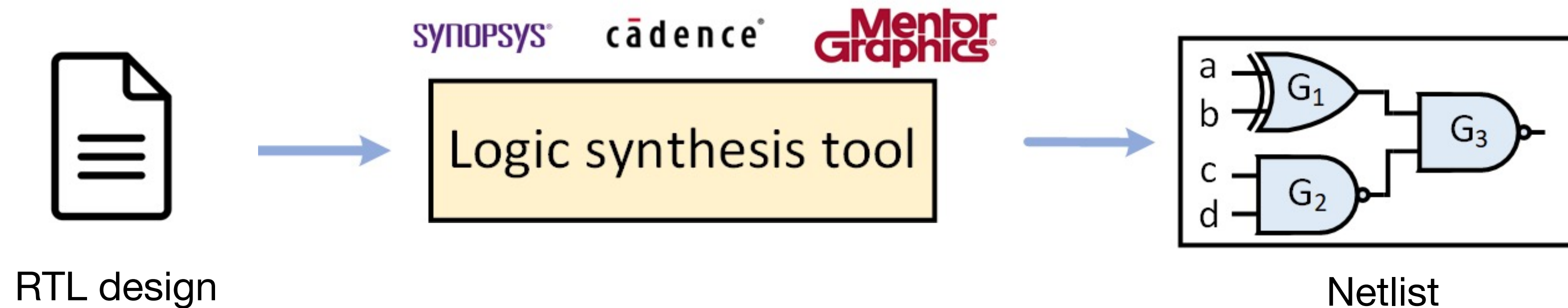
[1]. Yasin et al., "Provably-secure logic locking: From theory to practice." CCS 2017.

[2]. Sengupta et al., "Truly stripping functionality for logic locking: A fault-based perspective." TCAD 2020.



# Logic Synthesis v.s. Logic Locking

- Logic synthesis process



- Example of logic synthesis in logic locking

|    |    |    |    |    |    |
|----|----|----|----|----|----|
|    |    | ab |    |    |    |
|    |    | 00 | 01 | 11 | 10 |
| cd | 00 |    |    | 1  |    |
|    | 01 |    |    | 1  |    |
|    | 11 | 1  | 1  | 1  | 1  |
|    | 10 |    |    | 1  |    |

K-map of original circuit

|    |    |    |    |    |    |
|----|----|----|----|----|----|
|    |    | ab |    |    |    |
|    |    | 00 | 01 | 11 | 10 |
| cd | 00 | 1  |    | 1  |    |
|    | 01 |    |    | 1  |    |
|    | 11 | 1  | 1  | 1  | 1  |
|    | 10 |    |    | 1  |    |

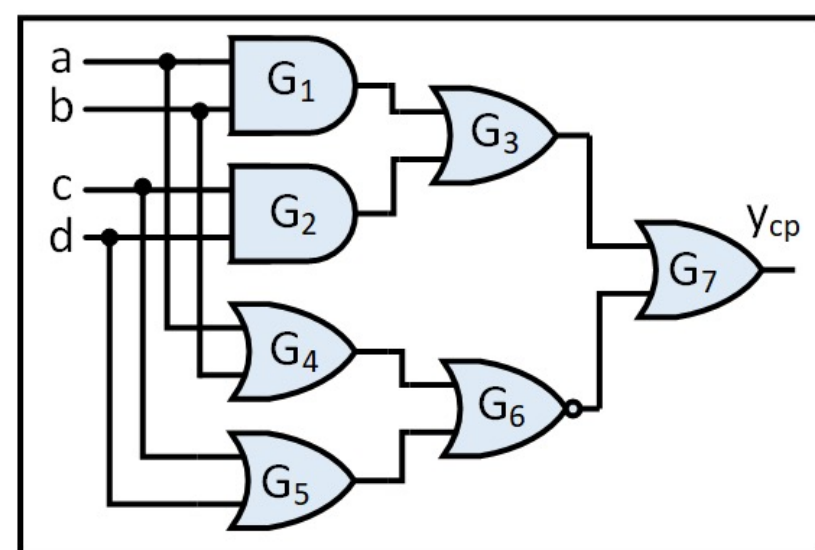
Adding a minterm, 0000

|    |    |    |    |              |    |
|----|----|----|----|--------------|----|
|    |    | ab |    |              |    |
|    |    | 00 | 01 | 11           | 10 |
| cd | 00 |    |    | 1            |    |
|    | 01 |    |    | 1            |    |
|    | 11 | 1  | 1  | <del>1</del> | 1  |
|    | 10 |    |    | 1            |    |

Removing a minterm, 1111

# Sparse Prime Implicant (SPI) Attack

- Using prime implicant table (PIT) to search for PIPs
  - Implicant: A cube that only covers ON-set minterms
  - Prime implicant: The implicant cannot be covered by any other implicant
  - Sparse Prime implicant: PIs are "far" away in rest of PIs in PIT
- SPI attack process



Corrupted circuit

| a | b | c | d | $y_{cp}$ |
|---|---|---|---|----------|
| 0 | 0 | 0 | 0 | 1        |
| 1 | 1 | - | - | 1        |
| - | - | 1 | 1 | 1        |

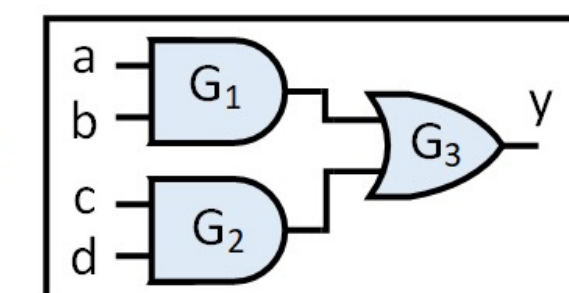
Corrupted circuit's PIT

| a | b | c | d |
|---|---|---|---|
| 0 | 0 | 0 | 0 |

Verified PIP

| a | b | c | d | y |
|---|---|---|---|---|
| 1 | 1 | - | - | 1 |
| - | - | 1 | 1 | 1 |

Original circuit's PIT



Original circuit

# Results

| Circuit \ Attack | Competition-small |     |     |     |     |     | Competition-large |     |     |
|------------------|-------------------|-----|-----|-----|-----|-----|-------------------|-----|-----|
|                  | b10               | b11 | b12 | b13 | b15 | b17 | b17L              | b20 | b22 |
| SAT              | ✓                 | ✓   | ✓   | ✓   | ✓   | X   | X                 | X   | X   |
| AppSAT           | X                 | X   | X   | ✓   | ✓   | X   | X                 | X   | X   |
| ATR, SPS, FALL   | X                 | X   | X   | X   | X   | X   | X                 | X   | X   |
| SPI (proposed)   | ✓                 | ✓   | ✓   | ✓   | ✓   | ✓   | ✓                 | ✓   | ✓   |

✓ denotes a successful attack, X denotes an unsuccessful attack

- CAC-rem locked circuit are from CSAW'19 logic locking competition
  - No one broke it during the competition
- SPI attack breaks all the locked circuits



# More Details (in the paper)

- Scalability of SPI attacks
- What makes SPI attacks hard?
  - PIPs that are far away from PITs of corrupted circuit (aka D2PIPs)
- Conventional benchmark circuits have only few D2PIPs (<100)
  - They are not secure; [should we even use them for logic locking research?](#)



## Future work

- Encode Boolean circuits such that # of D2PIPs are increased
- Tradeoff between security and overhead

I have  
a secret  
plan

# Thank you!

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